

IBM XL C Enterprise Edition for AIX



Compiler Reference

Version 7.0

IBM XL C Enterprise Edition for AIX



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Version 7.0

Before using this information and the product it supports, be sure to read the information in "Notices" on page 395.

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This edition applies to Version 7 Release 0 of IBM XL C Enterprise Edition for AIX (Program Number 5724-I10) and to all subsequent releases and modifications until otherwise indicated in new editions.

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How to read syntax diagrams

This section describes how to read syntax diagrams. It defines syntax diagram symbols, items that may be contained within the diagrams (keywords, variables, delimiters, operators, fragment references, operands) and provides syntax examples that contain these items.

Syntax diagrams pictorially display the order and parts (options and arguments) that comprise a command statement. They are read from left to right and from top to bottom, following the main path of the horizontal line.

Symbols

The following symbols may be displayed in syntax diagrams:

Symbol	Definition
▶—	Indicates the beginning of the syntax diagram.
—→	Indicates that the syntax diagram is continued to the next line.
▶—	Indicates that the syntax is continued from the previous line.
—▶	Indicates the end of the syntax diagram.

Syntax items

Syntax diagrams contain many different items. Syntax items include:

- Keywords - a command name or any other literal information.
- Variables - variables are italicized, appear in lowercase and represent the name of values you can supply.
- Delimiters - delimiters indicate the start or end of keywords, variables, or operators. For example, a left parenthesis is a delimiter.
- Operators - operators include add (+), subtract (-), multiply (*), divide (/), equal (=), and other mathematical operations that may need to be performed.
- Fragment references - a part of a syntax diagram, separated from the diagram to show greater detail.
- Separators - a separator separates keywords, variables or operators. For example, a comma (,) is a separator.

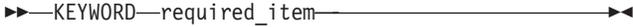
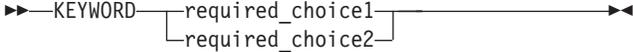
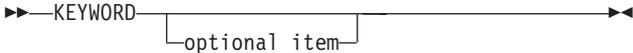
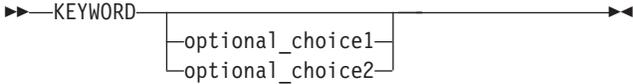
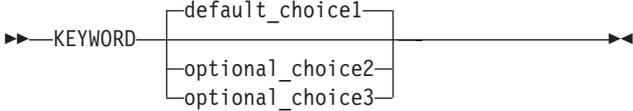
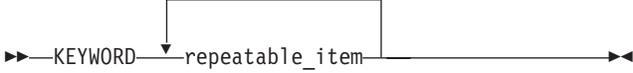
Keywords, variables, and operators may be displayed as required, optional, or default. Fragments, separators, and delimiters may be displayed as required or optional.

Item type	Definition
Required	Required items are displayed on the main path of the horizontal line.
Optional	Optional items are displayed below the main path of the horizontal line.
Default	Default items are displayed above the main path of the horizontal line.

Syntax examples

The following table provides syntax examples.

Table 1. Syntax examples

Item	Syntax example
Required item.	
Required items appear on the main path of the horizontal line. You must specify these items.	
Required choice.	
A required choice (two or more items) appears in a vertical stack on the main path of the horizontal line. You must choose one of the items in the stack.	
Optional item.	
Optional items appear below the main path of the horizontal line.	
Optional choice.	
A optional choice (two or more items) appear in a vertical stack below the main path of the horizontal line. You may choose one of the items in the stack.	
Default.	
Default items appear above the main path of the horizontal line. The remaining items (required or optional) appear on (required) or below (optional) the main path of the horizontal line. The following example displays a default with optional items.	
Variable.	
Variables appear in lowercase italics. They represent names or values.	
Repeatable item.	
An arrow returning to the left above the main path of the horizontal line indicates an item that can be repeated.	
An arrow returning to the left above a group of repeatable items indicates that one of the items can be selected, or a single item can be repeated.	

Part 1. Overview

XL C Enterprise Edition

You can use IBM XL C Enterprise Edition for AIX to compile C program source files into executable object modules.

Note: Throughout these pages, the `xlc` command invocation is used to describe the actions of the compiler. You can, however, substitute other forms of the compiler invocation command if your particular environment requires it, and compiler option usage will remain the same unless otherwise specified.

For more information about the XL C Enterprise Edition compiler, see the following topics in this section:

- “Compiler Modes”
- “Compiler Options” on page 5
- “Types of Input Files” on page 6
- “Types of Output Files” on page 7
- “Compiler Message and Listing Information” on page 8

Compiler Modes

Different forms of XL C Enterprise Edition compiler invocation commands support various levels of the C languages.

In most cases, you should use the `xlc` command to compile C source files.

You can use other forms of the command if your particular environment requires it. The various compiler invocation commands are:

Compiler Invocations	
Basic	Special
<code>xlc</code>	<code>xlc_r xlc128_r4 xlc128_r7 xlc128 xlc128_r xlc128_r4 xlc128_r7</code>
<code>cc</code>	<code>cc_r cc_r4 cc_r7 cc128 cc128_r cc128_r4 cc128_r7</code>
<code>c99</code>	<code>c99_r c99_r4 c99_r7 c99_128 c99_128_r c99_128_r4 c99_128_r7</code>
<code>c89</code>	<code>c89_r c89_r4 c89_r7 c89_128 c89_128_r c89_128_r4 c89_128_r7</code>

The basic compiler invocation commands appear as the first entry of each line above. Select a basic invocation using the following criteria:

xlc	<p>Invokes the compiler for C source files. The following compiler options are implied with this invocation:</p> <ul style="list-style-type: none"> • -qlanglvl=extc89 • -qalias=ansi • -qcpluscmt • -qkeyword=inline
cc	<p>Invokes the compiler for C source files. The following compiler options are implied with this invocation:</p> <ul style="list-style-type: none"> • -qlanglvl=extended • -qnoro • -qnoroconst
c99	<p>Invokes the compiler for C source files, with support for ISO C99 language features. Full ISO C99 conformance requires the presence of C99-compliant header files and run-time libraries. The following options are implied with this invocation:</p> <ul style="list-style-type: none"> • -qlanglvl=stdc99 • -qalias=ansi • -qstrict_induction • -D_ANSI_C_SOURCE • -D_ISOC99_SOURCE
c89	<p>Invokes the compiler for C source files, with support for ISO C89 language features. The following options are implied with this invocation:</p> <ul style="list-style-type: none"> • -qlanglvl=stdc89 • -qalias=ansi • -qstrict_induction • -qnolonglong • -D_ANSI_C_SOURCE <p>Use this invocation for strict conformance to the ANSI standard (ISO/IEC 9899:1990).</p>

IBM XL C Enterprise Edition for AIX provides variations on the basic compiler invocations. These variations are described below:

128-suffixed Invocations All **128**-suffixed invocation commands are functionally similar to their corresponding base compiler invocations. They specify the **-qldb128** option, which increases the length of **long double** types in your program from 64 to 128 bits. They also link with the 128 versions of the C run-times.

_r-suffixed Invocations

All `_r`-suffixed invocations additionally define the macro names `__THREAD_SAFE` and `__VACPP_MULTI__`, and add the libraries `-lc` and `-lpthreads`. The compiler option `-qthreaded` is also added. Use these commands if you want to create Posix threaded applications.

The `_r7` invocations are provided to help migrate programs based on Posix Draft 7 to Posix Draft 10. See `-qthreaded` for additional information. The `_r4` invocations should be used for DCE threaded applications.

Migrating AIX® Version 3.2.5 DCE Applications to AIX Version 4.3.3 and higher

The main invocation commands (except `c89`) have additional `_r4`-suffixed forms. These forms provide compatibility between DCE applications written for AIX Version 3.2.5 and AIX Version 4. They link your application to the correct AIX Version 4 DCE libraries, providing compatibility between the latest version of the `pthreads` library and the earlier versions supported on AIX Version 3.2.5.

Related Tasks

“Invoke the Compiler” on page 27

Related References

“Compiler Command Line Options” on page 51

“General Purpose Pragmas” on page 271

“Pragmas to Control Parallel Processing” on page 321

“threaded” on page 249

Compiler Options

Compiler options perform a wide variety of functions, such as setting compiler characteristics, describing the object code and compiler output to be produced, and performing some preprocessor functions. You can specify compiler options in one or more of three ways:

- on the command line
- in a configuration file (`.cfg`)
- in your source program

The compiler assumes default settings for most compiler options not explicitly set by you in the ways listed above.

When specifying compiler options, it is possible for option conflicts and incompatibilities to occur. IBM® XL C Enterprise Edition resolves most of these conflicts and incompatibilities in a consistent fashion, as follows:

In most cases, the compiler uses the following order in resolving conflicting or incompatible options:

1. Pragma statements in source code will override compiler options specified on the command line.
2. Compiler options specified on the command line will override compiler options specified in a configuration file. If conflicting or incompatible compiler options are specified in the same command line compiler invocation, the option appearing later in the invocation takes precedence.
3. Compiler options specified in a configuration file will override compiler default settings.

Option conflicts that do not follow this priority sequence are described in “Resolving Conflicting Compiler Options” on page 35.

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options” on page 29

“Resolving Conflicting Compiler Options” on page 35

Related References

“Compiler Command Line Options” on page 51

“General Purpose Pragmas” on page 271

“Pragmas to Control Parallel Processing” on page 321

Types of Input Files

You can input the following types of files to the XL C Enterprise Edition compiler:

C Source Files These are files containing C source code.

To use the C compiler to compile a C language source file, the source file must have a `.c` (lowercase c) suffix, for example, `mysource.c`.

The compiler will also accept source files with the `.i` suffix. This extension designates preprocessed source files.

The compiler processes the source files in the order in which they appear. If the compiler cannot find a specified source file, it produces an error message and the compiler proceeds to the next specified file. However, the link editor will not be run and temporary object files will be removed.

Your program can consist of several source files. All of these source files can be compiled at once using only one invocation of the compiler. Although more than one source file can be compiled using a single invocation of the compiler, you can specify only one set of compiler options on the command line per invocation. Each distinct set of command-line compiler options that you want to specify requires a separate invocation.

By default, the compiler preprocesses and compiles all the specified source files. Although you will usually want to use this default, you can use the compiler to preprocess the source file without compiling by specifying either the `-E` or the `-P` option. If you specify the `-P` option, a preprocessed source file, `file_name.i`, is created and processing ends.

The `-E` option preprocesses the source file, writes to standard output, and halts processing without generating an output file.

Preprocessed Source Files

Preprocessed source files have a `.i` suffix, for example, `file_name.i`. The compiler sends the preprocessed source file, `file_name.i`, to the compiler where it is preprocessed again in the same way as a `.c` file. Preprocessed files are useful for checking macros and preprocessor directives.

Object Files

Object files must have a `.o` suffix, for example, `file_name.o`. Object files, library files, and nonstripped executable files serve as input to the linkage editor. After compilation, the linkage editor links all of the specified object files to create an executable file.

Assembler Files	Assembler files must have a .s suffix, for example, <i>file_name.s</i> . Assembler files are assembled to create an object file.
Shared Library Files	Shared library files must have a .so suffix, for example <i>file_name.so</i> .
Nonstripped Executable Files	Extended Common Object File Format (XCOFF) files that have not been stripped with the AIX strip command can be used as input to the compiler. See the strip command in the <i>AIX Commands Reference</i> and the description of a.out file format in the <i>AIX Files Reference</i> for more information.

Related Concepts

“Types of Output Files”

Related References

“E” on page 105

“P” on page 199

See:

strip command in Commands Reference, Volume 5: s through u
Files Reference

Types of Output Files

You can specify the following types of output files when invoking the XL C Enterprise Edition compiler.

Executable Files By default, executable files are named **a.out**. To name the executable file something else, use the **-ofile_name** option with the invocation command. This option creates an executable file with the name you specify as *file_name*. The name you specify can be a relative or absolute path name for the executable file.

The format of the **a.out** file is described in the *AIX Files Reference*.

Object Files The compiler gives object files a **.o** suffix, for example, *file_name.o*, unless the **-ofile_name** option is specified giving a different suffix or no suffix at all.

If you specify the **-c** option, an output object file, *file_name.o*, is produced for each input source file *file_name.x*, where *x* is a recognized C filename extension. The linkage editor is not invoked, and the object files are placed in your current directory. All processing stops at the completion of the compilation.

You can link-edit the object files later into a single executable file by invoking the compiler.

Assembler Files Assembler files must have a **.s** suffix, for example, *file_name.s*.

They are created by specifying the **-S** option. Assembler files are assembled to create an object file.

Preprocessed Source Files	<p>Preprocessed source files have a <code>.i</code> suffix, for example, <code>file_name.i</code>.</p> <p>To make a preprocessed source file, specify the <code>-P</code> option. The source files are preprocessed but not compiled. You can also redirect the output from the <code>-E</code> option to generate a preprocessed file that contains <code>#line</code> directives.</p> <p>A preprocessed source file, <code>file_name.i</code>, is produced for each source file and has the same file name (with a <code>.i</code> extension) as the source file from which it was produced.</p>
Listing Files	<p>Listing files have a <code>.lst</code> suffix, for example, <code>file_name.lst</code>.</p> <p>Specifying any one of the listing-related options to the invocation command produces a compiler listing (unless you have specified the <code>-qnoprint</code> option). The file containing this listing is placed in your current directory and has the same file name (with a <code>.lst</code> extension) as the source file from which it was produced.</p>
Shared Library Files	<p>Shared library files have a <code>.so</code> suffix, for example, <code>my_shrplib.so</code>.</p>
Target Files	<p>Output files associated with the <code>-M</code> or <code>-qmakedep</code> options have a <code>.u</code> suffix, for example, <code>conversion.u</code>.</p> <p>The file contains targets suitable for inclusion in a description file for the <code>make</code> command.</p>

Related Concepts

“Types of Input Files” on page 6

Related References

“c” on page 84
 “E” on page 105
 “M” on page 180
 “makedep” on page 186
 “o” on page 198
 “P” on page 199
 “print” on page 210
 “S” on page 225

Related References

“c” on page 84
 “E” on page 105
 “M” on page 180
 “makedep” on page 186
 “o” on page 198
 “P” on page 199
 “print” on page 210
 “S” on page 225

Compiler Message and Listing Information

When the compiler encounters a programming error while compiling a C source program, it issues a diagnostic message to the standard error device and if the appropriate options have been selected, to the listing file.

Compiler Messages

The compiler issues messages specific to the C language.

If you specify the compiler option **-qsrcmsg** and the error is applicable to a particular line of code, the reconstructed source line or partial source line is included with the error message in the stderr file. A reconstructed source line is a preprocessed source line that has all the macros expanded.

If you specify the **-qsource** compiler option, the compiler will place messages in the source listing. For example, if you compile your file using the command line invocation **xlC -qsource filename.c**, then you will find a file called **filename.lst** in your current directory.

You can control the diagnostic messages issued, according to their severity, using either the **-qflag** option or the **-w** option. To get additional informational messages about potential problems in your program, use the **-qinfo** option.

Compiler Listings

The listings produced by the compiler are useful for debugging. By specifying appropriate options, you can request information on all aspects of a compilation. The listing consists of a combination of the following sections:

- Header section that lists the compiler name, version, and release, as well as the source file name and the date and time of the compilation
- Source section that lists the input source code with line numbers. If there is an error at a line, the associated error message appears after the source line.
- Options section that lists the options that were in effect during the compilation
- Attribute and cross-reference listing section that provides information about the variables used in the compilation unit
- File table section that shows the file number and file name for each main source file and include file
- Compilation epilogue section that summarizes the diagnostic messages, lists the number of source lines read, and indicates whether the compilation was successful
- Object section that lists the object code

Each section, except the header section, has a section heading that identifies it. The section heading is enclosed by angle brackets.

Related References

“Message Severity Levels and Compiler Response” on page 355

“Compiler Message Format” on page 356

“flag” on page 117

“info” on page 142

“source” on page 234

“srcmsg” on page 238

“w” on page 264

Program Parallelization

The compiler offers you three methods of implementing shared memory program parallelization. These are:

- Automatic parallelization of program loops.
- Explicit parallelization of C program code using IBM SMP pragma directives.
- Explicit parallelization of C program code using pragma directives compliant to the **OpenMP Application Program Interface** specification.

All methods of program parallelization are enabled when the **-qsmp** compiler option is in effect without the **omp** suboption. You can enable strict OpenMP compliance with the **-qsmp=omp** compiler option, but doing so will disable automatic parallelization.

Parallel regions of program code are executed by multiple threads, possibly running on multiple processors. The number of threads created is determined by the run-time options and calls to library functions. Work is distributed among available threads according to the specified scheduling algorithm. For a complete discussion on how threads are created and utilized, refer to the OpenMP Specification, section 2.3 Parallel Construct.

Note: The **-qsmp** option must only be used together with thread-safe compiler invocation modes.

For more information about parallel programming support offered by the XL C Enterprise Edition compiler, see the following topics in this section:

- “IBM SMP Directives”
- “OpenMP Directives” on page 12
- “Countable Loops” on page 13
- “Reduction Operations in Parallelized Loops” on page 14
- “Shared and Private Variables in a Parallel Environment” on page 15

For complete information about the OpenMP Specification, see:

- OpenMP Web site (www.openmp.org)
- OpenMP Specification (www.openmp.org/specs)

IBM SMP Directives

C IBM SMP directives exploit shared memory parallelism through the parallelization of *countable loops*. A loop is considered to be *countable* if it has any of the forms described in *Countable Loops*.

The compiler can automatically locate and where possible parallelize all countable loops in your program code. In general, a countable loop is automatically parallelized only if all of the follow conditions are met:

- the order in which loop iterations start or end does not affect the results of the program.
- the loop does not contain I/O operations.
- floating point reductions inside the loop are not affected by round-off error, unless the **-qnostrict** option is in effect.

- the `-qnostrict_induction` compiler option is in effect.
- the `-qsmp` compiler option is in effect without the `omp` suboption. The compiler must be invoked using a thread-safe compiler mode.

You can also explicitly instruct the compiler to parallelize selected countable loops.

The XL C Enterprise Edition compiler provides pragma directives that you can use to improve on automatic parallelization performed by the compiler. Pragmas fall into two general categories.

1. The first category of pragmas lets you give the compiler information on the characteristics of a specific countable loop. The compiler uses this information to perform more efficient automatic parallelization of the loop.
2. The second category gives you explicit control over parallelization. Use these pragmas to force or suppress parallelization of a loop, apply specific parallelization algorithms to a loop, and synchronize access to shared variables using critical sections.

Related Concepts

“OpenMP Directives”

“Countable Loops” on page 13

“Reduction Operations in Parallelized Loops” on page 14

“Shared and Private Variables in a Parallel Environment” on page 15

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

“Control Parallel Processing with Pragmas” on page 39

Related References

“smp” on page 232

“Pragmas to Control Parallel Processing” on page 321

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

For complete information about the OpenMP Specification, see:

- OpenMP Web site (www.openmp.org)
- OpenMP Specification (www.openmp.org/specs)

OpenMP Directives

C **C++** OpenMP directives exploit shared memory parallelism by defining various types of *parallel regions*. Parallel regions can include both iterative and non-iterative segments of program code.

Pragmas fall into four general categories:

1. The first category of pragmas lets you define parallel regions in which work is done by threads in parallel. Most of the OpenMP directives either statically or dynamically bind to an enclosing parallel region.
2. The second category lets you define how work will be distributed or shared across the threads in a parallel region.
3. The third category lets you control synchronization among threads.
4. The fourth category lets you define the scope of data visibility across threads.

Related Concepts

“IBM SMP Directives” on page 11

“Countable Loops”

“Reduction Operations in Parallelized Loops” on page 14

“Shared and Private Variables in a Parallel Environment” on page 15

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

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Related References

“smp” on page 232

“Pragmas to Control Parallel Processing” on page 321

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

See also Parallel Programming in General Programming Concepts: Writing and Debugging Programs.

For complete information about the OpenMP Specification, see:

- OpenMP Web site (www.openmp.org)
- OpenMP Specification (www.openmp.org/specs)

Countable Loops

A loop is considered to be *countable* if it has any of the forms shown below, and:

- there is no branching into or out of the loop.
- the *incr_expr* expression is not within a critical section.

The following are examples of countable loops.

```
for ([iv]; exit_cond; incr_expr)
    statement

for ([iv]; exit_cond; [expr]) {
    [declaration_list]
    [statement_list]
    incr_expr;
    [statement_list]
}

while (exit_cond) {
    [declaration_list]
    [statement_list]
    incr_expr;
    [statement_list]
}

do {
    [declaration_list]
    [statement_list]
    incr_expr;
    [statement_list]
} while (exit_cond)
```

The following definitions apply to the above examples:

<i>exit_cond</i>	takes	<i>iv</i> <= <i>ub</i>
	form:	<i>iv</i> < <i>ub</i>
		<i>iv</i> >= <i>ub</i>
		<i>iv</i> > <i>ub</i>
<i>incr_expr</i>	takes	<i>++iv</i>
	form:	<i>iv++</i>
		<i>--iv</i>
		<i>iv--</i>
		<i>iv += incr</i>
		<i>iv -= incr</i>
		<i>iv = iv + incr</i>
		<i>iv = incr + iv</i>
		<i>iv = iv - incr</i>

iv Iteration variable. The iteration variable is a signed integer that has either automatic or register storage class, does not have its address taken, and is not modified anywhere in the loop except in *incr_expr*.

incr Loop invariant signed integer expression. The value of the expression is known at run-time and is not 0. *incr* cannot reference extern or static variables, pointers or pointer expressions, function calls, or variables that have their address taken.

ub Loop invariant signed integer expression. *ub* cannot reference extern or static variables, pointers or pointer expressions, function calls, or variables that have their address taken.

Related Concepts

“IBM SMP Directives” on page 11
“OpenMP Directives” on page 12
“Reduction Operations in Parallelized Loops”
“Shared and Private Variables in a Parallel Environment” on page 15

Related Tasks

“Set Parallel Processing Run-time Options” on page 24
“Control Parallel Processing with Pragmas” on page 39

Related References

“smp” on page 232
“Pragmas to Control Parallel Processing” on page 321
“IBM SMP Run-time Options for Parallel Processing” on page 359
“OpenMP Run-time Options for Parallel Processing” on page 362
“Built-in Functions Used for Parallel Processing” on page 364

Reduction Operations in Parallelized Loops

The compiler can recognize and properly handle most reduction operations in a loop during both automatic and explicit parallelization. In particular, it can handle reduction statements that have either of the following forms:

```
var = var op expr;  
var assign_op expr;
```

where:

<i>var</i>	Is an identifier designating an automatic or register variable that does not have its address taken and is not referenced anywhere else in the loop, including all loops that are nested. For example, in the following code, only S in the nested loop is recognized as a reduction:
	<pre>int i,j, S=0; #pragma ibm parallel_loop for (i= 0 ;i < N; i++) { S = S+ i; #pragma ibm parallel_loop for (j=0;j< M; j++) { S = S + j; } }</pre>
<i>op</i>	Is one of the following operators: + - * ^ &
<i>assign_op</i>	Is one of the following operators: += -= *= ^= = &=
<i>expr</i>	Is any valid expression.

Recognized reductions are listed by the **-qinfo=reduction** option. When using IBM directives, use critical sections to synchronize access to all reduction variables not recognized by the compiler. OpenMP directives provide you with mechanisms to specify reduction variables explicitly.

Shared and Private Variables in a Parallel Environment

Variables can have either shared or private context in a parallel environment.

- Variables in shared context are visible to all threads running in associated parallel loops.
- Variables in private context are hidden from other threads. Each thread has its own private copy of the variable, and modifications made by a thread to its copy are not visible to other threads.

The default context of a variable is determined by the following rules:

- Variables with **static** storage duration are shared.
- Dynamically allocated objects are shared.
- Variables with automatic storage duration are private.
- Variables in heap allocated memory are shared. There can be only one shared heap.
- All variables defined outside a parallel construct become shared when the parallel loop is encountered.
- Loop iteration variables are private within their loops. The value of the iteration variable after the loop is the same as if the loop were run sequentially.
- Memory allocated within a parallel loop by the **alloca** function persists only for the duration of one iteration of that loop, and is private for each thread.

The following code segments show examples of these default rules:

```

int E1;                                /* shared static */
void main (argc,...) {                 /* argc is shared */
  int i;                                /* shared automatic */
  void *p = malloc(...);              /* memory allocated by malloc */
                                      /* is accessible by all threads */
                                      /* and cannot be privatized */

#pragma omp parallel firstprivate (p)
{
  int b;                                /* private automatic */
  static int s;                        /* shared static */

#pragma omp for
  for (i =0;...) {
    b = 1;                              /* b is still private here ! */
    foo (i);                            /* i is private here because it */
                                        /* is an iteration variable */
  }

#pragma omp parallel
  {
    b = 1;                              /* b is shared here because it */
                                        /* is another parallel region */
  }
}

int E2;                                /*shared static */
void foo (int x) {                    /* x is private for the parallel */
                                      /* region it was called from */

  int c;                                /* the same */
  ... }

```

The compiler can privatize some shared variables if it is possible to do so without changing the semantics of the program. For example, if each loop iteration uses a unique value of a shared variable, that variable can be privatized. Privatized shared variables are reported by the **-qinfo=private** option. Use critical sections to synchronize access to all shared variables not listed in this report.

Some OpenMP preprocessor directives let you specify visibility context for selected data variables. A brief summary of data scope attribute clauses are listed below:

Data Scope Attribute Clause	Description
private	The private clause declares the variables in the list to be private to each thread in a team.
firstprivate	The firstprivate clause provides a superset of the functionality provided by the private clause.
lastprivate	The lastprivate clause provides a superset of the functionality provided by the private clause.
shared	The shared clause shares variables that appear in the list among all the threads in a team. All threads within a team access the same storage area for shared variables.
reduction	The reduction clause performs a reduction on the scalar variables that appear in the list, with a specified operator.
default	The default clause allows the user to affect the data scope attributes of variables.

For more information, see OpenMP directive descriptions or the OpenMP C and C++ Application Program Interface specification.

Related Concepts

“IBM SMP Directives” on page 11

“OpenMP Directives” on page 12

“Countable Loops” on page 13

“Reduction Operations in Parallelized Loops” on page 14

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

“Control Parallel Processing with Pragmas” on page 39

Related References

“smp” on page 232

“Pragmas to Control Parallel Processing” on page 321

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

For complete information about the OpenMP Specification, see:

- OpenMP Web site at www.openmp.org
- OpenMP Specification at www.openmp.org/specs

Using XL C Enterprise Edition with Other Programming Languages

With XL C Enterprise Edition, you can call functions written in other XL languages from your C and C++ programs. Similarly, the other XL language programs can call functions written in C. You should already be familiar with the syntax of the languages you are using.

See “Use C with Other Programming Languages” on page 41 for more information.

Part 2. Configuration and Use

Set Up the Compilation Environment

Before you compile your C programs, you must set up the environment variables and the configuration file for your system.

For information about environment variables used by the XL C Enterprise Edition compiler, see the following topics in this section:

- “Set Environment Variables”
- “Set Environment Variables to Select 64- or 32-bit Modes” on page 24
- “Set Parallel Processing Run-time Options” on page 24
- “Set Other Environment Variables” on page 24

Set Environment Variables

You use different commands to set the environment variables depending on whether you are using the Bourne shell (**bsh** or **sh**), Korn shell (**ksh**), or C shell (**csh**). To determine the current shell, use the **echo** command:

```
echo $SHELL
```

The Bourne-shell path is **/bin/bsh** or **/bin/sh**. The Korn shell path is **/bin/ksh**. The C-shell path is **/bin/csh**.

For more information about the **NLSPATH** and **LANG** environment variables, see *System User's Guide: Operating System and Devices*. The AIX international language facilities are described in the *AIX General Programming Concepts*.

Set Environment Variables in bsh, ksh, or sh Shells

The following statements, either typed at the command line or inserted into a command file script, show how you can set environment variables in the Bourne or Korn shells. Paths shown assume that you are installing the compiler in the default installation location.

```
LANG=en_US
NLSPATH=/usr/lib/nls/msg/%L/%N:/usr/lib/nls/msg/L/%N
export LANG NLSPATH
```

To set the variables so that all users have access to them, add the commands to the file **/etc/profile**. To set them for a specific user only, add the commands to the file **.profile** in the user's home directory. The environment variables are set each time the user logs in.

Set Environment Variables in csh Shell

The following statements show how you can set environment variables in the csh shell. Paths shown assume that you are installing the compiler in the default installation location.

```
setenv LANG en_US
setenv NLSPATH /usr/lib/nls/msg/%L/%N:/usr/lib/nls/msg/L/%N
```

In the csh shell, you cannot set the environment variables so that all users have access to them. To set them for a specific user only, add the commands to the file **.cshrc** in the user's home directory. The environment variables are set each time the user logs in.

Related Tasks

“Set Environment Variables to Select 64- or 32-bit Modes”
“Set Parallel Processing Run-time Options”
“Set Other Environment Variables”

Related References

See also:

System User’s Guide: Operating System and Devices
AIX 5L™ Version 5.1 General Programming Concepts: Writing and Debugging Programs

Set Environment Variables to Select 64- or 32-bit Modes

You can set the OBJECT_MODE environment variable to specify a default compilation mode. Permissible values for the OBJECT_MODE environment variable are:

<i>(unset)</i>	Compiler programs generate and/or use 32-bit objects.
32	Compiler programs generate and/or use 32-bit objects.
64	Compiler programs generate and/or use 64-bit objects.
32_64	Set the compiler programs to accept both 32- and 64-bit objects. However, you should not use this value. The compiler never functions in this mode, and using this choice may generate an error message depending on other compilation options set at compile-time.

Related Tasks

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33
“Set Environment Variables” on page 23

Set Parallel Processing Run-time Options

The XLSMPOPTS environment variable sets options for programs using loop parallelization. For example, to have a program run-time create 4 threads and use dynamic scheduling with chunk size of 5, you would set the XLSMPOPTS environment variable as shown below:

```
XLSMPOPTS=PARTHDS=4:SCHEDULE=DYNAMIC=5
```

Additional environment variables set options for program parallelization using OpenMP-compliant directives.

Related Tasks

“Set Environment Variables” on page 23

Related References

“IBM SMP Run-time Options for Parallel Processing” on page 359
“OpenMP Run-time Options for Parallel Processing” on page 362

Set Other Environment Variables

Before using the compiler, ensure the following environment variables are set.

PATH	Specifies the directory search path for the executable files of the compiler.
-------------	---

MANPATH	Optionally specifies the directory search path for finding man pages. MANPATH must contain before the default man path.
OBJECT_MODE	Optionally specifies the default compilation mode to be either 32 or 64 bits.
LANG	Specifies the national language for message and help files. The LANG environment variable can be set to any of the locales provided on the system. See the description of locales in <i>AIX General Programming Concepts</i> for more information. The national language code for United States English is en_US . If the appropriate message catalogs have been installed on your system, any other valid national language code can be substituted for en_US . To determine the current setting of the national language on your system, use the both of the following echo commands: <pre>echo \$LANG echo \$NLSPATH</pre>
NLSPATH	Specifies the path name of the message and help files.
PDFDIR	Optionally specifies the directory in which the profile data file is created. The default value is unset, and the compiler places the profile data file in the current working directory. Setting this variable to an absolute path is recommended for profile-directed feedback.
TMPDIR	Optionally specifies the directory in which temporary files are created. The default location, /tmp, may be inadequate at high levels of optimization, where paging and temporary files can require significant amounts of disk space.

Note: The **LANG** and **NLSPATH** environment variables are initialized when the operating system is installed, and might differ from the ones you want to use.

Related Tasks

“Set Environment Variables” on page 23

Related References

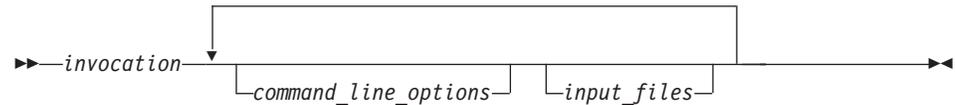
“Set Environment Variables to Select 64- or 32-bit Modes” on page 24

See also:

AIX 5L Version 5.1 General Programming Concepts: Writing and Debugging Programs

Invoke the Compiler

The IBM XL C Enterprise Edition compiler is invoked using the following syntax, where *invocation* can be replaced with any valid XL C Enterprise Edition invocation command:



The parameters of the compiler invocation command can be the names of input files, compiler options, and linkage-editor options.

Compiler options perform a wide variety of functions, such as setting compiler characteristics, describing the object code and compiler output to be produced, and performing some preprocessor functions.

By default, the invocation command calls *both* the compiler and the linkage editor. It passes linkage editor options to the linkage editor. Consequently, the invocation commands also accept all linkage editor options.

To compile without link-editing, use the **-c** compiler option. The **-c** option stops the compiler after compilation is completed and produces as output, an object file *file_name.o* for each *file_name.c* input source file, unless the **-o** option was used to specify a different object filename. The linkage editor is not invoked. You can link-edit the object files later using the same invocation command, specifying the object files without the **-c** option.

Invoke the Linkage Editor

The linkage editor link-edits specified object files to create one executable file. Invoking the compiler with one of the invocation commands automatically calls the linkage editor unless you specify one of the following compiler options: **-E**, **-P**, **-c**, **-S**, **-qsyntaxonly** or **-#**.

Input Files

Object files, unstripped executable files, and library files serve as input to the linkage editor. Object files must have a **.o** suffix, for example, **year.o**. Static library file names have a **.a** suffix, for example, **libold.a**. Dynamic library file names have a **.so** suffix, for example, **libold.so**.

Library files are created by combining one or more files into a single archive file with the AIX **ar** command. For a description of the **ar** command, refer to the *AIX Commands Reference*.

Output Files

The linkage editor generates an *executable file* and places it in your current directory. The default name for an executable file is **a.out**. To name the executable file explicitly, use the **-o file_name** option with the compiler invocation command, where *file_name* is the name you want to give to the executable file. For example, to compile **myfile.c** and generate an executable file called **myfile**, enter:

```
xlc myfile.c -o myfile
```

If you use the **-qmkshrobj** option to create a shared library, the default name of the shared object created is **shr.o**.

You can invoke the linkage editor explicitly with the **ld** command. However, the compiler invocation commands set several linkage-editor options, and link some standard files into the executable output by default. In most cases, it is better to use one of the compiler invocation commands to link-edit your object files.

Note: When link-editing object files, *do not* use the **-e** option of the **ld** command. The default entry point of the executable output is `__start`. Changing this label with the **-e** flag can cause erratic results.

Related Concepts

“Compiler Modes” on page 3

Related Tasks

“Specify Compiler Options” on page 29

“Invoke the Compiler” on page 27

Related References

“Compiler Command Line Options” on page 51

“Message Severity Levels and Compiler Response” on page 355

See also:

ld command in Commands Reference, Volume 5: s through u

Specify Compiler Options

You can specify compiler options in any of the following ways:

- On the command line (see page 29)
- In your source program (see page 31)
- In a configuration (.cfg) file (see page 32)

The compiler assumes default settings for most compiler options not explicitly set by you in the ways listed above.

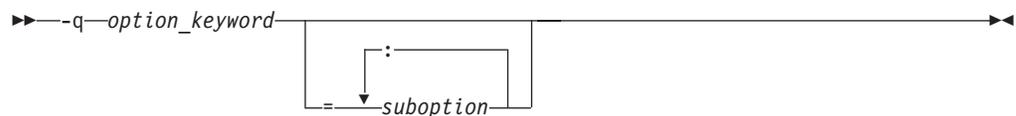
Specify Compiler Options on the Command Line

Most options specified on the command line override both the default settings of the option and options set in the configuration file. Similarly, most options specified on the command line are in turn overridden by pragma directives, which provide you a means of setting compiler options right in the source file. Options that do not follow this scheme are listed in *Resolving Conflicting Compiler Options*.

There are two kinds of command-line options:

- `-qoption_keyword` (compiler-specific)
- Flag options

-q Options



Command-line options in the `-qoption_keyword` format are similar to on and off switches. For *most* `-q` options, if a given option is specified more than once, the last appearance of that option on the command line is the one recognized by the compiler. For example, `-qsource` turns on the source option to produce a compiler listing, and `-qnosource` turns off the source option so no source listing is produced. For example:

```
xlc -qnosource MyFirstProg.c -qsource MyNewProg.c
```

would produce a source listing for both `MyNewProg.c` and `MyFirstProg.c` because the last `source` option specified (`-qsource`) takes precedence.

You can have multiple `-qoption_keyword` instances in the same command line, but they must be separated by blanks. Option keywords can appear in either uppercase or lowercase, but you must specify the `-q` in lowercase. You can specify any `-qoption_keyword` before or after the file name. For example:

```
xlc -qLIST -qfloat=nomaf file.c  
xlc file.c -qxref -qsource
```

You can also abbreviate many compiler options. For example, specifying `-qopt` is equivalent to specifying `-qoptimize` on the command line.

Some options have suboptions. You specify these with an equal sign following the **-qoption**. If the option permits more than one suboption, a colon (:) must separate each suboption from the next. For example:

```
xlc -qflag=w:e -qattr=full file.c
```

compiles the C source file `file.c` using the option **-qflag** to specify the severity level of messages to be reported. The **-qflag** suboption `w` (warning) sets the minimum level of severity to be reported on the listing, and suboption `e` (error) sets the minimum level of severity to be reported on the terminal. The **-qflag** option **-qattr** with suboption `full` will produce an attribute listing of all identifiers in the program.

Flag Options

The compilers available on AIX systems use a number of common conventional flag options. IBM XL C Enterprise Edition supports these flags. Lowercase flags are different from their corresponding uppercase flags. For example, **-c** and **-C** are two different compiler options: **-c** specifies that the compiler should only preprocess and compile and not invoke the linkage editor, while **-C** can be used with **-P** or **-E** to specify that user comments should be preserved.

IBM XL C Enterprise Edition also supports flags directed to other AIX programming tools and utilities (for example, the AIX `ld` command). The compiler passes on those flags directed to `ld` at link-edit time.

Some flag options have arguments that form part of the flag. For example:

```
xlc stem.c -F/home/tools/test3/new.cfg:xlc
```

where `new.cfg` is a custom configuration file.

You can specify flags that do not take arguments in one string. For example:

```
xlc -0cv file.c
```

has the same effect as:

```
xlc -0 -c -v file.c
```

and compiles the C source file `file.c` with optimization (**-O**) and reports on compiler progress (**-v**), but does not invoke the linkage editor (**-c**).

A flag option that takes arguments can be specified as part of a single string, but you can only use one flag that takes arguments, and it must be the last option specified. For example, you can use the **-o** flag (to specify a name for the executable file) together with other flags, only if the **-o** option and its argument are specified last. For example:

```
xlc -0vo test test.c
```

has the same effect as:

```
xlc -0 -v -otest test.c
```

Most flag options are a single letter, but some are two letters. Note that specifying **-pg** (extended profiling) is not the same as specifying **-p -g** (**-p** for profiling, and **-g** for generating debug information). Take care not to specify two or more options in a single string if there is another option that uses that letter combination.

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options in Your Program Source Files”

“Specify Compiler Options in a Configuration File” on page 32

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33

“Resolving Conflicting Compiler Options” on page 35

Related References

“Compiler Command Line Options” on page 51

Specify Compiler Options in Your Program Source Files

You can specify compiler options within your program source by using `#pragma` directives.

A pragma is an implementation-defined instruction to the compiler. It has the general form given below, where *character_sequence* is a series of characters that give specific compiler instruction and arguments, if any.



The *character_sequence* on a pragma is subject to macro substitutions, unless otherwise stated. More than one pragma construct can be specified on a single `#pragma` directive. The compiler ignores unrecognized pragmas, issuing an informational message indicating this.

Options specified with pragma directives in program source files override all other option settings, except other pragma directives. The effect of specifying the same pragma directive more than once varies. See the description for each pragma for specific information.

Pragma settings can carry over into included files. To avoid potential unwanted side-effects from pragma settings, you should consider resetting pragma settings at the point in your program source where the pragma-defined behavior is no longer required. Some pragma options offer **reset** or **pop** suboptions to help you do this.

These **#pragma** directives are listed in the detailed descriptions of the options to which they apply. For complete details on the various **#pragma** preprocessor directives, see *General Purpose Pragmas*.

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options on the Command Line” on page 29

“Specify Compiler Options in a Configuration File” on page 32

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33

“Resolving Conflicting Compiler Options” on page 35

Related References

“General Purpose Pragmas” on page 271

“Pragmas to Control Parallel Processing” on page 321

Specify Compiler Options in a Configuration File

The default configuration file (`/etc/vac.cfg`) defines values and compiler options for the compiler. The compiler refers to this file when compiling C programs. The configuration file is a plain text file, and you can make entries to this file to support specific compilation requirements or to support other C compilation environments.

Most options specified in the configuration file override the default settings of the option. Similarly, most options specified in the configuration file are in turn overridden by options set in the source file and on the command line. Options that do not follow this scheme are listed in *Resolving Conflicting Compiler Options*.

Tailor a Configuration File

The default configuration file is installed to `/etc/vac.cfg`.

You can copy this file and make changes to the copy to support your specific compilation requirements or to support other C compilation environments.

You can modify existing stanzas or add new stanzas to a configuration file. For example, to make `-qnor` the default for the `xl` compiler invocation command, add `-qnor` to the `xl` stanza in your copied version of the configuration file.

You can link the compiler invocation command to several different names. The name you specify when you invoke the compiler determines which stanza of the configuration file the compiler uses. You can add other stanzas to your copy of the configuration file to customize your own compilation environment. You can use the `-F` option with the compiler invocation command to make links to select additional stanzas or to specify a specific stanza in another configuration file. For example:

```
xl myfile.c -Fmyconfig:SPECIAL
```

would compile `myfile.c` using the `SPECIAL` stanza in a `myconfig.cfg` configuration file that you had created.

Configuration File Attributes

A configuration file includes several stanzas. The following are just some of the items defined by stanzas in the configuration file:

<code>as</code>	Path name to be used for the assembler. The default is <code>/bin/as</code> .
<code>asopt</code>	List of options for the assembler and not for the compiler. These override all normal processing by the compiler and are directed to the assembler specified in the <code>as</code> stanza. The string is formatted for the AIX <code>getopt()</code> subroutine as a concatenation of flag letters, with a letter followed by a colon (:) if the corresponding flag takes a parameter.
<code>ccomp</code>	C Front end. The default is <code>/usr/vac/exe/xlcentry</code> .
<code>codeopt</code>	List of options for the code-generation phase of the compiler.
<code>cppcode</code>	Path name to be used for the code generation phase of the compiler. The default is <code>/usr/vac/exe/xlccode</code>
<code>cppopt</code>	List of options for the lexical analysis phase of the compiler.

crt	Path name of the object file passed as the first parameter to the linkage editor. If you do not specify either the -p or the -pg option, the crt value is used. The default is /lib/crt0.o .
csuffix	Suffix for source programs. The default is c (lowercase c).
dis	Path name of the disassembler. The default is /usr/vacpp/exe/dis .
gcr	Path name of the object file passed as the first parameter to the linkage editor. If you specify the -pg option, the gcr value is used. The default is /lib/gcr0.o .
ld	Path name to be used to link C programs. The default is /bin/ld .
ldopt	List of options that are directed to the linkage editor part of the compiler. These override all normal processing by the compiler and are directed to the linkage editor. If the corresponding flag takes a parameter, the string is formatted for the getopt() subroutine as a concatenation of flag letters, with a letter followed by a colon (:).
libraries2	Library options, separated by commas, that the compiler passes as the last parameters to the linkage editor. libraries2 specifies the libraries that the linkage editor is to use at link-edit time for both profiling and nonprofiling. The default is empty.
mcrt	Path name of the object file passed as the first parameter to the linkage editor if you have specified the -p option. The default is /lib/mcrt0.o .
options	A string of option flags, separated by commas, to be processed by the compiler as if they had been entered on the command line.
osuffix	The suffix for object files. The default is .o .
proflibs	Library options, separated by commas, that the compiler passes to the linkage editor when profiling options are specified. proflibs specifies the profiling libraries used by the linkage editor at link-edit time. The default is -L/lib/profiled and -L/usr/lib/profiled .
ssuffix	The suffix for assembler files. The default is .s .
use	Values for attributes are taken from the named stanza and from the local stanza. For single-valued attributes, values in the use stanza apply if no value is provided in the local, or default, stanza. For comma-separated lists, the values from the use stanza are added to the values from the local stanza.
xc	The path name of the xc compiler component. The default is /usr/vac/bin/xc .

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options on the Command Line” on page 29

“Specify Compiler Options in Your Program Source Files” on page 31

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation”

“Resolving Conflicting Compiler Options” on page 35

Related References

“Compiler Command Line Options” on page 51

See also the *Configure the compiler* section in *Getting Started with* .

Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation

You can use IBM XL C Enterprise Edition compiler options to optimize compiler output for use on specific processor architectures. You can also instruct the compiler to compile in either 32- or 64-bit mode.

The compiler evaluates compiler options in the following order, with the last allowable one found determining the compiler mode:

1. Internal default (32-bit mode)
2. OBJECT_MODE environment variable setting, as follows:

OBJECT_MODE Setting	User-selected compilation-mode behavior, unless overridden by configuration file or command-line options
not set	32-bit compiler mode.
32	32-bit compiler mode.
64	64-bit compiler mode.
32_64	Fatal error and stop with following message, 1501-254 OBJECT_MODE=32_64 is not a valid setting for the compiler unless an explicit configuration file or command-line compiler-mode setting exists.
any other	Fatal error and stop with following message, 1501-255 OBJECT_MODE setting is not recognized and is not a valid setting for the compiler unless an explicit configuration file or command-line compiler-mode setting exists.

3. Configuration file settings
4. Command line compiler options (**-q32**, **-q64**, **-qarch**, **-qtune**)
5. Source file statements (**#pragma options tune=suboption**)

The compilation mode actually used by the compiler depends on a combination of the settings of the **-q32**, **-q64**, **-qarch** and **-qtune** compiler options, subject to the following conditions:

- *Compiler mode* is set according to the last-found instance of the **-q32** or **-q64** compiler options. If neither of these compiler options is set, the compiler mode is set by the value of the OBJECT_MODE environment variable. If the OBJECT_MODE environment variable is also not set, the compiler assumes 32-bit compilation mode.
- *Architecture target* is set according to the last-found instance of the **-qarch** compiler option, provided that the specified **-qarch** setting is compatible with the *compiler mode* setting. If the **-qarch** option is not set, the compiler assumes a **-qarch** setting of **com**.
- Tuning of the architecture target is set according to the last-found instance of the **-qtune** compiler option, provided that the **-qtune** setting is compatible with the *architecture target* and *compiler mode* settings. If the **-qtune** option is not set, the compiler assumes a default **-qtune** setting according to the **-qarch** setting in use. If **-qarch** is not specified, the compiler assumes a **-qtune** setting of **pwr3**.

Allowable combinations of these options are found in the **Acceptable Compiler Mode and Processor Architecture Combinations** table.

Possible option conflicts and compiler resolution of these conflicts are described below:

- **-q32** or **-q64** setting is incompatible with user-selected **-qarch** option.
Resolution: **-q32** or **-q64** setting overrides **-qarch** option; compiler issues a warning message, sets **-qarch** to its default setting, and sets the **-qtune** option accordingly to its default value.
- **-q32** or **-q64** setting is incompatible with user-selected **-qtune** option.
Resolution: **-q32** or **-q64** setting overrides **-qtune** option; compiler issues a warning message, and sets **-qtune** to the **-qarch** setting's default **-qtune** value.
- **-qarch** option is incompatible with user-selected **-qtune** option.

Resolution: Compiler issues a warning message, and sets **-qtune** to the **-qarch** setting's default **-qtune** value.

- Selected **-qarch** or **-qtune** options are not known to the compiler.

Resolution: Compiler issues a warning message, sets **-qarch** and **-qtune** to their default settings. The compiler mode (32- or 64-bit) is determined by the **OBJECT_MODE** environment variable or **-q32/-q64** compiler settings.

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options on the Command Line” on page 29

“Set Environment Variables to Select 64- or 32-bit Modes” on page 24

Related References

“Compiler Command Line Options” on page 51

“Resolving Conflicting Compiler Options”

“Acceptable Compiler Mode and Processor Architecture Combinations” on page 350

Resolving Conflicting Compiler Options

In general, if more than one variation of the same option is specified (with the exception of **xref** and **attr**), the compiler uses the setting of the last one specified. Compiler options specified on the command line must appear in the order you want the compiler to process them.

Two exceptions to the rules of conflicting options are the **-Idirectory** and **-Ldirectory** options, which have cumulative effects when they are specified more than once.

In most cases, the compiler uses the following order in resolving conflicting or incompatible options:

1. Pragma statements in source code will override compiler options specified on the command line.
2. Compiler options specified on the command line will override compiler options specified in a configuration file. If conflicting or incompatible compiler options are specified on the command line, the option appearing later on the command line takes precedence.
3. Compiler options specified in a configuration file will override compiler default settings.

Not all option conflicts are resolved using the above rules. The table below summarizes exceptions and how the compiler handles conflicts between them.

Option	Conflicting Options	Resolution
-qhalt	Multiple severities specified by -qhalt	Lowest severity specified
-qnoprint	-qxref -qattr -qsource -qlistopt -qlist	-qnoprint
-qfloat=rsqrt	-qnoignerrno	Last option specified
-qxref	-qxref=FULL	-qxref=FULL
-qattr	-qattr=FULL	-qattr=FULL

Option	Conflicting Options	Resolution
-p -pg -qprofile	-p -pg -qprofile	Last option specified
-qhsflt	-qrndsngl -qspnans	-qhsflt
-qhssngl	-qrndsngl -qspnans	-qhssngl
-E	-P -o -S	-E
-P	-c -o -S	-P
-#	-v	-#
-F	-B -t -W -qpath configuration file settings	-B -t -W -qpath
-qpath	-B -t	-qpath overrides -B and -t
-S	-c	-S

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Invoke the Compiler” on page 27

“Specify Compiler Options on the Command Line” on page 29

Related References

“Compiler Command Line Options” on page 51

Specify Path Names for Include Files

When you imbed one source file in another using the **#include** preprocessor directive, you must supply the name of the file to be included. You can specify a file name either by using a full path name or by using a relative path name.

- **Use a Full Path Name to Imbed Files**

The *full path name*, also called the *absolute path name*, is the file's complete name starting from the root directory. These path names start with the / (slash) character. The full path name locates the specified file regardless of the directory you are presently in (called your *working* or *current* directory).

The following example specifies the full path to file *mine.h* in John Doe's subdirectory *example_prog*:

```
/u/johndoe/example_prog/mine.h
```

- **Use a Relative Path Name to Imbed Files**

The *relative path name* locates a file relative to the directory that holds the current source file or relative to directories defined using the **-Idirectory** option.

Directory Search Sequence for Include Files Using Relative Path Names

C defines two versions of the **#include** preprocessor directive. IBM XL C Enterprise Edition supports both. With the **#include** directive, the include filename is enclosed between either the < > or " " delimiter characters.

Your choice of delimiter characters will determine the search path used to locate a given include filename. The compiler will search for that include file in all directories in the search path until the include file is found, as follows:

#include type	Directory Search Order
#include <file_name>	<ol style="list-style-type: none">1. The compiler first searches for <i>file_name</i> in each user directory specified by the -Idirectory compiler option, in the order that they appear on the command line.2. Finally, the compiler searches the standard search path for C headers. The default standard search path for C headers is /usr/include, but this path can be changed with the -qc_stdinc compiler option.
#include "file_name"	<ol style="list-style-type: none">1. The compiler first searches for the include file in the directory where your current source file resides. The current source file is the file that contains the directive #include "file_name".2. The compiler then searches for the include file according to the search order described above for #include <file_name>.

Notes:

1. *file_name* specifies the name of the file to be included, and can include a full or partial directory path to that file if you desire.
 - If you specify a file name by itself, the compiler searches for the file in the directory search list.

- If you specify a file name together with a partial directory path, the compiler appends the partial path to each directory in the search path, and tries to find the file in the completed directory path.
 - If you specify a full path name, the two versions of the **#include** directive have the same effect because the location of the file to be included is completely specified.
2. The only difference between the two versions of the **#include** directive is that the " " (user include) version first begins a search from the directory where your current source file resides. Typically, standard header files are included using the < > (system include) version, and header files that you create are included using the " " (user include) version.
 3. You can change the search order by specifying the **-qstdinc** and **-qidirfirst** options along with the **-Idirectory** option.

Use the **-qnostdinc** option to search only the directories specified with the **-Idirectory** option and the current source file directory, if applicable. For C programs, the **/usr/include** directory is not searched.

Use the **-qidirfirst** option with the **#include "file_name"** directive to search the directories specified with the **-Idirectory** option before searching other directories.

Use the **-I** option to specify the directory search paths.

Related References

"I" on page 138

"c_stdinc" on page 85

"idirfirst" on page 139

"stdinc" on page 240

Control Parallel Processing with Pragmas

Parallel processing operations are controlled by pragma directives in your program source. The pragmas have effect only when parallelization is enabled with the `-qsmp` compiler option.

You can use IBM SMP or OpenMP directives in C programs. Each have their own usage characteristics.

IBM SMP Directives

Syntax

```
#pragma ibm pragma_name_and_args  
<countable for|while|do loop>
```

Pragma directives must appear immediately before the section of code to which they apply. For most parallel processing pragma directives this section of code must be a countable loop, and the compiler will report an error if one is not found.

More than one parallel processing pragma directive can be applied to a countable loop. For example:

```
#pragma ibm independent_loop  
#pragma ibm independent_calls  
#pragma ibm schedule(static,5)  
<countable for|while|do loop>
```

Some pragma directives are mutually-exclusive of each other. If mutually-exclusive pragmas are specified for the same loop, the pragma last specified applies to the loop. In the example below, the **parallel_loop** pragma directive is applied to the loop, and the **sequential_loop** pragma directive is ignored.

```
#pragma ibm sequential_loop  
#pragma ibm parallel_loop
```

Other pragmas, if specified repeatedly for a given loop, have an additive effect. For example:

```
#pragma ibm permutation (a,b)  
#pragma ibm permutation (c)
```

is equivalent to:

```
#pragma ibm permutation (a,b,c)
```

OpenMP Directives

Syntax

```
#pragma omp pragma_name_and_args  
statement_block
```

Pragma directives generally appear immediately before the section of code to which they apply. The **omp parallel** directive is used to define the region of program code to be parallelized. Other OpenMP directives define visibility of data variables in the defined parallel region and how work within that region is shared and synchronized.

For example, the following example defines a parallel region in which iterations of a **for** loop can run in parallel:

```
#pragma omp parallel
{
  #pragma omp for
  for (i=0; i<n; i++)
    ...
}
```

This example defines a parallel region in which two or more non-iterative sections of program code can run in parallel:

```
#pragma omp sections
{
  #pragma omp section
  structured_block_1
  ...
  #pragma omp section
  structured_block_2
  ...
  ....
}
```

Related Concepts

“Program Parallelization” on page 11

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

Related References

“smp” on page 232

“Pragmas to Control Parallel Processing” on page 321

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

For complete information about the OpenMP Specification, see:

- OpenMP Web site
- OpenMP Specification

Use C with Other Programming Languages

You can use objects created in other programming languages in your C programs. The following topics in this section give an overview of programming considerations to follow when doing so.

- “Interlanguage Calling Conventions”
- “Corresponding Data Types”
- “Use the Subroutine Linkage Conventions in Interlanguage Calls” on page 43
- “Sample Program: C/C++ Calling Fortran” on page 48

Interlanguage Calling Conventions

You should follow these recommendations when writing C code to call functions written in other languages:

- Avoid using uppercase letters in identifiers. Fortran and Pascal use only lowercase letters for all external names. Although both fold external identifiers to lowercase by default, the Fortran compiler can be set to distinguish external names by case.
- Avoid using the underscore (`_`) and dollar sign (`$`) as the first character in identifiers, to prevent conflict with the naming conventions for the C and C++ language libraries.
- Avoid using long identifier names. The maximum number of significant characters in identifiers is 250 characters.

Related Concepts

“Using XL C Enterprise Edition with Other Programming Languages” on page 19

Related Tasks

“Corresponding Data Types”

“Use the Subroutine Linkage Conventions in Interlanguage Calls” on page 43

“Sample Program: C/C++ Calling Fortran” on page 48

Corresponding Data Types

The following table shows the correspondence between the data types available in C/C++, Fortran, and Pascal. Several data types in C have no equivalent representation in Pascal or Fortran. Do not use them when programming for interlanguage calls. Blank table cells indicate that no matching data type exists.

Correspondence of Data Types among C, C++, Fortran, and Pascal

C and C++ Data Types	Fortran Data Types	Pascal Data Types
bool (C++) _Bool (C)	LOGICAL*4	—
char	CHARACTER	CHAR
signed char	INTEGER*1	PACKED -128..127
unsigned char	LOGICAL*1	PACKED 0..255
signed short int	INTEGER*2	PACKED -32768..32767
unsigned short int	LOGICAL*2	PACKED 0..65535

C and C++ Data Types	Fortran Data Types	Pascal Data Types
signed long int	INTEGER*4	INTEGER
unsigned long int	LOGICAL*4	—
signed long long int	INTEGER*8	—
unsigned long long int	LOGICAL*8	—
float	REAL REAL*4	SHORTREAL
double	REAL*8 DOUBLE PRECISION	REAL
long double (default)	REAL*8 DOUBLE PRECISION	REAL
long double (with -qlongdouble or -qldb128)	REAL*16	—
structure of two floats	COMPLEX COMPLEX*4	RECORD of two SHORTREALS
structure of two doubles	COMPLEX*16 DOUBLE COMPLEX	RECORD of two REALS
structure of two long doubles (default)	COMPLEX*16	—
structure	—	RECORD (see notes below)
enumeration	INTEGER*4	Enumeration
char[n]	CHARACTER*n	PACKED ARRAY[1..n] OF CHAR
array pointer (*) to type, or type []	Dimensioned variable (transposed)	ARRAY
pointer (*) to function	Functional Parameter	Functional Parameter
structure (with -qalign=packed)	Sequence derived type	PACKED RECORD

Special Treatment of Character and Aggregate Data

Most numeric data types have counterparts across the three languages. Character and aggregate data types require special treatment:

- Because of padding and alignment differences, C structures do not exactly correspond to the Pascal **RECORD** data type.
- C character strings are delimited by a '\0' character. In Fortran, all character variables and expressions have a length that is determined at compile time. If Fortran passes a string argument to another routine, it adds a hidden argument giving the length to the end of the argument list. This length argument must be explicitly declared in C. The C code should not assume a null terminator; the supplied or declared length should always be used. Use the **strncat**, **strncpm**, and **strncpy** functions of the C runtime library. These functions are described in the *Technical Reference, Volumes 1 and 2: Base Operating System and Extensions*.
- Pascal's **STRING** data type corresponds to a C structure. For example.:

```
VAR s: STRING(10);
```

is equivalent to:

```
struct {
    int length;
    char str [10];
};
```

where length contains the actual length of STRING.

- The **-qmacpstr** option converts Pascal string literals into null-terminated strings, where the first byte contains the length of the string.
- C and Pascal store array elements in row-major order (array elements in the same row occupy adjacent memory locations). Fortran stores array elements in ascending storage units in column-major order (array elements in the same column occupy adjacent memory locations). The following example shows how a two-dimensional array declared by A[3][2] in C, A[1..3,1..2] in Pascal, and by A(3,2) in Fortran, is stored:

Storage of a Two-Dimensional Array			
Storage Unit	C and C++ Element Name	Pascal Element Name	Fortran Element Name
Lowest	A[0][0]	A[1,1]	A(1,1)
	A[0][1]	A[1,2]	A(2,1)
	A[1][0]	A[2,1]	A(3,1)
	A[1][1]	A[2,2]	A(1,2)
	A[2][0]	A[3,1]	A(2,2)
Highest	A[2][1]	A[3,2]	A(3,2)

- In general, for a multidimensional array, if you list the elements of the array in the order they are laid out in memory, a row-major array will be such that the rightmost index varies fastest, while a column-major array will be such that the leftmost index varies fastest.

Related Concepts

“Using XL C Enterprise Edition with Other Programming Languages” on page 19

Related Tasks

“Interlanguage Calling Conventions” on page 41

“Use the Subroutine Linkage Conventions in Interlanguage Calls”

“Sample Program: C/C++ Calling Fortran” on page 48

Related References

See also:

AIX 5L Version 5.1 Technical Reference: Base Operating System and Extensions, Volume 1 (A-P)

AIX 5L Version 5.1 Technical Reference: Base Operating System and Extensions, Volume 2 (Q-Z)

Use the Subroutine Linkage Conventions in Interlanguage Calls

Subroutine linkage conventions describe the machine state at subroutine entry and exit. Routines written in different languages and compiled into separate object files can be linked and executed as a single process if they follow the subroutine linkage conventions described in this section.

These linkage conventions provide fast and efficient subroutine linkage between languages. They specify how parameters are passed taking full advantage of floating-point registers (FPRs) and general-purpose registers (GPRs), and minimize the saving and restoring of registers on subroutine entry and exit.

- “Interlanguage Calls - Parameter Passing”
- “Interlanguage Calls - Call by Reference Parameters” on page 45
- “Interlanguage Calls - Call by Value Parameters” on page 45
- “Interlanguage Calls - Rules for Passing Parameters by Value” on page 45
- “Interlanguage Calls - Pointers to Functions” on page 46
- “Interlanguage Calls - Function Return Values” on page 47
- “Interlanguage Calls - Stack Floor” on page 47
- “Interlanguage Calls - Stack Overflow” on page 47
- “Interlanguage Calls - Traceback Table” on page 47
- “Interlanguage Calls - Type Encoding and Checking” on page 47
- “Sample Program: C/C++ Calling Fortran” on page 48

Related Concepts

“Using XL C Enterprise Edition with Other Programming Languages” on page 19

Related Tasks

“Interlanguage Calling Conventions” on page 41

“Corresponding Data Types” on page 41

Interlanguage Calls - Parameter Passing

Linkage conventions specify the methods for parameter passing and whether return values are to be in FPRs, GPRs, or both. The GPRs and FPRs available for argument passing are specified in two fixed lists: R3-R10 and FP1-FP13.

Prototyping affects how parameters are passed and whether widening occurs:

Nonprototyped functions

In nonprototyped functions in the C language, floating-point arguments are widened to **double** and integral types are widened to **int**.

Prototyped functions

No widening conversions occur except in arguments passed to an ellipsis function. Floating-point **double** arguments are only passed in FPRs. If an ellipsis is present in the prototype, floating-point **double** arguments are passed in both FPRs and GPRs.

When there are more argument words than available parameter GPRs and FPRs, the remaining words are passed in storage on the stack. The values in storage are the same as if they were in registers. Space for more than 8 words of arguments (float and nonfloat) must be reserved on the stack even if all the arguments were passed in registers.

The size of the parameter area is sufficient to contain all the arguments passed on any call statement from a procedure associated with the stack frame. Although not all the arguments for a particular call actually appear in storage, they can be regarded as forming a list in this area, each one occupying one or more words.

The methods of passing parameters are as follows:

- In C, all function arguments are passed by value, and the called function receives a copy of the value passed to it.
- In Fortran, by default, arguments are passed by reference, and the called function receives the address of the value passed to it. You can use the **%VAL** Fortran built-in function to pass by value. Refer to the *AIX XL Fortran User's Guide* for more information about using **%VAL** and interlanguage calls.

- In Pascal, the function declaration determines whether a parameter is expected to be passed by value or by reference.

Interlanguage Calls - Call by Reference Parameters

For call-by-reference (as in Fortran), the address of the parameter is passed in a register.

When passing parameters by reference, if you write C or C++ functions that:

- you want to call from a Fortran program, declare all parameters as pointers.
- call a program written in Fortran, all arguments must be pointers or scalars with the address operator.
- you want to call from a Pascal program, declare all parameters as pointers that the Pascal program treats as reference parameters.
- call a program written in Pascal, all arguments corresponding to reference parameters must be pointers.

Interlanguage Calls - Call by Value Parameters

In prototype functions with a variable number of arguments— specified with an ellipsis, as in *function(...)*— the compiler widens all floating-point arguments to double precision. Integral arguments (except for **long int**) are widened to **int**. Because of this widening, some data types cannot be passed between Pascal and C without explicit conversions, and Pascal routines cannot have value parameters of certain data types.

The following information refers to call by value, as in C. In the following list, arguments are classified as floating values or nonfloating values:

- Each nonfloating scalar argument requires 1 word and appears in that word exactly as it would appear in a GPR. It is right-justified, if language semantics specify, and is word aligned.
- Each float value occupies 1 word, float doubles occupy 2 successive words in the list, and long doubles occupy either 2 or 4 words, depending on the setting of the **-qldb1128/-qldouble** option.
- Structure values appear in successive words as they would anywhere in storage, satisfying all appropriate alignment requirements. Structures are aligned to a fullword and occupy $(\text{sizeof}(\text{struct } X) + (\text{wordsize} - 1)) / \text{wordsize}$ fullwords, with any padding at the end. A structure smaller than a word is left-justified within its word or register. Larger structures can occupy multiple registers and can be passed partly in storage and partly in registers.
- Other aggregate values are passed *val-by-ref*; that is, the compiler actually passes their addresses and arranges for a copy to be made in the invoked program.
- A function pointer is passed as a pointer to the routine's function descriptor. The first word contains the entry-point address. See *Interlanguage Calls - Pointers to Functions* for more information.

Interlanguage Calls - Rules for Passing Parameters by Value

The following is a 32-bit example of a call to a prototyped function:

```
int i, j; //32 bits each
long k; //32 bits
double d1, d2;
float f1;
short int s1;
```

```

char c;
...
void f(int, int, int, double, float, char, double, short);
f( i, j, k, d1, f1, c, d2, s1 );

```

The function call results in the following storage mapping:

Will be passed in:	Storage mapping of the PARM area on the Stack
R3	0 i
R4	4 j
R5	8 k
FP1 (R6, R7 unused)	12 d1
FP2 (R8 unused)	16 f1
R9	24 /// /// /// c ← right justified for a nonprototyped C function
FP3 (R10 unused)	28 d2
Stack	32
	36 /// /// /// s1 ← right justified for a nonprototyped C function

Notes:

1. A parameter is guaranteed to be mapped only if its address is taken.
2. Data with less than fullword alignment is copied into high-order bytes. Because the function in the example is prototyped, the mapping of parameters c and s1 is right-justified.
3. The parameter list is a conceptually contiguous piece of storage containing a list of words. For efficiency, the first 8 words of the list are not actually stored in the space reserved for them, but passed in GPR3-GPR10. Furthermore, the first 13 floating point value parameter values are not passed in GPRs, but are passed in FPR1-FPR13. In all cases, parameters beyond the first 8 words of the list are also stored in the space reserved for them.
4. If the called procedure intends to treat the parameter list as a contiguous piece of storage (for example, if the address of a parameter is taken in C), the parameter registers are stored in the space reserved for them in the stack.
5. A register image is stored on the stack.
6. The argument area (P₁ ... P_n) must be large enough to hold the largest parameter list.

Interlanguage Calls - Pointers to Functions

A function pointer is a data type whose values range over function addresses. Variables of this type appear in several programming languages such as C and Fortran. In Fortran, a dummy argument that appears in an **EXTERNAL** statement is a function pointer. Function pointers are supported in contexts such as the target of a call statement or an actual argument of such a statement.

A function pointer is a fullword quantity that is the address of a function descriptor. The function descriptor is a 3-word object. The first word contains the address of the entry point of the procedure, the second has the address of the TOC of the module in which the procedure is bound, and the third is the environment pointer for languages such as Pascal. There is only one function descriptor per

entry point. It is bound into the same module as the function it identifies, if the function is external. The descriptor has an external name, which is the same as the function name, but without a leading . (dot). This descriptor name is used in all import and export operations.

Interlanguage Calls - Function Return Values

Functions pass their return values according to type:

- Pointers, enumerated types, and integral values (**int**, **short**, **long**, **char**, and unsigned types) of any length are returned, right-justified, in R3; **long long** values are returned in R3 and R4. (R3 in 64-bit mode)
- **floats** and **doubles** are returned in FP1; 128-bit **long doubles** are returned in FP1 and FP2.
- Calling functions supply a pointer to a memory location where the called function stores the returned value.
- **long doubles** are returned in R1 and R2.

Interlanguage Calls - Stack Floor

The *stack floor* is a system-defined address below which the stack cannot grow.

Other system invariants related to the stack must be maintained by all compilers and assemblers:

- No data is saved or accessed from an address lower than the stack floor.
- The stack pointer is always valid. When the stack frame size is more than 32767 bytes, take care to ensure that its value is changed in a single instruction, so that there is no timing window in which a signal handler would either overlay the stack data or erroneously appear to overflow the stack segment.

Interlanguage Calls - Stack Overflow

Linkage conventions require no explicit inline check for overflow. The operating system uses a storage-protect mechanism to detect stores past the end of the stack segment.

Interlanguage Calls - Traceback Table

The compiler supports the traceback mechanism, which is required by the AIX Operating System symbolic debugger to unravel the call or return stack. Each function has a traceback table in the text segment at the end of its code. This table contains information about the function, including the type of function as well as stack frame and register information.

Interlanguage Calls - Type Encoding and Checking

Detecting errors before a program is run is a key objective of IBM XL C Enterprise Edition. Runtime errors are hard to find, and many are caused by mismatching subroutine interfaces or conflicting data definitions.

XL C Enterprise Edition uses a scheme for early detection that encodes information about all external symbols (data and programs). If the **-qextchk** option has been specified, this information about external symbols is checked at bind or load time for consistency.

The *AIX 5L for POWER-based Systems: Assembler Language Reference* book describes the following details of the Subroutine Linkage Convention:

- Register usage (general-purpose, floating-point, and special-purpose registers)

- Stack
- The calling routine's responsibilities
- The called routine's responsibilities

Sample Program: C/C++ Calling Fortran

A C or C++ program can call a Fortran function or subroutine.

The following example illustrates how program units written in different languages can be combined to create a single program. It also demonstrates parameter passing between C/C++ and Fortran subroutines with different data types as arguments.

```
#include <iostream>
extern double add(int *, double [],
int *, double []);

double ar1[4]={1.0, 2.0, 3.0, 4.0};
double ar2[4]={5.0, 6.0, 7.0, 8.0};

main()
{
int x, y;
double z;

x = 3;
y = 3;

z = add(&x, ar1, &y, ar2); /* Call Fortran add routine */
/* Note: Fortran indexes arrays 1..n*/
/* C indexes arrays 0..(n-1) */

printf("The sum of %1.0f and %1.0f is %2.0f \n",
ar1[x-1], ar2[y-1], z);
}
```

The Fortran subroutine is:

C Fortran function add.f - for C/C++ interlanguage call example
C Compile separately, then link to C/C++ program

```
REAL FUNCTION ADD*8 (A, B, C, D)
REAL*8 B,D
INTEGER*4 A,C
DIMENSION B(4), D(4)
ADD = B(A) + D(C)
RETURN
END
```

Related Concepts

“Using XL C Enterprise Edition with Other Programming Languages” on page 19

Related Tasks

“Interlanguage Calling Conventions” on page 41

“Corresponding Data Types” on page 41

“Use the Subroutine Linkage Conventions in Interlanguage Calls” on page 43

Part 3. Reference

Compiler Options

This section describes the compiler options available in XL C Enterprise Edition. Options fall into three general groups, as described in the following topics in this section.

- “Compiler Command Line Options”
- “General Purpose Pragmas” on page 271
- “Pragmas to Control Parallel Processing” on page 321

Compiler Command Line Options

This section lists and describes XL C Enterprise Edition command line options.

To get detailed information on any option listed, see the full description page(s) for that option. Those pages describe each of the compiler options, including:

- The command-line syntax of the compiler option. The first line under the **Syntax** heading specifies the command-line or configuration-file method of specification. The second line, if one appears, is the **#pragma options** keyword for use in your source file.
- The default setting of the option if you do not specify the option on the command line, in the configuration file, or in a **#pragma** directive within your program.
- The purpose of the option and additional information about its behavior.

Related Concepts

“Compiler Options” on page 5

Related Tasks

“Specify Compiler Options on the Command Line” on page 29

“Specify Compiler Options in Your Program Source Files” on page 31

“Specify Compiler Options in a Configuration File” on page 32

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33

“Resolving Conflicting Compiler Options” on page 35

Related References

“General Purpose Pragmas” on page 271

“Pragmas to Control Parallel Processing” on page 321

“Acceptable Compiler Mode and Processor Architecture Combinations” on page 350

Summary of Command Line Compiler Options

Option Name	Type	Default	Description
# (pound sign)	<i>-flag</i>	-	Traces the compilation without doing anything.
32, 64	<i>-qopt</i>	32	Selects 32- or 64-bit compiler mode.
aggrcopy	<i>-qopt</i>	See aggrcopy .	Enables destructive copy operations for structures and unions.
alias	<i>-qopt</i>	See alias .	Specifies which type-based aliasing is to be used during optimization.
align	<i>-qopt</i>	align=full	Specifies what aggregate alignment modes the compiler uses for file compilation.
alloca	<i>-qopt</i>	-	Substitutes inline code for calls to function alloca as if #pragma alloca directives are in the source code.
ansialias	<i>-qopt</i>	See -qansialias .	Specifies whether type-based aliasing is to be used during optimization.
arch	<i>-qopt</i>	arch=com	Specifies the architecture on which the executable program will be run.
asm	<i>-qopt</i>	noasm	Instructs the compiler to accept (or not) assembler language extensions.
asm_as	<i>-qopt</i>	—	Specifies the path and flags used to invoke the assembler.
assert	<i>-qopt</i>	noassert	Instructs the compiler to apply aliasing assertions to your compilation unit.
attr	<i>-qopt</i>	noattr	Produces a compiler listing that includes an attribute listing for all identifiers.
B	<i>-flag</i>	-	Determines substitute path names for the compiler, assembler, linkage editor, and preprocessor.
b	<i>-flag</i>	bdynamic	Instructs the linker to process subsequent shared objects as either dynamic, shared or static.
bitfields	<i>-qopt</i>	bitfields=unsigned	Specifies if bit fields are signed.
bmaxdata	<i>-flag</i>	0	Sets the size of the heap in bytes.
brtl	<i>-flag</i>	-	Enables runtime linking.
C	<i>-flag</i>	-	Preserves comments in preprocessed output.
c	<i>-flag</i>	-	Instructs the compiler to pass source files to the compiler only.
c_stdinc	<i>-qopt</i>	-	Changes the standard search location for the C headers.
cache	<i>-qopt</i>	-	Specify a cache configuration for a specific execution machine.
chars	<i>-qopt</i>	chars=unsigned	Instructs the compiler to treat all variables of type char as either signed or unsigned .
check	<i>-qopt</i>	nocheck	Generates code which performs certain types of run-time checking.
compact	<i>-qopt</i>	nocompact	When used with optimization, reduces code size where possible, at the expense of execution speed.

Option Name	Type	Default	Description
cplusplus	-qopt	See cplusplus .	Use this option if you want C++ comments to be recognized in C source files.
D	-flag	-	Defines the identifier <i>name</i> as in a #define preprocessor directive.
dataimported	-qopt	-	Mark data as imported.
datalocal	-qopt	-	Marks data as local.
mbsc, dbcs	-qopt	nodbcs	Use the -qdbcs option if your program contains multibyte characters.
dbxextra	-qopt	nodbxextra	Specifies that all typedef declarations, struct , union , and enum type definitions are included for debugger processing.
digraph	-qopt	See digraph .	Enables the use of digraph character sequences in your program source.
directstorage	-qopt	nodirectstorage .	Informs the compiler that write-through enabled or cache-inhibited storage may be referenced.
dollar	-qopt	nodollar	Allows the \$ symbol to be used in the names of identifiers.
dpcl	-qopt	nodpcl	Generates block scopes to support the IBM Dynamic Probe Class Library.
E	-flag	-	Instructs the compiler to preprocess the source files.
e	-flag	-	Specifies the entry name for the shared object. Equivalent to using ld -e name . See your system documentation for additional information about ld options.
enum	-qopt	See enum .	Specifies the amount of storage occupied by the enumerations.
expfile	-qopt	-	Saves all exported symbols in a file.
extchk	-qopt	noextchk	Generates bind-time type checking information and checks for compile-time consistency.
F	-flag	-	Names an alternative configuration file for the compiler.
f	-flag	-	Names a file to store a list of object files.
fdpr	-qopt	nofdpr	Collects program information for use with the AIX fdpr performance-tuning utility.
flag	-qopt	flag=ii	Specifies the minimum severity level of diagnostic messages to be reported.
float	-qopt	See float .	Specifies various floating point options to speed up or improve the accuracy of floating point operations.
fltrap	-qopt	nofltrap	Generates extra instructions to detect and trap floating point exceptions.
fold	-qopt	fold	Specifies that constant floating point expressions are to be evaluated at compile time.
format	-qopt	noformat	Warns of possible problems with string input and output format specifications.

Option Name	Type	Default	Description
fullpath	-qopt	nofullpath	Specifies what path information is stored for files when you use the -g option.
funcsect	-qopt	nofuncsect	Place instructions for each function in a separate object file control section or csect.
G	-flag	-	Linkage editor (ld command) option only. Used to generate a dynamic library file.
g	-flag	-	Generates debugging information for use by a debugger such as the Distributed Debugger.
genproto	-qopt	nogenproto	Produces ANSI prototypes from K&R function definitions.
halt	-qopt	halt=s	Instructs the compiler to stop after the compilation phase when it encounters errors of specified <i>severity</i> or greater.
heapdebug	-qopt	noheapdebug	Enables debug versions of memory management functions.
hot	-qopt	nohot	Instructs the compiler to perform high-order loop analysis and transformations during optimization.
hsflt	-qopt	nohsflt	Speeds up calculations by removing range checking on single-precision float results and on conversions from floating point to integer.
hssngl	-qopt	nohssngl	Specifies that single-precision expressions are rounded only when the results are stored into float memory locations.
I	-flag	-	Specifies an additional search path for #include filenames that do not specify an absolute path.
idirfirst	-qopt	noidirfirst	Specifies the search order for files included with the #include "file_name" directive.
ignerrno	-qopt	noignerrno	Allows the compiler to perform optimizations that assume errno is not modified by system calls.
ignprag	-qopt	-	Instructs the compiler to ignore certain pragma statements.
info	-qopt	noinfo	Produces informational messages.
initauto	-qopt	noinitauto	Initializes automatic storage to a specified two-digit hexadecimal byte value.
inlglue	-qopt	noinlglue	Generates fast external linkage by inlining the pointer glue code necessary to make a call to an external function or a call through a function pointer.
inline	-qopt	See inline .	Attempts to inline functions instead of generating calls to a function.
ipa	-qopt	See ipa .	Turns on or customizes a class of optimizations known as interprocedural analysis (IPA).
isolated_call	-qopt	-	Specifies functions in the source file that have no side effects.
keepparm	-qopt	nokeepparm	C++ Ensures that function parameters are stored on the stack even if the application is optimized.

Option Name	Type	Default	Description
keyword	<i>-qopt</i>	See keyword .	Controls whether a specified string is treated as a keyword or an identifier.
L	<i>-flag</i>	See L .	Searches the specified directory for library files specified by the -l option.
l	<i>-flag</i>	See l .	Searches a specified library for linking.
langlvl	<i>-qopt</i>	See langlvl .	Selects the C language level for compilation.
largepage	<i>-qopt</i>	nolargepage .	Instructs the compiler to exploit large page heaps available on POWER4 and POWER5 systems running AIX v5.1D or later.
ldbl128, longdouble	<i>-qopt</i>	noldbl128	Increases the size of long double type from 64 bits to 128 bits.
libansi	<i>-qopt</i>	nolibansi	Assumes that all functions with the name of an ANSI C library function are in fact the system functions.
linedebug	<i>-qopt</i>	nolinedebug	Generates abbreviated line number and source file name information for the debugger.
list	<i>-qopt</i>	nolist	Produces a compiler listing that includes an object listing.
listopt	<i>-qopt</i>	nolistopt	Produces a compiler listing that displays all options in effect.
longlit	<i>-qopt</i>	nolonglit	Makes unsuffixed literals the long type for 64-bit mode.
longlong	<i>-qopt</i>	See longlong .	Allows long long types in your program.
M	<i>-flag</i>	-	Creates an output file that contains targets suitable for inclusion in a description file for the make command.
ma	<i>-flag</i>	-	Substitutes inline code for calls to function alloca as if #pragma alloca directives are in the source code.
macpstr	<i>-qopt</i>	nomacpstr	Converts Pascal string literals into null-terminated strings where the first byte contains the length of the string.
maf	<i>-qopt</i>	maf	Specifies whether the floating-point multiply-add instructions are to be generated.
makedep	<i>-qopt</i>	-	Creates an output file that contains targets suitable for inclusion in a description file for the make command.
maxerr	<i>-qopt</i>	nomaxerr	Instructs the compiler to halt compilation when a specified number of errors of specified or greater severity is reached.
maxmem	<i>-qopt</i>	maxmem=8192	Limits the amount of memory used for local tables of specific, memory-intensive optimizations.
mbs, ducs	<i>-qopt</i>	nombcs	Use the -qmbcs option if your program contains multibyte characters.
mkshrobj	<i>-qopt</i>	-	Creates a shared object from generated object files.
O, optimize	<i>-qopt, -flag</i>	nooptimize	Optimizes code at a choice of levels during compilation.

Option Name	Type	Default	Description
o	<i>-flag</i>	-	Specifies an output location for the object, assembler, or executable files created by the compiler.
P	<i>-flag</i>	-	Preprocesses the C source files named in the compiler invocation and creates an output preprocessed source file for each input source file.
p	<i>-flag</i>	-	Sets up the object files produced by the compiler for profiling.
pascal	<i>-qopt</i>	nopascal	Ignores the word pascal in type specifiers and function declarations.
path	<i>-qopt</i>	-	Constructs alternate program and path names.
pdf1, pdf2	<i>-qopt</i>	nopdf1, nopdf2	Tunes optimizations through Profile-Directed Feedback.
pg	<i>-flag</i>	-	Sets up the object files for profiling, but provides more information than is provided by the -p option.
phsinfo	<i>-qopt</i>	nophsinfo	Reports the time taken in each compilation phase.
pic	<i>-qopt</i>	pic=small	Instructs the compiler to generate Position-Independent Code suitable for use in shared libraries.
prefetch	<i>-qopt</i>	prefetch	Enables generation of prefetching instructions in compiled code.
print	<i>-qopt</i>	print	-qnoprint suppresses listings.
proclocal, procimported, procunknown	<i>-qopt</i>	See proclocal .	Mark functions as local, imported, or unknown.
proto	<i>-qopt</i>	noproto	Assumes all functions are prototyped.
Q	<i>-flag</i>	See Q .	Attempts to inline functions.
r	<i>-flag</i>	-	Produces a relocatable object.
report	<i>-qopt</i>	noreport	Instructs the compiler to produce transformation reports that show how program loops are parallelized and optimized.
ro	<i>-qopt</i>	See ro .	Specifies the storage type for string literals.
roconst	<i>-qopt</i>	See roconst .	Specifies the storage location for constant values.
roptr	<i>-qopt</i>	noroptr	Specifies the storage location for constant pointers.
rrm	<i>-qopt</i>	norm	Prevents floating-point optimizations that are incompatible with run-time rounding to plus and minus infinity modes.
S	<i>-flag</i>	-	Generates an assembly language file (.s) for each source file.
s	<i>-flag</i>	-	Strips symbol table.
saveopt	<i>-qopt</i>	nosaveopt	Saves the command line compiler options into an object file.

Option Name	Type	Default	Description
showinc	<i>-qopt</i>	noshowinc	Used together with -qsource to selectively show user header files (includes using " ") or system header files (includes using < >) in the program source listing.
showpdf	<i>-qopt</i>	noshowpdf	Used together with -qpdf1 to add additional call and block count profiling information to an executable.
smallstack	<i>-qopt</i>	nosmallstack	Instructs the compiler to reduce the size of the stack frame.
smp	<i>-qopt</i>	nosmp	Enables parallelization of program code.
source	<i>-qopt</i>	nosource	Produces a compiler listing and includes source code.
sourcetype	<i>-qopt</i>	sourcetype=default	Instructs the compiler to treat all source files as if they are the source type specified by this option, regardless of actual source filename suffix.
spill	<i>-qopt</i>	spill=512	Specifies the size of the register allocation spill area.
spnans	<i>-qopt</i>	nospnans	Generates extra instructions to detect signalling NaN on conversion from single precision to double precision.
srcmsg	<i>-qopt</i>	nosrcmsg	Adds the corresponding source code lines to the diagnostic messages in the stderr file.
statsym	<i>-qopt</i>	nostatsym	Adds user-defined, non-external names that have a persistent storage class to the name list.
stdinc	<i>-qopt</i>	stdinc	Specifies which files are included with #include <file_name> and #include "file_name" directives.
strict	<i>-qopt</i>	See strict .	Turns off aggressive optimizations of the -O3 option that have the potential to alter the semantics of your program.
strict_induction	<i>-qopt</i>	See strict_induction .	Disables loop induction variable optimizations that have the potential to alter the semantics of your program.
suppress	<i>-qopt</i>	See suppress .	Specifies compiler message numbers to be suppressed.
symtab	<i>-qopt</i>	-	Determines what information appears in the symbol table.
syntaxonly	<i>-qopt</i>	-	Causes the compiler to perform syntax checking without generating an object file.
t	<i>-flag</i>	See t .	Adds the prefix specified by the -B option to designated programs.
tabsize	<i>-qopt</i>	tabsize=8	Changes the length of tabs as perceived by the compiler.
tbtable	<i>-qopt</i>	See tbtable .	Sets traceback table characteristics.
threaded	<i>-qopt</i>	See threaded .	Indicates that the program will run in a multi-threaded environment.
tocdata	<i>-qopt</i>	notocdata .	Specifies the thread-local storage model to be used by the application.

Option Name	Type	Default	Description
tocmerge	<i>-qopt</i>	notocmerge.	Enables TOC merging to reduce TOC pointer loads and improves the scheduling of external loads.
trigraph	<i>-qopt</i>	See trigraph.	Enables the use of trigraph character sequences in your program source.
tune	<i>-qopt</i>	See tune.	Specifies the architecture for which the executable program is optimized.
U	<i>-flag</i>	-	Undefines a specified identifier defined by the compiler or by the -D option.
unroll	<i>-qopt</i>	unroll=auto	Unrolls inner loops in the program.
unwind	<i>-qopt</i>	unwind	Informs the compiler that the application does not rely on any program stack unwinding mechanism.
upconv	<i>-qopt</i>	noupconv	Preserves the unsigned specification when performing integral promotions.
utf	<i>-qopt</i>	noutf	Enables recognition of UTF literal syntax.
V	<i>-flag</i>	-	Instructs the compiler to report information on the progress of the compilation in a command-like format.
v	<i>-flag</i>	-	Instructs the compiler to report information on the progress of the compilation.
W	<i>-flag</i>	-	Passes the listed words to a designated compiler program.
w	<i>-flag</i>	-	Requests that warning messages be suppressed.
warn64	<i>-qopt</i>	nowarn64	Enables checking for possible data conversion problems between 32-bit and 64-bit compiler modes.
weaksymbol	<i>-qopt</i>	noweaksymbol	Instructs the compiler to generate weak symbols.
xcall	<i>-qopt</i>	noxcall	Generates code to treat static routines within a compilation unit as if they were external calls.
xref	<i>-qopt</i>	noxref	Produces a compiler listing that includes a cross-reference listing of all identifiers.
y	<i>-flag</i>	-yn	Specifies the compile-time rounding mode of constant floating-point expressions.
Z	<i>-flag</i>	-	Specifies a search path for library names.

(pound sign)

Purpose

Traces the compilation without invoking anything. This option previews the compilation steps specified on the command line. When the `xlc` command is issued with this option, it names the programs within the preprocessor, compiler, and linkage editor that would be invoked, and the options that would be specified to each program. The preprocessor, compiler, and linkage editor are not invoked.

Syntax

►► — `-#` ————— ◀◀

Notes

Use this command to determine the commands and files that will be involved in a particular compilation. It avoids the overhead of compiling the source code and overwriting any existing files, such as `.lst` files. Information is displayed to standard output.

This option displays the same information as `-v`, but does not invoke the compiler. The `-#` option overrides the `-v` option.

Example

To preview the steps for the compilation of the source file `myprogram.c`, enter:

```
xlc myprogram.c -#
```

Related References

“Compiler Command Line Options” on page 51

“v” on page 262

32, 64

Purpose

Selects either 32- or 64-bit compiler mode.

Syntax



Notes

The **-q32** and **-q64** options override the compiler mode set by the value of the `OBJECT_MODE` environment variable, if it exists. If this option is not explicitly specified on the command line, and the `OBJECT_MODE` environment variable is not set, the compiler will default to 32-bit output mode.

If the compiler is invoked in 64-bit mode, the `__64BIT__` preprocessor macro is defined.

Use **-q32** and **-q64** options, along with the **-qarch** and **-qtune** compiler options, to optimize the output of the compiler to the architecture on which that output will be used. Refer to the *Acceptable Compiler Mode and Processor Architecture Combinations* table for valid combinations of the **-q32**, **-q64**, **-qarch**, and **-qtune** compiler options. In 64-bit mode, **-qarch=com** is treated the same as **-qarch=ppc**.

Using **-qarch=ppc** or any ppc family architecture with **-qfloat=hssngl** or **-qfloat=hsflt** may produce incorrect results on rs64II or future systems.

Example

To specify that the executable program testing compiled from `myprogram.c` is to run on a computer with a 32-bit PowerPC® architecture, enter:

```
xlc -o testing myprogram.c -q32 -qarch=ppc
```

Important!

1. If you mix 32- and 64-bit compilation modes for different source files, your XCOFF objects will not bind. You must recompile completely to ensure that all objects are in the same mode.
2. Your link options must reflect the type of objects you are linking. If you compiled 64-bit objects, you must link these objects using 64-bit mode.

Related References

“Compiler Command Line Options” on page 51

“arch” on page 70

“float” on page 118

“tune” on page 253

“warn64” on page 265

“Acceptable Compiler Mode and Processor Architecture Combinations” on page 350

aggrcopy

Purpose

Enables destructive copy operations for structures and unions.

Syntax

►► — -q—aggrcopy—= —nooverlap— —overlap— ►►

Default Setting

The default setting of this option is **-qaggrcopy=overlap** when compiling with **-qlanglvl=extended** or **-qlanglvl=classic** in effect. Otherwise, the default is **-qaggrcopy=nooverlap**.

Programs that do not comply to the ANSI C standard as it pertains to non-overlap of source and destination assignment may need to be compiled with the **-qaggrcopy=overlap** compiler option.

Notes

If the **-qaggrcopy=nooverlap** compiler option is enabled, the compiler assumes that the source and destination for structure and union assignments do not overlap. This assumption lets the compiler generate faster code.

Example

```
xlc myprogram.c -qaggrcopy=nooverlap
```

Related References

“Compiler Command Line Options” on page 51

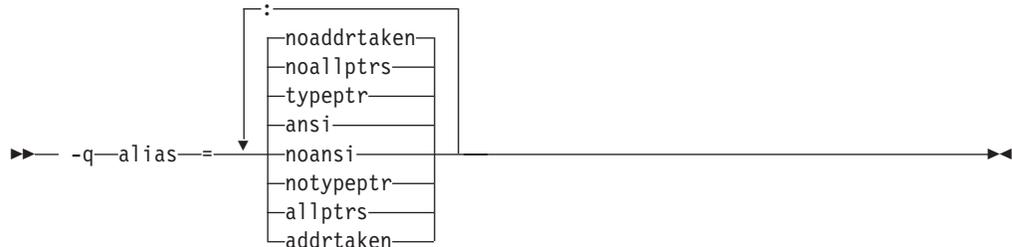
“langlvl” on page 165

alias

Purpose

Instructs the compiler to apply aliasing assertions to your compilation unit. The compiler will take advantage of the aliasing assertions to improve optimizations where possible, unless you specify otherwise.

Syntax



where available aliasing options are:

[no]typeptr	If notypeptr is specified, pointers to different types are never aliased. In other words, in the compilation unit, no two pointers of different types will point to the same storage location.
[no]allptrs	If noallptrs is specified, pointers are never aliased (this also implies -qalias=typeptr). Therefore, in the compilation unit, no two pointers will point to the same storage location.
[no]addrtaken	If noaddrtaken is specified, variables are disjoint from pointers unless their address is taken. Any class of variable for which an address has <i>not</i> been recorded in the compilation unit will be considered disjoint from indirect access through pointers.
[no]ansi	If ansi is specified, type-based aliasing is used during optimization, which restricts the lvalues that can be safely used to access a data object. The optimizer assumes that pointers can <i>only</i> point to an object of the same type. This (ansi) is the default for the xlc and c89 compilers . This option has no effect unless you also specify the -O option.

If you select **noansi**, the optimizer makes worst-case aliasing assumptions. It assumes that a pointer of a given type can point to an external object or any object whose address is already taken, regardless of type. This is the default for the **cc** compiler.

Notes

The following are not subject to type-based aliasing:

- Signed and unsigned types. For example, a pointer to a **signed int** can point to an **unsigned int**.
- Character pointer types can point to any type.
- Types qualified as **volatile** or **const**. For example, a pointer to a **const int** can point to an **int**.

Example

To specify worst-case aliasing assumptions when compiling `myprogram.c`, enter:

```
xlc myprogram.c -O -qalias=noansi
```

Related References

“Compiler Command Line Options” on page 51

“ansialias” on page 69

“#pragma disjoint” on page 279

align

Purpose

Specifies what aggregate alignment rules the compiler uses for file compilation.

Syntax



where available alignment options are:

full	The compiler uses the RISC System/6000 [®] alignment rules. This is the default. Note: The full suboption is the default to ensure backward-compatibility with existing objects. If backward-compatibility is not necessary, you should consider using natural alignment to improve potential application performance.
power	The compiler uses the RISC System/6000 alignment rules. The <i>power</i> option is the same as <i>full</i> .
mac68k	The compiler uses the Macintosh** alignment rules. This suboption is valid only for 32-bit compilations.
twobyte	The compiler uses the Macintosh alignment rules. The <i>twobyte</i> option is the same as <i>mac68k</i> . This suboption is valid only for 32-bit compilations.
packed	The compiler uses the packed alignment rules.
bit_packed	The compiler uses the bit_packed alignment rules. Alignment rules for bit_packed are the same as that for packed alignment except that bitfield data is packed on a bit-wise basis without respect to byte boundaries.
natural	The compiler maps structure members to their natural boundaries. This has the same effect as the <i>power</i> suboption, except that it also applies alignment rules to doubles and long doubles that are not the first member of a structure or union.

See also “#pragma align” on page 273 and “#pragma options” on page 299.

Notes

If you use the **-qalign** option more than once on the command line, the last alignment rule specified applies to the file.

You can control the alignment of a subset of your code by using **#pragma align=*alignment_mode*** to override the setting of the **-qalign** compiler option. Use **#pragma align=reset** to revert to a previous alignment rule. The compiler stacks alignment directives, so you can go back to using the previous alignment directive, without knowing what it is, by specifying the **#pragma align=reset** directive. For example, you can use this option if you have a structure declaration within an include file and you do not want the alignment rule specified for the structure to apply to the file in which the structure is included.

Examples

Example 1 - Affecting Only Aggregate Definition

Using the compiler invocation:

```
xlc file2.c /* <-- default alignment rule for file is */
           /*          full because no alignment rule specified */
```

Where file2.c has:

```
extern struct A A1;
typedef struct A A2;

#pragma options align=bit_packed /* <-- use bit_packed alignment rules*/
struct A {
    int a;
    char c;
};
#pragma options align=reset /* <-- Go back to default alignment rules */

struct A A1; /* <-- aligned using bit_packed alignment rules since */
A2 A3;      /*      this mode applied when struct A was defined   */
```

Example 2 - Imbedded #pragmas

Using the compiler invocation:

```
xlc -qalign=mac68k file.c /* <-- default alignment rule for file is */
                          /*      Macintosh                          */
```

Where file.c has:

```
struct A {
    int a;
    struct B {
        char c;
        double d;
    };
    #pragma options align=power /* <-- B will be unaffected by this */
                               /*      #pragma, unlike previous behavior; */
                               /*      Macintosh alignment rules still */
                               /*      in effect                          */
} BB;
#pragma options align=reset /* <-- A is unaffected by this #pragma; */
} AA;                       /*      Macintosh alignment rules still */
                               /*      in effect                          */
```

Using the `__align` specifier

You can use the `__align` specifier to explicitly specify data alignment when declaring or defining a data item.

`__align` Specifier:

Purpose: Use the `__align` specifier to explicitly specify alignment and padding when declaring or defining data items.

Syntax:

```
declarator __align (int_const) identifier;
```

```
__align (int_const) struct_or_union_specifier [identifier] {struct_decln_list}
```

where:

int_const Specifies a byte-alignment boundary. *int_const* must be an integer greater than 0 and equal to a power of 2.

Notes: The `__align` specifier can only be used with declarations of first-level variables and aggregate definitions. It ignores parameters and automatics.

The `__align` specifier can be used on an aggregate definition nested within another aggregate definition.

The `__align` specifier cannot be used in the following situations:

- Individual elements within an aggregate definition.
- Variables declared with incomplete type.
- Aggregates declared without definition.
- Individual elements of an array.
- Other types of declarations or definitions, such as **typedef**, **function**, and **enum**.
- Where the size of variable alignment is smaller than the size of type alignment.

Not all alignments may be representable in an object file.

Examples: Applying `__align` to first-level variables:

```
int __align(1024) varA;          /* varA is aligned on a 1024-byte boundary
                                and padded with 1020 bytes          */
static int __align(512) varB; /* varB is aligned on a 512-byte boundary
                                and padded with 508 bytes          */
int __align(128) functionB( ); /* An error                      */
typedef int __align(128) T;    /* An error                      */
__align enum C {a, b, c};     /* An error                      */
```

Applying `__align` to align and pad aggregate tags without affecting aggregate members:

```
__align(1024) struct structA {int i; int j;}; /* struct structA is aligned
                                                on a 1024-byte boundary
                                                with size including padding
                                                of 1024 bytes          */
__align(1024) union unionA {int i; int j;}; /* union unionA is aligned
                                                on a 1024-byte boundary
                                                with size including padding
                                                of 1024 bytes          */
```

Applying `__align` to a structure or union, where the size and alignment of the aggregate using the structure or union is affected:

```
__align(128) struct S {int i;}; /* sizeof(struct S) == 128          */
struct S sarray[10];           /* sarray is aligned on 128-byte boundary
                                with sizeof(sarray) == 1280      */
struct S __align(64) svar;     /* error - alignment of variable is
                                smaller than alignment of type    */
struct S2 {struct S s1; int a;} s2; /* s2 is aligned on 128-byte boundary
                                with sizeof(s2) == 256          */
```

Applying `__align` to an array:

```
AnyType __align(64) arrayA[10]; /* Only arrayA is aligned on a 64-byte
                                boundary, and elements within that array
                                are aligned according to the alignment
                                of AnyType. Padding is applied after the
                                back of the array and does not affect
                                the size of the array member itself. */
```

Applying `__align` where size of variable alignment differs from size of type alignment:

```
__align(64) struct S {int i;};

struct S __align(32) s1;      /* error, alignment of variable is smaller
                              than alignment of type */

struct S __align(128) s2;   /* s2 is aligned on 128-byte boundary */

struct S __align(16) s3[10]; /* error */

int __align(1) s4;          /* error */

__align(1) struct S {int i;}; /* error */
```

Related References

“Compiler Command Line Options” on page 51

“#pragma align” on page 273

“#pragma pack” on page 306

See also *Aligning data in aggregates* in the *XL C/C++ Programming Guide*.

alloca

Purpose

Substitutes inline code for calls to function `alloca`, as if `#pragma alloca` directives were in the source code.

Syntax

▶— `-q—alloca—` —▶

Notes

If `#pragma alloca` is unspecified, and if you do not use `-ma`, `alloca` is treated as a user-defined identifier rather than as a built-in function.

You may want to consider using a C99 variable length array in place of `alloca`.

Example

To compile `myprogram.c` so that calls to the function `alloca` are treated as inline, enter:

```
xlc myprogram.c -qalloca
```

Related References

“Compiler Command Line Options” on page 51

“D” on page 96

“ma” on page 182

“#pragma alloca” on page 274

ansialias

Purpose

Specifies whether type-based aliasing is to be used during optimization. Type-based aliasing restricts the lvalues that can be used to access a data object safely.

Syntax

►► -q ansialias
noansialias ◀◀

See also “#pragma options” on page 299.

Notes

This option is obsolete. Use **-qalias=** in your new applications.

The default with **xlC**, **c99** and **c89** is **ansialias**. The optimizer assumes that pointers can *only* point to an object of the same type.

The default with **cc** is **noansialias**.

This option has no effect unless you also specify the **-O** option.

If you select **noansialias**, the optimizer makes worst-case aliasing assumptions. It assumes that a pointer of a given type can point to an external object or any object whose address is already taken, regardless of type.

The following are not subject to type-based aliasing:

- Signed and unsigned types; for example, a pointer to a **signed int** can point to an **unsigned int**.
- Character pointer types can point to any type.
- Types qualified as **volatile** or **const**; for example, a pointer to a **const int** can point to an **int**.

Example

To specify worst-case aliasing assumptions when compiling `myprogram.c`, enter:

```
xlC myprogram.c -O -qnoansialias
```

Related References

“Compiler Command Line Options” on page 51

“alias” on page 62

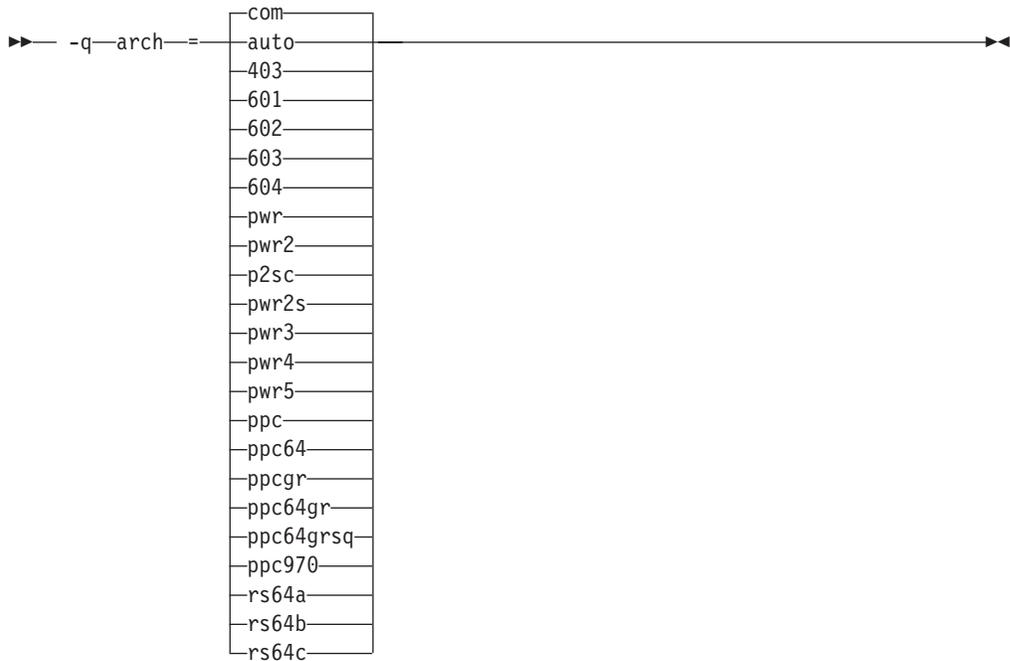
“#pragma options” on page 299

arch

Purpose

Specifies the general processor architecture for which the code (instructions) should be generated.

Syntax



where available options specify broad families of processor architectures or subgroups of those architecture families, described below.

- com**
 - This is the default if **-q32** is set or implied.
 - In 32-bit execution mode, produces object code containing instructions that will run on any of the POWER, POWER2*, and PowerPC* hardware platforms (that is, the instructions generated are *common* to all platforms. Using **-qarch=com** is referred to as compiling in *common mode*).
 - In 64-bit mode, produces object code that will run on all the 64-bit PowerPC hardware platforms but not 32-bit-only platforms.
 - Defines the `_ARCH_COM` macro.
 - This is the default option unless the **-O4** or **-O5** compiler options are specified.
- auto**
 - This is implied if **-O4** or **-O5** is set or implied.
 - Produces object code containing instructions that will run on the hardware platform on which it is compiled.
- 403**
 - Produces object code containing instructions that will run on the PowerPC 403 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC` and `_ARCH_403` macros.

- 601
 - Produces object code containing instructions that will run on the PowerPC 601 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PWR`, `_ARCH_PPC`, and `_ARCH_601` macros.
 - This option is not valid if `-q64` is set or implied.
- 602
 - Produces object code containing instructions that will run on the PowerPC 602 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC` and `_ARCH_602` macros.
 - This option is not valid if `-q64` is set or implied.
- 603
 - Produces object code containing instructions that will run on the PowerPC 603 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, and `_ARCH_603` macros.
 - This option is not valid if `-q64` is set or implied.
- 604
 - Produces object code containing instructions that will run on the PowerPC 604 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, and `_ARCH_604` macros.
 - This option is not valid if `-q64` is set or implied.
- pwr
 - Produces object code containing instructions that will run on any of the POWER, POWER2, and PowerPC 601 hardware platforms.
 - Defines the `_ARCH_COM` and `_ARCH_PWR` macros.
 - This option is not valid if `-q64` is set or implied.
- pwr2
 - Produces object code containing instructions that will run on any POWER2 platform.
 - Defines the `_ARCH_COM`, `_ARCH_PWR` and `_ARCH_PWR2` macros.
 - This option is not valid if `-q64` is set or implied.
- pwr2s
 - Produces object code containing instructions that will run on POWER2 Chip desktop implementations.
 - Defines the `_ARCH_COM`, `_ARCH_PWR`, `_ARCH_PWR2`, and `_ARCH_PWR2S` macros.
 - This option is not valid if `-q64` is set or implied.
- p2sc
 - Produces object code containing instructions that will run on the POWER2 Super Chip hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PWR`, `_ARCH_PWR2`, and `_ARCH_P2SC` macros.
 - This option is not valid if `-q64` is set or implied.
- pwr3
 - Produces object code containing instructions that will run on any POWER3™, POWER4™, POWER5™, or PowerPC 970 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, `_ARCH_PPC64GRSQ`, and `_ARCH_PWR3` macros.
- pwr4
 - Produces object code containing instructions that will run on any POWER4, POWER5, or PowerPC 970 hardware platform.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, `_ARCH_PPC64GRSQ`, `_ARCH_PWR3`, and `_ARCH_PWR4` macros.

- pwr5
 - Produces object code containing instructions that will run on the POWER5 hardware platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, `_ARCH_PPC64GRSQ`, `_ARCH_PWR3`, `_ARCH_PWR4`, and `_ARCH_PWR5` macros.
- ppc
 - In 32-bit mode, produces object code containing instructions that will run on any of the 32-bit PowerPC hardware platforms. This suboption will cause the compiler to produce single-precision instructions to be used with single-precision data.
 - Defines the `_ARCH_COM` and `_ARCH_PPC` macros.
 - Specifying `-qarch=ppc` together with `-q64` implies `-qarch=ppc64`.
- ppc64
 - Produces object code that will run on any of the 64-bit PowerPC hardware platforms.
 - This suboption can be selected when compiling in 32-bit mode, but the resulting object code may include instructions that are not recognized or behave differently when run on 32-bit PowerPC platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC` and `_ARCH_PPC64` macros.
- ppcgr
 - In 32-bit mode, produces object code for PowerPC processors that support optional graphics instructions.
 - If the `-q64` compiler option is also in effect, this option is interpreted as `ppc64gr`, described below.
 - Defines the `_ARCH_COM`, `_ARCH_PPC` and `_ARCH_PPCGR` macros.
- ppc64gr
 - Produces code for any 64-bit PowerPC hardware platform that supports optional graphics instructions.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, and `_ARCH_PPC64GR` macros.
- ppc64grsq
 - Produces code for any 64-bit PowerPC hardware platform that supports optional graphics and square root instructions.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, and `_ARCH_PPC64GRSQ` macros.
- ppc970
 - Generates instructions specific to PowerPC 970 platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC970`, `_ARCH_PWR3`, `_ARCH_PWR4`, `_ARCH_PPC64GR`, and `_ARCH_PPC64GRSQ` macros.
- rs64a
 - Produces object code that will run on RS64I platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPC64`, and `_ARCH_RS64A` macros.
- rs64b
 - Produces object code that will run on RS64II platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, `_ARCH_PPC64GRSQ`, and `_ARCH_RS64B` macros.
- rs64c
 - Produces object code that will run on RS64III platforms.
 - Defines the `_ARCH_COM`, `_ARCH_PPC`, `_ARCH_PPCGR`, `_ARCH_PPC64`, `_ARCH_PPC64GR`, `_ARCH_PPC64GRSQ`, and `_ARCH_RS64C` macros.

Notes

If you want maximum performance on a specific architecture and will not be using the program on other architectures, use the appropriate architecture option.

Using `-qarch=ppc` or any ppc family architecture with `-qfloat=hssngl` or `-qfloat=hsflt` may produce incorrect results on RS64II or future systems.

You can use `-qarch=suboption` with `-qtune=suboption`. `-qarch=suboption` specifies the architecture for which the instructions are to be generated, and `-qtune=suboption` specifies the target platform for which the code is optimized.

Example

To specify that the executable program `testing` compiled from `myprogram.c` is to run on a computer with a 32-bit PowerPC architecture, enter:

```
xlc -o testing myprogram.c -q32 -qarch=ppc
```

Related Tasks

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“O, optimize” on page 194

“tune” on page 253

“Acceptable Compiler Mode and Processor Architecture Combinations” on page 350

asm

Purpose

Instructs the compiler to accept assembler language extensions.

Syntax



Notes

This feature lets you insert inline assembler statements into your source code.

The **-qasm=gcc** compiler option instructs the compiler to recognize the **asm** keyword together with extended gcc syntax and semantics. The default setting of **-qnoasm** disables this support.

The system assembler program must be available for this command to have effect. See the **-qasm_as** compiler option for more information.

Example

The following code snippet shows a simple use of the **-qasm** compiler option:

```
int a, b, c;
int main() {
    asm("add %0, %1, %2" : "=r"(a) : "r"(b), "r"(c) );
}
```

Related References

“Compiler Command Line Options” on page 51

“asm_as” on page 75

asm_as

Purpose

Specifies the path and flags used to invoke the assembler in order to handle assembler code in an **asm** statement.

Syntax

►► — `-qasm_as=invocation_string` —►►

where **invocation_string** is the path and flags used to invoke the assembler for **asm** statements.

Notes

Use this option to specify an alternate assembler program and the flags required to invoke that assembler.

This option overrides the default setting of the **as** command defined in the compiler configuration file.

Example

To instruct the compiler to use the assembler program at `/bin/as` when it encounters inline assembler code in `myprog.c`, specify the following on the command line:

```
xlc myprog.c -qasm_as=/bin/as
```

Related References

“Compiler Command Line Options” on page 51

“asm” on page 74

assert

Purpose

Requests the compiler to apply aliasing assertions to your compilation unit. The compiler will take advantage of the aliasing assertions to improve optimizations where possible.

Syntax



where available aliasing options include:

noassert	No aliasing assertions are applied.
ASsert=TYPEptr	Pointers to different types are never aliased. In other words, in the compilation unit no two pointers of different types will point to the same storage location.
ASsert=ALLPtrs	Pointers are never aliased (this implies -qassert=typeptr). Therefore, in the compilation unit, no two pointers will point to the same storage location.
ASsert=ADDRtaken	Variables are disjoint from pointers unless their address is taken. Any class of variable for which an address has <i>not</i> been recorded in the compilation unit will be considered disjoint from indirect access through pointers.

See also “#pragma options” on page 299.

Notes

This option is obsolete. Use -qalias= in your new applications.

Related References

“Compiler Command Line Options” on page 51

“alias” on page 62

attr

Purpose

Produces a compiler listing that includes an attribute listing for all identifiers.

Syntax



where:

- qnoattr Does not produce an attribute listing for identifiers in the program.
- qattr=full Reports all identifiers in the program.
- qattr Reports only those identifiers that are used.

See also “#pragma options” on page 299.

Notes

This option does not produce a cross-reference listing unless you also specify -qxref.

The -qnoprint option overrides this option.

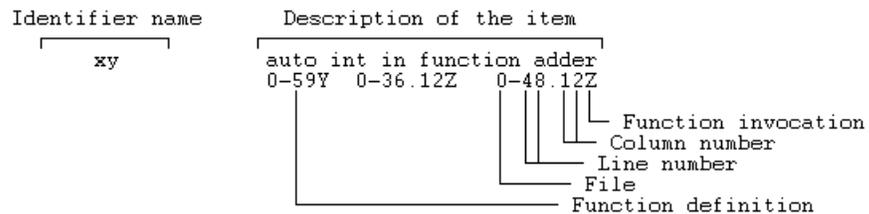
If -qattr is specified after -qattr=full, it has no effect. The full listing is produced.

Example

To compile the program myprogram.c and produce a compiler listing of all identifiers, enter:

```
xlc myprogram.c -qxref -qattr=full
```

A typical cross-reference listing has the form:



Related References

- “Compiler Command Line Options” on page 51
- “print” on page 210
- “xref” on page 268
- “#pragma options” on page 299

B

Purpose

Determines substitute path names for programs such as the compiler, assembler, linkage editor, and preprocessor.

Syntax

► -B prefix -t-program ►

where *program* can be any program name recognized by the **-t** compiler option. See the documentation for **t** for more information about specifying programs.

Notes

The optional *prefix* defines part of a path name to the new programs. The compiler does not add a / between the prefix and the program name.

To form the complete path name for each program, IBM XL C Enterprise Edition adds prefix to the standard program names for the compiler, assembler, linkage editor and preprocessor.

Use this option if you want to keep multiple levels of some or all of IBM XL C Enterprise Edition executables and have the option of specifying which one you want to use.

If **-Bprefix** is not specified, the default path is used.

-B -tprograms specifies the programs to which the **-B** prefix name is to be appended.

The **-Bprefix -tprograms** options override the **-Fconfig_file** option.

Example

To compile myprogram.c using a substitute **xlc** compiler in **/lib/tmp/mine/** enter:

```
xlc myprogram.c -B/lib/tmp/mine/ -tc
```

To compile myprogram.c using a substitute linkage editor in **/lib/tmp/mine/**, enter:

```
xlc myprogram.c -B/lib/tmp/mine/ -tl
```

Related References

“Compiler Command Line Options” on page 51

“path” on page 202

“t” on page 246

b

Purpose

Controls how shared objects are processed by the linkage editor.

Syntax



where options are:

dynamic, shared	Causes the linker to process subsequent shared objects in dynamic mode. This is the default. In dynamic mode, shared objects are not statically included in the output file. Instead, the shared objects are listed in the loader section of the output file.
static	Causes the linker to process subsequent shared objects in static mode. In static mode, shared objects are statically linked in the output file.

Notes

The default option, **-bdynamic**, ensures that the C library (**lib.c**) links dynamically. To avoid possible problems with unresolved linker errors when linking the C library, you must add the **-bdynamic** option to the end of any compilation sections that use the **-bstatic** option.

For more information about this and other **ld** options, see the *AIX Commands Reference*.

Related References

“Compiler Command Line Options” on page 51

“brtl” on page 82

See also:

ld command in *Commands Reference, Volume 5: s through u*

bmaxdata

Purpose

This option sets the maximum size of the area shared by the static data (both initialized and uninitialized) and the heap to *number* bytes. This value is used by the system loader to set the soft ulimit.

The default setting is **-bmaxdata=0**.

Syntax

►► -bmaxdata=0
number►►

Notes

Valid values for *number* are 0 and multiples of 0x10000000 (0x10000000, 0x20000000, 0x30000000, ...). The maximum value allowed by the system is 0x80000000.

If the value of *size* is 0, a single 256MB (0x10000000 byte) data segment (segment 2) will be shared by the static data, the heap, and the stack. If the value is non-zero, a data area of the specified size (starting in segment 3) will be shared by the static data and the heap, while a separate 256 MB data segment (segment 2) will be used by the stack. So, the total data size when 0 is specified is 256MB, and the total size when 0x10000000 is specified is 512MB, with 256MB for the stack and 256MB for static data and the heap.

Related References

“Compiler Command Line Options” on page 51

brtl

Purpose

Enables run-time linking for the output file.

Syntax

►► — -brtl —————►►

Notes

DCE thread libraries and heap debug libraries are not compatible with run-time linking. Do not specify the **-brtl** compiler option if you are invoking the compiler with **xlC_r4**, or if the **-qheapdebug** compiler option is specified.

Run-time linking is the ability to resolve undefined and non-deferred symbols in shared modules after the program execution has already begun. It is a mechanism for providing run-time definitions (these function definitions are not available at link-time) and symbol rebinding capabilities. The *main* application must be built to enable run-time linking. You cannot simply link any module with the run-time linker.

To include run-time linking in your program, compile using the **-brtl** compiler option. This will add a reference to the run-time linker to your program, which will be called by your program's start-up code (`/lib/crt0.o`) when program execution begins. Shared object input files are listed as dependents in the program loader section in the same order as they are specified on the command line. When the program execution begins, the system loader loads these shared objects so their definitions are available to the run-time linker.

The system loader must be able to load and resolve all symbols referenced in the the main program and called modules, or the program will not execute.

Related References

"Compiler Command Line Options" on page 51

"b" on page 79

"G" on page 129

Also, see the Shared Objects and Run-time Linking chapter in General Programming Concepts: Writing and Debugging Programs.

C

Purpose

Instructs the compiler to pass source files to the compiler only.

Syntax

▶▶ -c ◀◀

Notes

The compiled source files are not sent to the linkage editor. The compiler creates an output object file, *file_name.o*, for each valid source file, such as *file_name.c*, *file_name.i*, *file_name.C*, *file_name.cpp*, etc.

The `-c` option is overridden if either the `-E`, `-P`, or `-qsyntaxonly` options are specified.

The `-c` option can be used in combination with the `-o` option to provide an explicit name of the object file that is created by the compiler.

Example

To compile `myprogram.c` to produce an object file **myprogram.o**, but no executable file, enter the command:

```
xlc myprogram.c -c
```

To compile `myprogram.c` to produce the object file **new.o** and no executable file, enter:

```
xlc myprogram.c -c -o new.o
```

Related References

“Compiler Command Line Options” on page 51

“E” on page 105

“o” on page 198

“P” on page 199

“syntaxonly” on page 245

c_stdinc

Purpose

Changes the standard search location for the C headers.

Syntax

```
►► -qc_stdinc= path ◀◀
```

Notes

The standard search path for C headers is determined by the **-qc_stdinc** compiler option. If this compiler option is not specified or specifies an empty string, the default header file search path is used. You can find the default search path for this option in the compiler default configuration file.

If this option is specified more than once, only the last instance of the option is used by the compiler. To specify multiple directories for a search path, specify this option once, using a **:** (colon) to separate multiple search directories.

This option is ignored if the **-qnostdinc** option is in effect.

Example

To specify **mypath/headers1** and **mypath/headers2** as being part of the standard search path, enter:

```
xlc myprogram.c -qc_stdinc=mypath/headers1:mypath/headers2
```

Related Tasks

“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

“Specify Compiler Options in a Configuration File” on page 32

Related References

“Compiler Command Line Options” on page 51

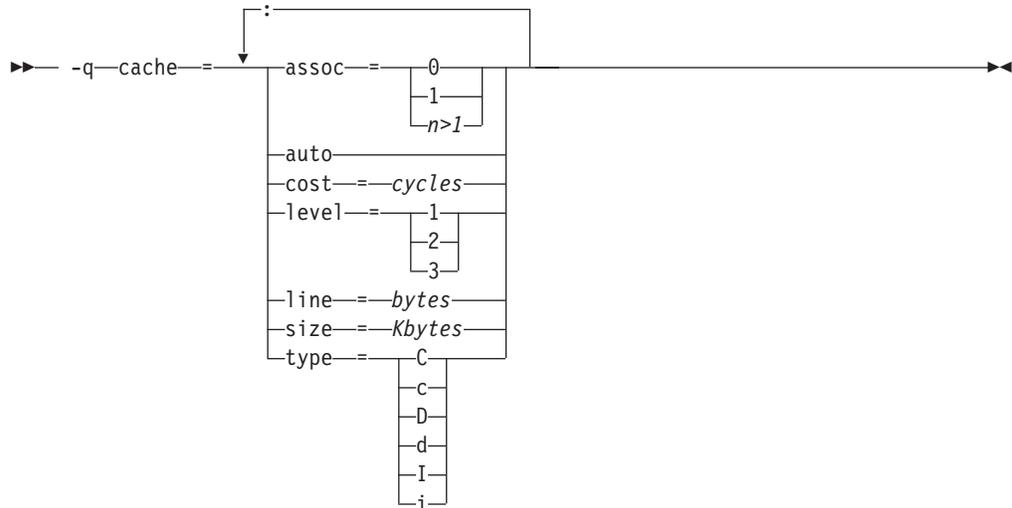
“stdinc” on page 240

cache

Purpose

The `-qcache` option specifies the cache configuration for a specific execution machine. If you know the type of execution system for a program, and that system has its instruction or data cache configured differently from the default case, use this option to specify the exact cache characteristics. The compiler uses this information to calculate the benefits of cache-related optimizations.

Syntax



where available cache options are:

<code>assoc=number</code>	Specifies the set associativity of the cache, where <i>number</i> is one of: 0 Direct-mapped cache 1 Fully associative cache N>1 n-way set associative cache
<code>auto</code>	Automatically detects the specific cache configuration of the compiling machine. This assumes that the execution environment will be the same as the compilation environment.
<code>cost=cycles</code>	Specifies the performance penalty resulting from a cache miss.
<code>level=level</code>	Specifies the level of cache affected, where <i>level</i> is one of: 1 Basic cache 2 Level-2 cache or, if there is no level-2 cache, the table lookaside buffer (TLB) 3 TLB If a machine has more than one level of cache, use a separate <code>-qcache</code> option.
<code>line=bytes</code>	Specifies the line size of the cache.
<code>size=Kbytes</code>	Specifies the total size of the cache.

`type=cache_type` The settings apply to the specified type of cache, where *cache_type* is one of:

- C or c** Combined data and instruction cache
- D or d** Data cache
- I or i** Instruction cache

Notes

The **-qtune** setting determines the optimal default **-qcache** settings for most typical compilations. You can use the **-qcache** to override these default settings. However, if you specify the wrong values for the cache configuration, or run the program on a machine with a different configuration, the program will work correctly but may be slightly slower.

You must specify **-O4**, **-O5**, or **-qipa** with the **-qcache** option.

Use the following guidelines when specifying **-qcache** suboptions:

- Specify information for as many configuration parameters as possible.
- If the target execution system has more than one level of cache, use a separate **-qcache** option to describe each cache level.
- If you are unsure of the exact size of the cache(s) on the target execution machine, specify an estimated cache size on the small side. It is better to leave some cache memory unused than it is to experience cache misses or page faults from specifying a cache size larger than actually present.
- The data cache has a greater effect on program performance than the instruction cache. If you have limited time available to experiment with different cache configurations, determine the optimal configuration specifications for the data cache first.
- If you specify the wrong values for the cache configuration, or run the program on a machine with a different configuration, program performance may degrade but program output will still be as expected.
- The **-O4** and **-O5** optimization options automatically select the cache characteristics of the compiling machine. If you specify the **-qcache** option together with the **-O4** or **-O5** options, the option specified last takes precedence.

Example

To tune performance for a system with a combined instruction and data level-1 cache, where cache is 2-way associative, 8 KB in size and has 64-byte cache lines, enter:

```
xlc -O4 -qcache=type=c:level=1:size=8:line=64:assoc=2 file.c
```

Related References

“Compiler Command Line Options” on page 51

“ipa” on page 151

“O, optimize” on page 194

chars

Purpose

Instructs the compiler to treat all variables of type **char** as either **signed** or **unsigned**.

Syntax

►► -qchars=signed | unsigned ►►

See also “#pragma chars” on page 277 and “#pragma options” on page 299.

Notes

You can also specify sign type in your source program using either of the following preprocessor directives:

```
#pragma options chars=sign_type
```

```
#pragma chars (sign_type)
```

where *sign_type* is either **signed** or **unsigned**.

Regardless of the setting of this option, the type of **char** is still considered to be distinct from the types **unsigned char** and **signed char** for purposes of type-compatibility checking.

Example

To treat all **char** types as **signed** when compiling myprogram.c, enter:

```
xlc myprogram.c -qchars=signed
```

Related References

“Compiler Command Line Options” on page 51

“#pragma chars” on page 277

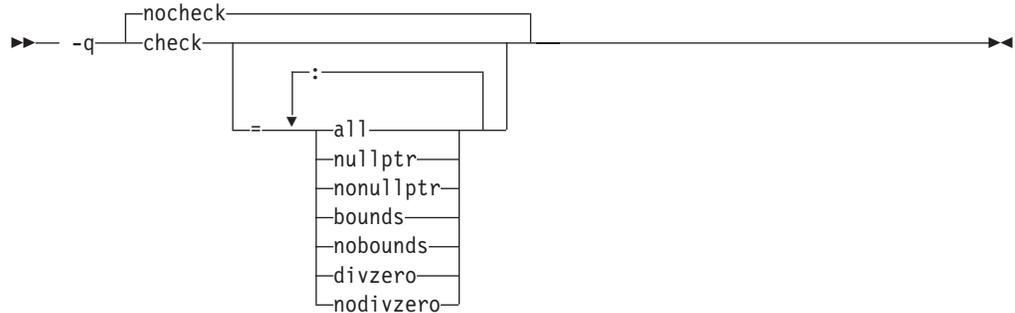
“#pragma options” on page 299

check

Purpose

Generates code that performs certain types of run-time checking. If a violation is encountered, a run-time exception is raised by sending a **SIGTRAP** signal to the process.

Syntax



where:

all

Switches on all the following suboptions. You can use the **all** option along with the **no...** form of one or more of the other options as a filter.

For example, using:

```
xlc myprogram.c -qcheck=all:nonnullptr
```

provides checking for everything except for addresses contained in pointer variables used to reference storage.

If you use **all** with the **no...** form of the options, **all** should be the first suboption.

nullptr | nonullptr

Performs run-time checking of addresses contained in pointer variables used to reference storage. The address is checked at the point of use; a trap will occur if the value is less than 512.

bounds | nobounds

Performs run-time checking of addresses when subscripting within an object of known size. The index is checked to ensure that it will result in an address that lies within the bounds of the object's storage. A trap will occur if the address does not lie within the bounds of the object.

This suboption has no effect on accesses to a variable length array.

divzero | nodivzero

Performs run-time checking of integer division. A trap will occur if an attempt is made to divide by zero.

See also “#pragma options” on page 299.

Notes

The **-qcheck** option has several suboptions, as described above. If you use more than one *suboption*, separate each one with a colon (:).

Specifying the **-qcheck** option without any suboptions, and without any other variations of **-qcheck** on the command line, turns all of the suboptions on.

Using the **-qcheck** option with suboptions turns the specified suboptions on if they do not have the no prefix, and off if they have the no prefix.

You can specify the **-qcheck** option more than once. The suboption settings are accumulated, but the later suboptions override the earlier ones.

The **-qcheck** option affects the run-time performance of the application. When checking is enabled, run-time checks are inserted into the application, which may result in slower execution.

Examples

1. For **-qcheck=nullptr:bounds**:

```
void func1(int* p) {
    *p = 42;          /* Traps if p is a null pointer */
}

void func2(int i) {
    int array[10];
    array[i] = 42;   /* Traps if i is outside range 0 - 9 */
}
```

2. For **-qcheck=divzero**:

```
void func3(int a, int b) {
    a / b;          /* Traps if b=0 */
}
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

compact

Purpose

When used with optimization, reduces code size where possible, at the expense of execution speed.

Syntax

►► -q nocompact
compact _____ ►►

See also “#pragma options” on page 299.

Notes

Code size is reduced by inhibiting optimizations that replicate or expand code inline, such as inlining or loop unrolling. Execution time may increase.

Example

To compile myprogram.c to reduce code size, enter:

```
xlc myprogram.c -O -qcompact
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

cplusplus

Purpose

Use this option if you want C++ comments to be recognized in C source files.

Syntax

► — -q — nocplusplus
cplusplus —►

Default

The default setting varies:

- **-qcplusplus** is implicitly selected when you invoke the compiler with **xlC** or **xlC_r**.
- **-qcplusplus** is also implicitly selected when **-qclanglvl** is set to **stdc99** or **extc99**. You can override these implicit selections by specifying **-qnocplusplus** after the **-qclanglvl** option on the command line; for example: **-qclanglvl=stdc99 -qnocplusplus** or **-qclanglvl=extc99 -qnocplusplus**.
- Otherwise, the default setting is **-qnocplusplus**.

Notes

The `__C99_CPLUSCMT` compiler macro is defined when **cplusplus** is selected.

The character sequence `//` begins a C++ comment, except within a header name, a character constant, a string literal, or a comment. Comments do not nest, and macro replacement is not performed within comments. The following character sequences are ignored within a C++ comment:

- `//`
- `/*`
- `*/`

C++ comments have the form `//text`. The two slashes (`//`) in the character sequence must be adjacent with nothing between them. Everything to the right of them until the end of the logical source line, as indicated by a new-line character, is treated as a comment. The `//` delimiter can be located at any position within a line.

`//` comments are *not* part of C89. The result of the following valid C89 program will be incorrect if **-qcplusplus** is specified:

```
main() {  
    int i = 2;  
    printf("%i\n", i /* 2 */  
          + 1);  
}
```

The correct answer is 2 (2 divided by 1). When **-qcplusplus** is specified, the result is 3 (2 plus 1).

The preprocessor handles all comments in the following ways:

- If the **-C** option is *not* specified, all comments are removed and replaced by a single blank.
- If the **-C** option *is* specified, comments are output unless they appear on a preprocessor directive or in a macro argument.

- If **-E** is specified, continuation sequences are recognized in all comments and are output
- If **-P** is specified, comments are recognized and stripped from the output, forming concatenated output lines.

A comment can span multiple physical source lines if they are joined into one logical source line through use of the backslash (\) character. You can represent the backslash character by a trigraph (??/).

Examples

1. Example of C++ Comments

The following examples show the use of C++ comments:

```
// A comment that spans two \
physical source lines

// A comment that spans two ??/
physical source lines
```

2. Preprocessor Output Example 1

For the following source code fragment:

```
int a;
int b; // A comment that spans two \
      physical source lines
int c;
      // This is a C++ comment
int d;
```

The output for the **-P** option is:

```
int a;
int b;
int c;

int d;
```

The C89 mode output for the **-P -C** options is:

```
int a;
int b; // A comment that spans two    physical source lines
int c;
      // This is a C++ comment
int d;
```

The output for the **-E** option is:

```
int a;
int b;

int c;

int d;
```

The C89 mode output for the **-E -C** options is:

```
#line 1 "fred.c"
int a;
int b; // a comment that spans two \
      physical source lines
int c;
      // This is a C++ comment
int d;
```

Extended mode output for the **-P -C** options or **-E -C** options is:

```

int a;
int b; // A comment that spans two \
      physical source lines
int c;
      // This is a C++ comment
int d;

```

3. Preprocessor Output Example 2 - Directive Line

For the following source code fragment:

```

int a;
#define mm 1 // This is a C++ comment on which spans two \
            physical source lines
int b;
            // This is a C++ comment
int c;

```

The output for the **-P** option is:

```

int a;
int b;

int c;

```

The output for the **-P -C** options:

```

int a;
int b;
            // This is a C++ comment
int c;

```

The output for the **-E** option is:

```

#line 1 "fred.c"
int a;
#line 4
int b;

int c;

```

The output for the **-E -C** options:

```

#line 1 "fred.c"
int a;
#line 4
int b;
            // This is a C++ comment
int c;

```

4. Preprocessor Output Example 3 - Macro Function Argument

For the following source code fragment:

```

#define mm(aa) aa
int a;
int b; mm(// This is a C++ comment
        int blah);
int c;
      // This is a C++ comment
int d;

```

The output for the **-P** option:

```

int a;
int b; int blah;
int c;

int d;

```

The output for the **-P -C** options:

```

int a;
int b; int blah;
int c;
      // This is a C++ comment
int d;

```

The output for the -E option is:

```
#line 1 "fred.c"
int a;
int b;
int blah;
int c;

int d;
```

The output for the -E -C option is:

```
#line 1 "fred.c"
int a;
int b;
int blah;
int c;
    // This is a C++ comment
int d;
```

5. Compile Example

To compile myprogram.c. so that C++ comments are recognized as comments, enter:

```
xlc myprogram.c -qcplusplus
```

Related References

"Compiler Command Line Options" on page 51

"C" on page 83

"E" on page 105

"langlvl" on page 165

"P" on page 199

D

Purpose

Defines the macro *name* as in a **#define** preprocessor directive. *definition* is an optional definition or value assigned to *name*.

Syntax

```
▶ -D name [= definition] ▶
```

Notes

You can also define a macro name in your source program using the **#define** preprocessor directive, provided that the macro name has not already been defined by the **-D** compiler option.

`-Dname=` is equivalent to `#define name`.

`-Dname` is equivalent to `#define name 1`. (This is the default.)

Using the **#define** directive to define a macro name already defined by the **-D** option will result in an error condition.

To aid in program portability and standards compliance, the operating system provides several header files that refer to macro names you can set with the **-D** option. You can find most of these header files either in the `/usr/include` directory or in the `/usr/include/sys` directory. See “Header Files Overview” in the *AIX Files Reference* for more information.

The configuration file uses the **-D** option to specify the following predefined macros:

Macro name	Applies to AIX v5.2	Applies to AIX v5.1	Applies to AIX v4.3
AIX	✓	✓	✓
AIX32	✓	✓	✓
AIX41	✓	✓	✓
AIX43	✓	✓	✓
AIX50	✓	✓	
AIX51	✓	✓	
AIX52	✓		
IBMR2	✓	✓	✓
POWER	✓	✓	✓
ANSI_C_SOURCE	✓	✓	✓

To ensure that the correct macros for your source file are defined, use the **-D** option with the appropriate macro name. If your source file includes the `/usr/include/sys/stat.h` header file, you must compile with the option **-D_POSIX_SOURCE** to pick up the correct definitions for that file.

If your source file includes the `/usr/include/standards.h` header file, `_ANSI_C_SOURCE`, `_XOPEN_SOURCE`, and `_POSIX_SOURCE` are defined if you have not defined any of them.

The `-Uname` option, which is used to undefine macros defined by the `-D` option, has a higher precedence than the `-Dname` option.

Examples

1. AIX v4.2 and later provides support for files greater than 2 gigabytes in size so you can store large quantities of data in a single file. To allow Large File manipulation in your application, compile with the `-D_LARGE_FILES` and `-qlonglong` compiler options. For example:

```
xlc myprogram.c -D_LARGE_FILES -qlonglong
```

2. To specify that all instances of the name `COUNT` be replaced by 100 in `myprogram.c`, enter:

```
xlc myprogram.c -DCOUNT=100
```

This is equivalent to having `#define COUNT 100` at the beginning of the source file.

Related References

“Compiler Command Line Options” on page 51

“U” on page 255

Appendix A, “Predefined Macros,” on page 369

See also:

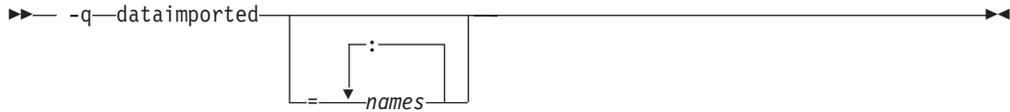
Header Files command in the Files Reference

dataimported

Purpose

Marks data as imported.

Syntax



Notes

When this option is in effect, imported variables are dynamically bound with a shared portion of a library.

- Specifying **-qdataimported** instructs the compiler to assume that all variables are imported.
- Specifying **-qdataimported=names** marks the named variables as being imported, where *names* is a list of variable names separated by colons (:). Variables not explicitly named are not affected.

Conflicts among the **-qdataimported** and **-qdatalocal** data-marking options are resolved in the following manner:

Options that list variable names:

The last explicit specification for a particular variable name is used.

Options that change the default:

This form does not specify a name list. The last option specified is the default for variables not explicitly listed in the name-list form.

Related References

“Compiler Command Line Options” on page 51

“datalocal” on page 99

datalocal

Purpose

Marks data as local.

Syntax



Notes

When this option is in effect, local variables are statically bound with the functions that use them.

- Specifying **-qdatalocal** instructs the compiler to assume that all variables are local.
- Specifying **-qdatalocal=names** marks the named variables as local, where *names* is a list of identifiers separated by colons (:). Variables not explicitly named are not affected.

Performance may decrease if an imported variable is assumed to be local.

Conflicts among the **-qdataimported** and **-qdatalocal** data-marking options are resolved in the following manner:

- | | |
|-----------------------------------|---|
| Options that list variable names: | The last explicit specification for a particular variable name is used. |
| Options that change the default: | This form does not specify a name list. The last option specified is the default for variables not explicitly listed in the name-list form. |

Related References

“Compiler Command Line Options” on page 51

“dataimported” on page 98

dbxextra

Purpose

Specifies that all **typedef** declarations, **struct**, **union**, and **enum** type definitions are included for debugging.

Syntax

►► — -q — nodbxextra
dbxextra —————►►

See also “#pragma options” on page 299.

Notes

Use this option with the **-g** option to produce additional debugging information for use with a debugger.

When you specify the **-g** option, debugging information is included in the object file. To minimize the size of object and executable files, the compiler only includes information for symbols that are referenced. Debugging information is not produced for unreferenced arrays, pointers, or file-scope variables unless **-qdbxextra** is specified.

Using **-qdbxextra** may make your object and executable files larger.

Example

To include all symbols in `myprogram.c` for debugging, enter:

```
xlc myprogram.c -g -qdbxextra
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

“#pragma options” on page 299

digraph

Purpose

Lets you use digraph key combinations or keywords to represent characters not found on some keyboards.

Syntax



See also “#pragma options” on page 299.

Defaults

- `-qnodigraph` when `-qlanglvl` is set to a value other than `extc99` or `stdc99`.
- `-qdigraph` when `-qlanglvl` is set to a value of `extc99` or `stdc99`.

Notes

A digraph is a keyword or combination of keys that lets you produce a character that is not available on all keyboards.

The digraph key combinations are:

Key Combination	Character Produced
<code><%</code>	<code>{</code>
<code>%></code>	<code>}</code>
<code><:</code>	<code>[</code>
<code>:></code>	<code>]</code>
<code>%%</code>	<code>#</code>

Example

To disable digraph character sequences when compiling your program, enter:

```
xlc myprogram.c -qnodigraph
```

Related References

“Compiler Command Line Options” on page 51

“langlvl” on page 165

“trigraph” on page 252

“#pragma options” on page 299

directstorage

Purpose

Informs the compiler that write-through-enabled or cache-inhibited storage may be referenced.

Syntax

►► -q nodirectstorage
directstorage ◀◀

Notes

The **-qdirectstorage** compiler option informs the compiler that write-through enabled or cache-inhibited storage may be referenced, and that appropriate compiler output should be generated.

The PowerPC architecture allows many different implementations of cache organization. To ensure that your application will execute correctly on all implementations, you should assume that separate instruction and data caches exist and program your application accordingly.

Depending on the storage control attributes specified by the program and the function being performed, your program may use cache instructions to guarantee that the function is performed correctly.

For example, the **dcbz** instruction allocates a block of data in the cache and then initializes it to a series of zeroes. Though it can be used to boost performance when zeroing a large block of data, the **dcbz** instruction should be used with caution because it will cause an alignment error to occur under any of the following conditions:

- The cache block specified by the instruction is in a memory region marked cache-inhibited.
- The cache is in write-through mode.
- The L1 Dcache or L2 cache is disabled.

Specifying **-qdirectstorage** will suppress generation of the **dcbz** instruction, and avoid the alignment errors mentioned above.

Related References

“Compiler Command Line Options” on page 51

dollar

Purpose

Allows the \$ symbol to be used in the names of identifiers.

Syntax

►► — -q — nodollar — dollar — ◀◀

See also “#pragma options” on page 299.

Example

To compile myprogram.c so that \$ is allowed in identifiers in the program, enter:

```
xlc myprogram.c -qdollar
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

dpcl

Purpose

Generates symbols that tools based on the Dynamic Probe Class Library (DPCL) can use to see the structure of an executable file.

Syntax

►► -q nodpcl
dpcl ►►

Notes

DPCL is an application program interface (API), or library, that can simplify the task of building programming tools. It is designed to support application performance analysis tools, but it may also be used for application steering, load balancing, and many other types of tools.

DPCL supports serial, shared memory, and message passing applications with equal ease. It uses a client/server architecture that allows a tool to create a new application or connect to a running application. DPCL inserts or removes instrumentation from the application while it is running, without recompiling the application.

When you specify the **-qdpcl** option, the compiler emits symbols to define blocks of code in a program. You can then use tools that use the DPCL interface to examine performance information such as memory usage for object files that you have compiled with this option.

You must specify **-qdpcl** together with the **-g** option to ensure that the compiler generates debugging information required by debugging and program analysis tools.

You cannot specify the **-qipa** or **-qsmp** options together with **-qdpcl**.

Related References

“Compiler Command Line Options” on page 51

“dpcl”

“g” on page 130

“ipa” on page 151

E

Purpose

Instructs the compiler to preprocess the source files.

Syntax

▶— -E —▶

Notes

The **-E** and **-P** options have different results. When the **-E** option is specified, the compiler assumes that the input is a C file and that the output will be recompiled or reprocessed in some way. These assumptions are:

- Original source coordinates are preserved. This is why **#line** directives are produced.
- All tokens are output in their original spelling, which, in this case, includes continuation sequences. This means that any subsequent compilation or reprocessing with another tool will give the same coordinates (for example, the coordinates of error messages).

The **-P** option is used for general-purpose preprocessing. No assumptions are made concerning the input or the intended use of the output. This mode is intended for use with input files that are not written in C. As such, all preprocessor-specific constructs are processed as described in the ANSI C standard. In this case, the continuation sequence is removed as described in the “Phases of Translation” of that standard. All non-preprocessor-specific text should be output as it appears.

Using **-E** causes **#line** directives to be generated to preserve the source coordinates of the tokens. Blank lines are stripped and replaced by compensating **#line** directives.

The line continuation sequence is removed and the source lines are concatenated with the **-P** option. With the **-E** option, the tokens are output on separate lines in order to preserve the source coordinates. The continuation sequence may be removed in this case.

The **-E** option overrides the **-P**, **-o**, and **-qsyntaxonly** options, and accepts any file name.

If used with the **-M** option, **-E** will work only for files with a **.c** (C source files), or a **.i** (preprocessed source files) filename suffix. Source files with unrecognized filename suffixes are treated and preprocessed as C files, and no error message is generated.

Unless **-C** is specified, comments are replaced in the preprocessed output by a single space character. New lines and **#line** directives are issued for comments that span multiple source lines, and when **-C** is not specified. Comments within a macro function argument are deleted.

The default is to preprocess, compile, and link-edit source files to produce an executable file.

Example

To compile `myprogram.c` and send the preprocessed source to standard output, enter:

```
xlc myprogram.c -E
```

If myprogram.c has a code fragment such as:

```
#define SUM(x,y) (x + y) ;
int a ;
#define mm 1 ; /* This is a comment in a
                preprocessor directive */
int b ;        /* This is another comment across
                two lines */
int c ;
                /* Another comment */
c = SUM(a, /* Comment in a macro function argument*/
        b) ;
```

the output will be:

```
#line 2 "myprogram.c"
int a;
#line 5
int b;

int c;

c =
(a + b);
```

Related References

“Compiler Command Line Options” on page 51

“M” on page 180

“o” on page 198

“P” on page 199

“syntaxonly” on page 245

e

Purpose

This option is used only together with the **-qmkshrobj** compiler option. See the description for the **-qmkshrobj** compiler option for more information.

Syntax

▶▶ — *-e—name—* —————▶▶

Related References

“Compiler Command Line Options” on page 51

“mkshrobj” on page 192

enum

Purpose

Specifies the amount of storage occupied by enumerations.

Syntax



where valid **enum** settings are:

- 1 Specifies that enumerations occupy 1 byte of storage. Enumerations are of type **char** if the range of enumeration values falls within the limits of **signed char**, and **unsigned char** otherwise.
- 2 Specifies that enumerations occupy 2 bytes of storage. Enumerations are of type **short** if the range of enumeration values falls within the limits of **signed short**, and **unsigned short** otherwise.
- 4 Specifies that enumerations occupy 4 bytes of storage. Enumerations are of type **int** if the range of enumeration values falls within the limits of **signed int**, and **unsigned int** otherwise.
- 8 Specifies that enumerations occupy 8 bytes of storage.
In 32-bit compilation mode, the enumeration is of type **long long** if the range of enumeration values falls within the limits of **signed long long**, and **unsigned long long** otherwise.
In 64-bit compilation mode, the enumeration is of type **long** if the range of enumeration values falls within the limits of **signed long**, and **unsigned long** otherwise.
- int Specifies that enumerations occupy 4 bytes of storage and are represented by **int**. Values cannot exceed the range of **signed int** in C compilations.
- intlong Specifies that enumerations will occupy 8 bytes of storage if the range of values in the enumeration exceeds the limit for **int**. See the description for **-qenum=8**.
If the range of values in the enumeration does not exceed the limit for **int**, the enumeration will occupy 4 bytes of storage and is represented by **int**.
- small Specifies that enumerations occupy the smallest amount of space (1, 2, 4, or 8 bytes of storage) that can accurately represent the range of values in the enumeration. Signage is *unsigned*, unless the range of values includes negative values.
If an 8-byte enum results, the actual enumeration type used is dependent on compilation mode. See the description for **-qenum=8**

See also “#pragma enum” on page 280 and “#pragma options” on page 299.

Notes

The **-qenum=small** option allocates to an **enum variable** the amount of storage that is required by the smallest predefined type that can represent that range of **enum** constants. By default, an unsigned predefined type is used. If any **enum** constant is negative, a signed predefined type is used.

The **-qenum=1|2|4|8** options allocate a specific amount of storage to an **enum variable**. If the specified storage size is smaller than that required by the range of **enum** variables, a Severe error message is issued and compilation stops.

The ISO C 1989 and ISO C1999 Standards require that enumeration values not exceed the range of **int**. When compiling with **-qlanglvl=stdc89** or **-qlanglvl=stdc99** in effect, the compiler will behave as follows if the value of an enumeration exceeds the range of **int**:

- If **-qenum=int** is in effect, a Severe error message is issued and compilation stops.
- For all other settings of **-qenum**, an Informational message is issued and compilation continues.

The tables that follow show the priority for selecting a predefined type. The table also shows the predefined type, the maximum range of **enum** constants for the corresponding predefined type, and the amount of storage that is required for that predefined type, that is, the value that the **sizeof** operator would yield when applied to the minimum-sized **enum**.

Related References

“Compiler Command Line Options” on page 51

“#pragma enum” on page 280

“#pragma options” on page 299

Enum sizes and types – All types are signed unless otherwise noted.

	enum=1			enum=2			enum=4			enum=8		
	var	const		var	const		var	const		var	const	
Range												
0..127	char	int		short	int		int	int		long	long	const
-128..127	char	int		short	int		int	int		long	long	long
0..255	unsigned char	int		short	int		int	int		long	long	long
0..32767	ERROR ¹	int		short	int		int	int		long	long	long
-32768..32767	ERROR ¹	int		short	int		int	int		long	long	long
0..65535	ERROR ¹	int		unsigned short	int		int	int		long	long	long
0..2147483647	ERROR ¹	int		ERROR ¹	int		int	int		long	long	long
-(2147483647+1) ..2147483647	ERROR ¹	int		ERROR ¹	int		int	int		long	long	long
0..4294967295	ERROR ¹	unsigned int		ERROR ¹	unsigned int		unsigned int	unsigned int		long	long	long
0..(2 ⁶³ -1)	ERROR ¹	long ^{2b}		ERROR ¹	long ^{2b}		ERROR ¹	long ^{2b}		long	long ^{2b}	long ^{2b}
-2 ⁶³ ..(2 ⁶³ -1)	ERROR ¹	long ^{2b}		ERROR ¹	long ^{2b}		ERROR ¹	long ^{2b}		long	long ^{2b}	long ^{2b}
0..2 ⁶⁴	ERROR ¹	unsigned long ^{2b}		ERROR ¹	unsigned long ^{2b}		ERROR ¹	unsigned long ^{2b}		unsigned long ^{2b}	unsigned long ^{2b}	unsigned long ^{2b}

Notes:

- These enumerations are too large for the **-qenum=1|2|4** settings. A Severe error is issued and compilation stops. To correct this condition, you should reduce the range of the enumerations, choose a larger **enum** setting, or choose a dynamic **enum** setting such as **small** or **intlong**.
- Enumeration types must not exceed the range of **int** when compiling C applications to ISO C 1989 and ISO C 1999 Standards. When compiling with **-qlanglvl=stdc89** or **-qlanglvl=stdc99** in effect, the compiler will behave as follows if the value of an enumeration exceeds the range of **int**:
 - If **-qenum=int** is in effect, a Severe error message is issued and compilation stops.
 - For all other settings of **-qenum**, an Informational message is issued and compilation continues.

Enum sizes and types – All types are signed unless otherwise noted.											
	enum=int			enum=intlong			enum=small				
	var	const		var	const		var	const		var	const
Range	int	int		int	int		int	int		int	int
0..127	int	int		int	int		int	int		int	int
-128..127	int	int		int	int		int	int		int	int
0..255	int	int		int	int		int	int		int	int
0..32767	int	int		int	int		int	int		int	int
-32768..32767	int	int		int	int		int	int		int	int
0..65535	int	int		int	int		int	int		int	int
0..2147483647	int	int		int	int		int	int		int	int
-(2147483647+1) ..2147483647	int	int		int	int		int	int		int	int
0..4294967295	unsigned int	unsigned int		unsigned int	unsigned int		unsigned int	unsigned int		unsigned int	unsigned int
0..(2 ⁶³ -1)	ERR ^{2a}	ERR ^{2a}		long long ^{2b}	long long ^{2b}		long long ^{2b}	long long ^{2b}		unsigned long long ^{2b}	unsigned long long ^{2b}
-2 ⁶³ ..(2 ⁶³ -1)	ERR ^{2a}	ERR ^{2a}		long long ^{2b}	long long ^{2b}		long long ^{2b}	long long ^{2b}		unsigned long long ^{2b}	unsigned long long ^{2b}
0..2 ⁶⁴	ERR ^{2a}	ERR ^{2a}		unsigned long long ^{2b}	unsigned long long ^{2b}		unsigned long long ^{2b}	unsigned long long ^{2b}		unsigned long long ^{2b}	unsigned long long ^{2b}

Notes:

- These enumerations are too large for the **-qenum=1|2|4** settings. A Severe error is issued and compilation stops. To correct this condition, you should reduce the range of the enumerations, choose a larger **enum** setting, or choose a dynamic **enum** setting, such as **small** or **intlong**.
- Enumeration types must not exceed the range of **int** when compiling C applications to ISO C 1989 and ISO C 1999 Standards. When compiling with **-qlanglvl=stdc89** or **-qlanglvl=stdc99** in effect, the compiler will behave as follows if the value of an enumeration exceeds the range of **int**:
 - If **-qenum=int** is in effect, a Severe error message is issued and compilation stops.
 - For all other settings of **-qenum**, an Informational message is issued and compilation continues.

expfile

Purpose

Saves all exported symbols in a designated file.

This option is used only together with the **-qmkshrobj** compiler option. See the description for the **-qmkshrobj** compiler option for more information.

Syntax

► — `-q-expfile=filename` —►

Related References

“Compiler Command Line Options” on page 51

“mkshrobj” on page 192

extchk

Purpose

Generates bind-time type checking information and checks for compile-time consistency.

Syntax

►► -q noextchk
extchk _____ ►►

See also “#pragma options” on page 299.

Notes

-qextchk checks for consistency at compile-time and detects mismatches across compilation units at link-time.

-qextchk does not perform type-checking on functions or objects that contain references to incomplete types.

Example

To compile myprogram.c so that bind-time checking information is produced, enter:

```
xlc myprogram.c -qextchk
```

Related References

“Compiler Command Line Options” on page 51

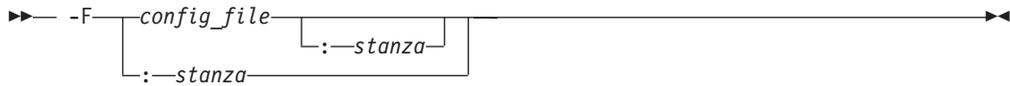
“#pragma options” on page 299

F

Purpose

Names an alternative configuration file (.cfg) for the compiler.

Syntax



where suboptions are:

<i>config_file</i>	Specifies the name of a compiler configuration file.
<i>stanza</i>	Specifies the name of the command used to invoke the compiler. This directs the compiler to use the entries under <i>stanza</i> in the <i>config_file</i> to set up the compiler environment.

Notes

The default is a configuration file supplied at installation time. Any file names or stanzas that you specify on the command line or within your source file override the defaults specified in the configuration file.

For information regarding the contents of the configuration file, refer to “Specify Compiler Options in a Configuration File” on page 32.

The `-B`, `-t`, and `-W` options override the `-F` option.

Example

To compile `myprogram.c` using a configuration file called `/usr/tmp/myvac.cfg`, enter:

```
xlc myprogram.c -F/usr/tmp/myvac.cfg:xlc
```

Related Tasks

“Specify Compiler Options in a Configuration File” on page 32

Related References

“Compiler Command Line Options” on page 51

“B” on page 78

“t” on page 246

“W” on page 263

f

Purpose

Names a file to store a list of object files for `xlc` to pass to the linker.

Syntax

▶▶ — `-f`—*filelistname*—————▶▶

Notes

The *filelistname* file should contain only the names of object files. There should be one object file per line.

This option is the same as the `-f` option for the `ld` command.

Example

To pass the list of files contained in `myobjlistfile` to the linker, enter:

```
xlc -f/usr/tmp/myobjlistfile
```

Related References

“Compiler Command Line Options” on page 51

fdpr

Purpose

Collects information about your program for use with the AIX **fdpr** (Feedback Directed Program Restructuring) performance-tuning utility.

Syntax

►► -q nofdpr
fdpr ◀◀

Notes

You should compile your program with **-qfdpr** before optimizing it with the **fdpr** performance-tuning utility. Optimization data is stored in the object file.

For more information on using the **fdpr** performance-tuning utility, refer to the *AIX Version 4 Commands Reference* or enter the command:

```
man fdpr
```

Example

To compile `myprogram.c` so it includes data required by the **fdpr** utility, enter:

```
xlc myprogram.c -qfdpr
```

Related References

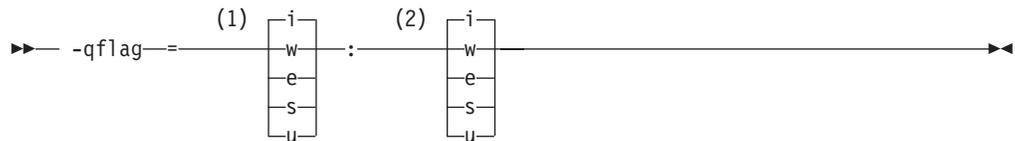
“Compiler Command Line Options” on page 51

flag

Purpose

Specifies the minimum severity level of diagnostic messages to be reported in a listing and displayed on a terminal. The diagnostic messages display with their associated sub-messages.

Syntax



Notes:

- 1 Minimum severity level messages reported in listing
- 2 Minimum severity level messages reported on terminal

where message severity levels are:

<i>severity</i>	Description
i	Information
w	Warning
e	Error
s	Severe error
u	Unrecoverable error

See also “#pragma options” on page 299.

Notes

You must specify a minimum message severity level for both listing and terminal reporting.

Specifying informational message levels does not turn on the **-qinfo** option.

Example

To compile myprogram.c so that the listing shows all messages that were generated and your workstation displays only error and higher messages (with their associated information messages to aid in fixing the errors), enter:

```
xlc myprogram.c -qflag=i:e
```

Related References

“Compiler Command Line Options” on page 51

“info” on page 142

“w” on page 264

“#pragma options” on page 299

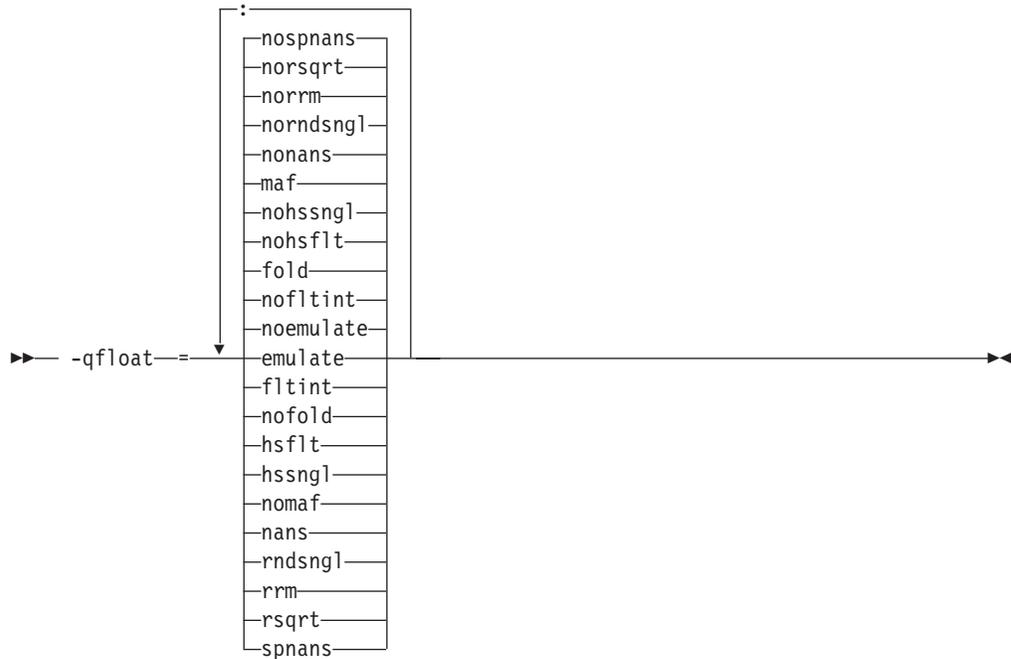
“Compiler Messages” on page 355

float

Purpose

Specifies various floating-point options. These options provide different strategies for speeding up or improving the accuracy of floating-point calculations.

Syntax



Option selections are described in the **Notes** section below. See also “#pragma options” on page 299.

Notes

Using **float** suboptions other than the default settings may produce varying results in floating point computations. Incorrect computational results may be produced if not all required conditions for a given suboption are met. For these reasons, you should only use this option if you are experienced with floating-point calculations involving IEEE floating-point values and can properly assess the possibility of introducing errors in your program.

You can specify one or more of the following **float** suboptions.

emulate noemulate	<p>Emulates the floating-point instructions omitted by the PowerPC 403™ processor. The default is float=noemulate.</p> <p>To emulate PowerPC 403 processor floating-point instructions, use -qfloat=emulate. Function calls are emitted in place of PowerPC 403 floating-point instructions. Use this option only in a single-threaded, stand-alone environment targeting the PowerPC 403 processor.</p> <p>Do not use -qfloat=emulate with any of the following:</p> <ul style="list-style-type: none"> • -qarch=pwr, -qarch=pwr2, -qarch=pwrx • -qlongdouble, -qldbl128 • xlC128 compiler invocation commands
fltint nofltint	<p>Speeds up floating-point-to-integer conversions by using faster inline code that does not check for overflows. The default is float=nofltint, which checks floating-point-to-integer conversions for out-of-range values.</p> <p>This suboption must only be used with an optimization option.</p> <ul style="list-style-type: none"> • With -O2 in effect, -qfloat=nofltint is the implied setting. • With -O3 and greater in effect, -qfloat=fltint is implied. <p>To include range checking in floating-point-to-integer conversions with the -O3 option, specify -qfloat=nofltint.</p> <ul style="list-style-type: none"> • -qnostrict sets -qfloat=fltint <p>Changing the optimization level will not change the setting of the fltint suboption if fltint has already been specified.</p> <p>For PWR2 and PPC family architectures, faster inline code is used that correctly handles out-of-range values.</p> <p>If the -qstrict -qnostrict and -qfloat= options conflict, the last setting is used.</p>
fold nofold	<p>Specifies that constant floating-point expressions are to be evaluated at compile time rather than at run time.</p> <p>The -qfloat=fold option replaces the obsolete -qfold option. Use -qfloat=fold in your new applications.</p>

hsflt
nohsflt

Note: The **hsflt** suboption is for specific applications in which floating-point computations have known characteristics. Using this option when you are compiling other application programs can produce incorrect results without warning. Also, using this option with **-qfloat=rndsngl**, **-q64**, or **-qarch=ppc** or any other PPC family architecture setting may produce incorrect results on rs64b or future systems.

The **hsflt** option speeds up calculations by truncating instead of rounding computed values to single precision before storing and on conversions from floating point to integer. The **nohsflt** suboption specifies that single-precision expressions are rounded after expression evaluation and that floating-point-to-integer conversions are to be checked for out-of-range values.

The **hsflt** suboption overrides the **rndsngl**, **nans**, and **spnans** suboptions.

The **-qfloat=hsflt** option replaces the obsolete **-qhsflt** option. Use **-qfloat=hsflt** in your new applications.

This option has little effect unless the **-qarch** option is set to **pwr**, **pwr2**, **pwrx**, **pwr2s** or, in 32-bit mode, **com**. For PPC family architectures, all single-precision (float) operations are rounded. This option only affects double-precision (double) expressions cast to single-precision (float).

hssngl
nohssngl

Note: Using this suboption with **-qfloat=rndsngl**, **-q64**, or **-qarch=ppc** or any other PPC family architecture setting may produce incorrect results on rs64b or future systems.

The **hssngl** option specifies that single-precision expressions are rounded only when the results are stored into **float** memory locations. The **nohssngl** option specifies that single-precision expressions are rounded after expression evaluation. Using **hssngl** can improve runtime performance and is safer than using **-qfloat=hsflt**.

This suboption has little effect unless the **-qarch** option is set to **pwr**, **pwr2**, **pwrx**, **pwr2s** or, in 32-bit mode, **com**. For PPC family architectures, all single-precision (float) operations are rounded. This option only affects double-precision (double) expressions cast to single-precision (float) and used in an assignment operator for which a store instruction is generated.

The **-qfloat=hssngl** option replaces the obsolete **-qhssngl** option. Use **-qfloat=hssngl** in your new applications.

maf
nomaf

Makes floating-point calculations faster and more accurate by using floating-point multiply-add instructions where appropriate. The results may not be exactly equivalent to those from similar calculations performed at compile time or on other types of computers. Negative zero results may be produced. This option may affect the precision of floating-point intermediate results.

The **-qfloat=maf** option replaces the obsolete **-qmaf** option. Use **-qfloat=maf** in your new applications.

nans nonans	<p>Generates extra instructions to detect signalling NaN (Not-a-Number) when converting from single-precision to double-precision at run time. The option nonans specifies that this conversion need not be detected. -qfloat=nans is required for full compliance to the IEEE 754 standard.</p> <p>The hsflt option overrides the nans option.</p> <p>When used with the -qflttrap or -qflttrap=invalid option, the compiler detects invalid operation exceptions in comparison operations that occur when one of the operands is a signalling NaN.</p> <p>The -qfloat=nans option replaces the obsolete -qfloat=spnans option and the -qspnans option. Use -qfloat=nans in your new applications.</p>
rndsngl norndsngl	<p>Specifies that the result of each single-precision (float) operation is to be rounded to single precision. -qfloat=norndsngl specifies that rounding to single-precision happens only after full expressions have been evaluated. Using this option may sacrifice speed for consistency with results from similar calculations on other types of computers.</p> <p>The hsflt suboption overrides the rndsngl option.</p> <p>This suboption has no effect unless the -qarch option is set to pwr, pwr2, pwrx, pwr2s or, in 32-bit mode, com. For PPC family architectures, all single-precision (float) operations are rounded.</p> <p>Using this option with -qfloat=hssngl or -qfloat=hsflt may produce incorrect results on rs64b or future systems.</p> <p>The -qfloat=rndsngl option replaces the obsolete -qrndsngl option. Use -qfloat=rndsngl in your new applications.</p>
rrm norrm	<p>Prevents floating-point optimizations that are incompatible with runtime rounding to plus and minus infinity modes. Informs the compiler that the floating-point rounding mode may change at run time or that the floating-point rounding mode is not <i>round to nearest</i> at run time.</p> <p>-qfloat=rrm must be specified if the Floating Point Status and Control register is changed at run time (as well as for initializing exception trapping).</p> <p>The -qfloat=rrm option replaces the obsolete -qrrm option. Use -qfloat=rrm in your new applications.</p>

`rsqrt`
`norsqrt` Specifies whether a sequence of code that involves division by the result of a square root can be replaced by calculating the reciprocal of the square root and multiplying. Allowing this replacement produces code that runs faster.

- For `-O2`, the default is `-qfloat=norsqrt`.
- For `-O3`, the default is `-qfloat=rsqrt`. Use `-qfloat=norsqrt` to override this default.
- `-qnostrict` sets `-qfloat=rsqrt`. (Note that `-qfloat=rsqrt` means that `errno` will *not* be set for any `sqrt` function calls.)
- `-qfloat=rsqrt` has no effect when `-qarch=pwr2` is also specified.
- `-qfloat=rsqrt` has no effect unless `-qignerrno` is also specified.

Changing the optimization level will not change the setting of the `rsqrt` option if `rsqrt` has already been specified. If the `-qstrict` | `-qnostrict` and `-qfloat= options` conflict, the last setting is used.

`spnans`
`nospnans` Generates extra instructions to detect signalling NaN on conversion from single-precision to double-precision. The option `nospnans` specifies that this conversion need not be detected.

The `hsflt` suboption overrides the `spnans` suboption.

The `-qfloat=spnans` option replaces the obsolete `-qfloat=spnans` and `-qspnans` options. Use `-qfloat=spnans` in your new applications.

Example

To compile `myprogram.c` so that constant floating point expressions are evaluated at compile time and multiply-add instructions are not generated, enter:

```
xlc myprogram.c -qfloat=fold:nomaf
```

Related References

“Compiler Command Line Options” on page 51

“arch” on page 70

“flttrap” on page 123

“ldbl128, longdouble” on page 172

“rrm” on page 224

“strict” on page 241

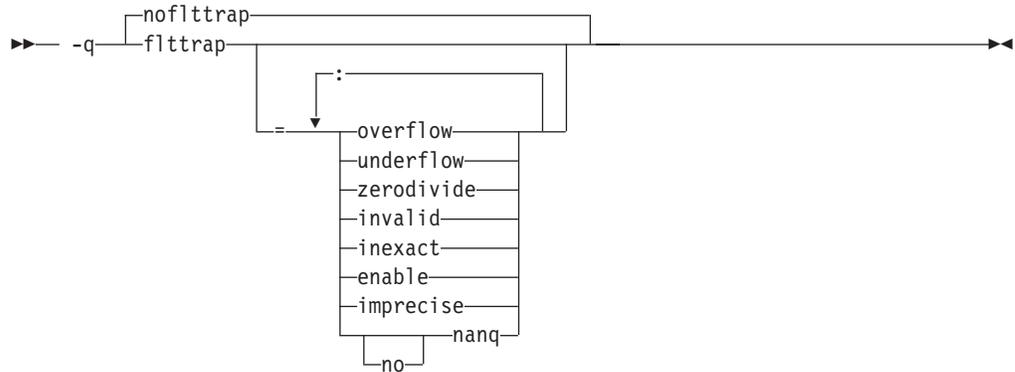
“#pragma options” on page 299

flttrap

Purpose

Generates extra instructions to detect and trap run-time floating-point exceptions.

Syntax



where suboptions do the following:

ENable	Enables the specified exceptions in the prologue of the main program. With the exception of nanq (described below), this suboption is required if you want to turn on exception trapping options listed below without modifying the source code.
Overflow	Generates code to detect and trap floating-point overflow.
UNDERflow	Generates code to detect and trap floating-point underflow.
ZERODivide	Generates code to detect and trap floating-point division by zero.
INValid	Generates code to detect and trap floating-point invalid operation exceptions.
INEXact	Generates code to detect and trap floating-point inexact exceptions.
IMPrecise	Generates code for imprecise detection of the specified exceptions. If an exception occurs, it is detected, but the exact location of the exception is not determined.
nanq	Generates code to detect and trap NaNQ (Not a Number Quiet) exceptions handled by or generated by floating point operations. The nanq and nonanq settings are not affected by -qnoflttrap , -qflttrap , or -qflttrap=enable .

See also “#pragma options” on page 299.

Notes

This option is recognized during linking. **-qnoflttrap** specifies that these extra instructions need not be generated.

Specifying the **-qflttrap** option with no suboptions is equivalent to setting **-qflttrap=overflow:underflow:zerodivide:invalid:inexact**. The exceptions are not automatically enabled, and all floating-point operations are checked to provide precise exception-location information.

If specified with **#pragma options**, the **-qnoflttrap** option *must* be the first option specified.

If your program contains signalling NaNs, you should use the `-qfloat=nans` along with `-qflttrap` to trap any exceptions.

The compiler exhibits behavior as illustrated in the following examples when the `-qflttrap` option is specified together with `-qoptimize` options:

- with `-O2`:
 - `1/0` generates a `div0` exception and has a result of infinity
 - `0/0` generates an invalid operation
- with `-O3` or greater:
 - `1/0` generates a `div0` exception and has a result of infinity
 - `0/0` returns zero multiplied by the result of the previous division.

Example

To compile `myprogram.c` so that floating-point overflow and underflow and divide by zero are detected, enter:

```
xlc myprogram.c -qflttrap=overflow:underflow:zerodivide:enable
```

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“O, optimize” on page 194

fold

Purpose

Specifies that constant floating-point expressions are to be evaluated at compile time.

Syntax

►► — -q —

fold
nofold

 —►►

See also “#pragma options” on page 299.

Notes

This option is obsolete. Use -qfloat=fold in your new applications.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

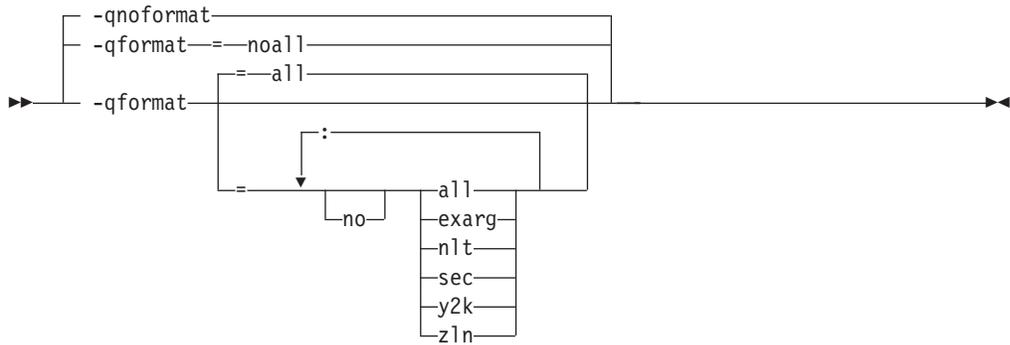
“#pragma options” on page 299

format

Purpose

Warns of possible problems with string input and output format specifications. Functions diagnosed are `printf`, `scanf`, `strftime`, `strfmon` family functions and functions marked with format attributes.

Syntax



where suboptions are:

- `all` Turns on all format diagnostic messages.
- `exarg` Warns if excess arguments appear in `printf` and `scanf` style function calls.
- `nlt` Warns if a format string is not a string literal, unless the format function takes its format arguments as a `va_list`.
- `sec` Warns of possible security problems in use of format functions.
- `y2k` Warns of `strftime` formats that produce a 2-digit year.
- `zln` Warns of zero-length formats.

Note: Specifying `no` in front of any of the above suboptions disables that group of diagnostic messages. For example, to turn off diagnostic messages for `y2k` warnings, specify `-qformat=noy2k` on the command line.

Notes

If `-qformat` is *not* specified on the command line, the compiler assumes a default setting of `-qnoformat`, which is equivalent to `-qformat=noall`.

If `-qformat` is specified on the command line without any suboptions, the compiler assumes a default setting of `-qformat=all`.

Examples

1. To enable all format string diagnostics, enter either of the following:

```
xlc myprogram.c -qformat=all
xlc myprogram.c -qformat
```
2. To enable all format diagnostic checking except that for `y2k` date diagnostics, enter:

```
xlc myprogram.c -qformat=all:noy2k
```

Related References

“Compiler Command Line Options” on page 51

fullpath

Purpose

Specifies what path information is stored for files when you use the **-g** compiler option.

Syntax

►► -q nofullpath
fullpath _____►►

Notes

Using **-qfullpath** causes the compiler to preserve the absolute (full) path name of source files specified with the **-g** option.

The relative path name of files is preserved when you use **-qnofullpath**.

-qfullpath is useful if the executable file was moved to another directory. If you specified **-qnofullpath**, the debugger would be unable to find the file unless you provide a search path in the debugger. Using **-qfullpath** would locate the file successfully.

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

funcsect

Purpose

Places instructions for each function in a separate object file control section or csect. By default, each object file will consist of a single control section combining all functions defined in the corresponding source file.

Syntax

►► — -q — nofuncsect
funcsect —————►►

Notes

Using multiple csects increases the size of the object file, but can reduce the size of the final executable by allowing the linkage editor to remove functions that are not called or that have been inlined by the optimizer at all places they are called.

If the file contains initialized static data or the pragma statement

```
#pragma comment (copyright)
```

some functions will be one machine word larger.

If this option is specified in **#pragma options**, the pragma directive must be specified before the first statement in the compilation unit.

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

G

Purpose

Tells the linkage editor to create a shared object enabled for runtime linking.

Syntax

►— -G —►

Notes

The compiler will automatically export all global symbols from the shared object unless you specify which symbols to export by using **-bE**;, **-bexport**;, **-bexpall** or **-bnoexpall**.

If you use **-G** to create a shared library, the compiler will:

1. If you don't specify **-bE**;, **-bexport**;, **-bexpall** or **-bnoexpall**, create an export list containing all global symbols using the CreateExportList script. You can specify another script with the **-tE/-B** or **-qpath=E**: options.
2. Save the export list if either the CreateExportList utility or the **-qmkshrobj** and **-qexpfile** compiler options were used to create an export list.
3. Call the linker with the appropriate options and object files to build a shared object.

This is a linkage editor (**ld**) option. Refer to your operating system documentation for a description of **ld** command usage and syntax.

For more information about the CreateExportList utility and creating shared objects, see Creating a Library in the Programming Guide.

Related References

"Compiler Command Line Options" on page 51

"B" on page 78

"b" on page 79

"brtl" on page 82

"expfile" on page 112

"path" on page 202

"t" on page 246

Also, on the Web see:

Shared Objects and Runtime Linking chapter in General Programming Concepts:

Writing and Debugging Programs

ld Command section in Commands Reference, Volume 3: i through m

g

Purpose

Generates information used for debugging tools such as the IBM Distributed Debugger.

Syntax

▶— -g —▶

Notes

Specifying **-g** will turn off all inlining unless you explicitly request it. For example:

Options	Effect on inlining
-g	No inlining.
-O	Inline declared functions.
-O -Q	Inline declared functions and auto inline others.
-g -O	No inlining.
-g -O -Q	Inline declared functions and auto inline others.

The default with **-g** is not to include information about unreferenced symbols in the debugging information.

To include information about both referenced and unreferenced symbols, use the **-qdbxextra** option with **-g**.

To specify that source files used with **-g** are referred to by either their absolute or their relative path name, use **-qfullpath**.

You can also use the **-qlinedebug** option to produce abbreviated debugging information in a smaller object size.

Example

To compile `myprogram.c` to produce an executable program testing so you can debug it, enter:

```
xlc myprogram.c -o testing -g
```

To compile `myprogram.c` to produce an executable program named `testing_all`, and containing additional information about unreferenced symbols so you can debug it, enter:

```
xlc myprogram.c -o testing_all -g -qdbxextra
```

Related References

“Compiler Command Line Options” on page 51

“dbxextra” on page 100

“fullpath” on page 127

“linedebug” on page 174

“O, optimize” on page 194

“Q” on page 214

genproto

Purpose

Produces ANSI prototypes from K&R function definitions. This should help to ease the transition from K&R to ANSI.

Syntax



Notes

Using **-qgenproto** without **parmnames** will cause prototypes to be generated without parameter names. Parameter names are included in the prototype when **parmnames** is specified.

Example

For the following function, `foo.c`:

```
foo(a,b,c)
float a;
int *b;
int c;
```

specifying

```
xlc -c -qgenproto foo.c
```

produces

```
int foo(double, int*, int);
```

The parameter names are dropped. On the other hand, specifying

```
xlc -c -qgenproto=parmnames foo.c
```

produces

```
int foo(double a, int* b, int c);
```

In this case the parameter names are kept.

Note that **float a** is represented as **double** or **double a** in the prototype, since ANSI states that all narrow-type arguments (such as **chars**, **shorts**, and **floats**) are widened before they are passed to K&R functions.

Related References

“Compiler Command Line Options” on page 51

halt

Purpose

Instructs the compiler to stop after the compilation phase when it encounters errors of specified *severity* or greater.

Syntax



where severity levels in order of increasing severity are:

<i>severity</i>	Description
i	Information
w	Warning
e	Error
s	Severe error
u	Unrecoverable error

See also “#pragma options” on page 299.

Notes

When the compiler stops as a result of the **-qhalt** option, the compiler return code is nonzero.

When **-qhalt** is specified more than once, the lowest severity level is used.

The **-qhalt** option can be overridden by the **-qmaxerr** option.

Diagnostic messages may be controlled by the **-qflag** option.

Example

To compile myprogram.c so that compilation stops if a **warning** or higher level message occurs, enter:

```
xlc myprogram.c -qhalt=w
```

Related References

“Compiler Command Line Options” on page 51

“flag” on page 117

“maxerr” on page 188

“#pragma options” on page 299

heapdebug

Purpose

Enables debug versions of memory management functions.

Syntax



Notes

By default, the compiler uses the regular memory management functions (**calloc**, **malloc**, **new**, etc.) and does not preinitialize their local storage. The header files for the regular memory management functions are found in the system include directory.

The header files declaring the debug versions of these memory management functions, such as **_debug_calloc**, **_debug_malloc**, **new**, and so on, are found in the product include directory at `usr/vac/include`.

Specifying **-qheapdebug** changes the header file search order so the compiler can find the headers declaring the debug versions of the memory management functions. The compiler will search for header files first in the product include directory, where header files for the debug versions of memory management functions are stored, and then in the system include directory.

Specifying **-qheapdebug** also defines the `__DEBUG_ALLOC__` macro.

Example

To compile `myprogram.c` with the debug versions of memory management functions, enter:

```
xlc -qheapdebug myprogram.c -o testing
```

Related References

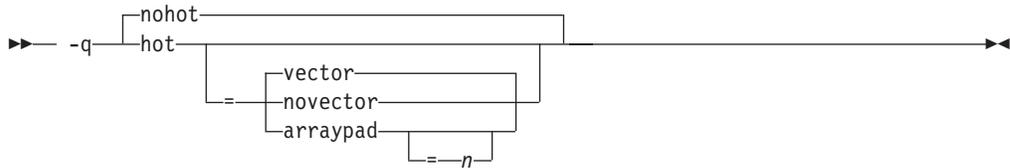
“Compiler Command Line Options” on page 51

hot

Purpose

Instructs the compiler to perform high-order loop analysis and transformations during optimization.

Syntax



where:

- arraypad** The compiler will pad any arrays where it infers there may be a benefit and will pad by whatever amount it chooses. Not all arrays will necessarily be padded, and different arrays may be padded by different amounts.
- arraypad=*n*** The compiler will pad every array in the code. The pad amount must be a positive integer value, and each array will be padded by an integral number of elements. Because *n* is an integral value, we recommend that pad values be multiples of the largest array element size, typically 4, 8, or 16.
- vector | novector** The compiler converts certain operations that are performed in a loop on successive elements of an array (for example, square root, reciprocal square root) into a call to a library routine. This call will calculate several results at one time, which is faster than calculating each result sequentially.

If you specify **-qhot=novector**, the compiler performs high-order transformations on loops and arrays, but avoids optimizations where certain code is replaced by calls to vector library routines. The **-qhot=vector** option may affect the precision of your program's results so you should specify either **-qhot=novector** or **-qstrict** if the change in precision is unacceptable to you.

Default

The **-qhot=vector** suboption is on by default when you specify the **-qhot**, **-qsmp**, **-O4**, or **-O5** options.

Notes

If you do not also specify optimization of at least level 2 when specifying **-qhot** on the command line, the compiler assumes **-O2**.

Because of the implementation of the cache architecture, array dimensions that are powers of two can lead to decreased cache utilization. The optional **arraypad** suboption permits the compiler to increase the dimensions of arrays where doing so might improve the efficiency of array-processing loops. If you have large arrays with some dimensions (particularly the first one) that are powers of 2, or if you find that your array-processing programs are slowed down by cache misses or page faults, consider specifying **-qhot=arraypad**.

The **vector** suboption optimizes array data to run mathematical operations in parallel where applicable. The compiler uses standard registers with no vector size restrictions. The **vector** suboption supports single and double-precision floating-point mathematics, and is useful for applications with significant mathematical processing demands.

Both **-qhot=arraypad** and **-qhot=arraypad=*n*** are unsafe options. They do not perform any checking for reshaping or equivalences that may cause the code to break if padding takes place.

Example

The following example turns on the **-qhot=arraypad** option:

```
xlc -qhot=arraypad myprogram.c
```

Related References

“Compiler Command Line Options” on page 51

“C” on page 83

“O, optimize” on page 194

“smp” on page 232

hsflt

Purpose

Speeds up calculations by removing range checking on single-precision **float** results, and on conversions from floating point to integer. **-qnohsflt** specifies that single-precision expressions are rounded after expression evaluation, and that floating-point-to-integer conversions are to be checked for out of range values.

Syntax

►► -q nohsflt
hsflt _____►►

See also “#pragma options” on page 299.

Notes

*This option is obsolete. Use **-qfloat=hsflt** in your new applications.*

The **-qhsflt** option overrides the **-qrndsngl** and **-qspnans** options.

The **-qhsflt** option is intended for specific applications in which floating-point computations have known characteristics. Using this option when compiling other application programs can produce incorrect results without warning.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“spnans” on page 237

“#pragma options” on page 299

hssngl

Purpose

Specifies that single-precision expressions are rounded only when the results are stored into **float** memory locations. **-qnohssngl** specifies that single-precision expressions are rounded after expression evaluation. Using **-qhssngl** can improve run-time performance.

Syntax

►► -q nohssngl hssngl _____►►

See also “#pragma options” on page 299.

Notes

This option is obsolete. Use **-qfloat=hssngl** in your new applications.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“#pragma options” on page 299

Purpose

Specifies an additional search path for **#include** filenames that do not specify an absolute path.

Syntax

►— `-I—directory`—►

Notes

The value for *directory* must be a valid path name (for example, `/u/golnaz`, or `/tmp`, or `./subdir`). The compiler appends a slash (`/`) to the directory and then concatenates it with the file name before doing the search. The path *directory* is the one that the compiler searches first for **#include** files whose names do not start with a slash (`/`). If *directory* is not specified, the default is to search the standard directories.

If the `-I` *directory* option is specified both in the configuration file and on the command line, the paths specified in the configuration file are searched first.

The `-I` *directory* option can be specified more than once on the command line. If you specify more than one `-I` option, directories are searched in the order that they appear on the command line. See *Directory Search Sequence for Include Files Using Relative Path Names* for more information about searching directories.

If you specify a full (absolute) path name on the **#include** directive, this option has no effect.

Example

To compile `myprogram.c` and search `/usr/tmp` and then `/oldstuff/history` for included files, enter:

```
xlc myprogram.c -I/usr/tmp -I/oldstuff/history
```

Related Tasks

“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

Related References

“Compiler Command Line Options” on page 51

idirfirst

Purpose

Specifies the search order for files included with the **#include** “file_name” directive.

Syntax

►— -q— noidirfirst
idirfirst —►

See also “#pragma options” on page 299.

Notes

Use **-qidirfirst** with the **-I** option.

The normal search order (for files included with the **#include** “file_name” directive) *without* the **idirfirst** option is:

1. Search the directory where the current source file resides.
2. Search the directory or directories specified with the **-I** option.
3. Search the standard include directory **/usr/include**.

With **-qidirfirst**, the directories specified with the **-I** option are searched before the directory where the current file resides.

-qidirfirst has no effect on the search order for the **#include** <file_name> directive.

-qidirfirst is independent of the **-qnostdinc** option, which changes the search order for both **#include** “file_name” and **#include** <file_name>.

The search order of files is described in *Directory Search Sequence for Include Files Using Relative Path Names*.

The last valid **#pragma options** [no]idirfirst remains in effect until replaced by a subsequent **#pragma options** [no]idirfirst.

Example

To compile myprogram.c and search **/usr/tmp/myinclude** for included files before searching the current directory (where the source file resides), enter:

```
xlc myprogram.c -I/usr/tmp/myinclude -qidirfirst
```

Related Tasks

“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

Related References

“Compiler Command Line Options” on page 51

“I” on page 138

“stdinc” on page 240

“#pragma options” on page 299

ignerrno

Purpose

Allows the compiler to perform optimizations that assume **errno** is not modified by system calls.

Syntax

►► -q noignerrno
ignerrno ◄◄

See also “#pragma options” on page 299.

Notes

Some system library routines set **errno** when an exception occurs. This setting and subsequent side effects of **errno** may be ignored by specifying **-qignerrno**.

Specifying a **-O3** or greater optimization option will also set **-qignerrno**. If you require both optimization and the ability to set **errno**, you should specify **-qnoignerrno** after the optimization option on the command line.

Related References

“Compiler Command Line Options” on page 51

“O, optimize” on page 194

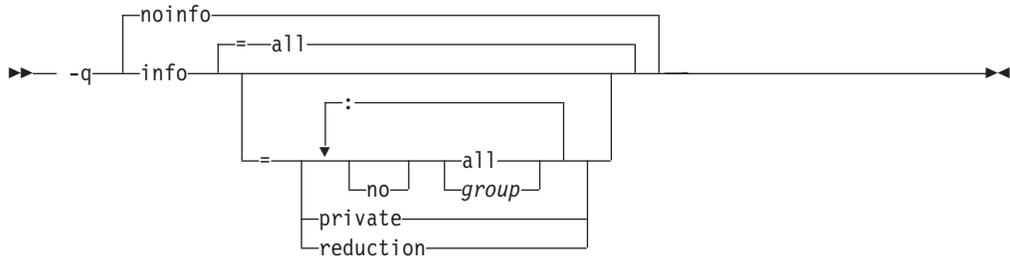
“#pragma options” on page 299

info

Purpose

Produces informational messages.

Syntax



where **-qinfo** options and diagnostic message groups are described in the **Notes** section below. See also “#pragma info” on page 287 and “#pragma options” on page 299.

Defaults

If you do not specify **-qinfo** on the command line, the compiler assumes:

1. **-qnoinfo**

If you specify **-qinfo** on the command line without any suboptions, the compiler assumes:

1. **-qinfo=all**

Notes

Specifying **-qinfo=all** or **-qinfo** with no suboptions turns on all diagnostic messages for all groups.

Specifying **-qnoinfo** or **-qinfo=noall** turns off all diagnostic messages for all groups.

You can use the **#pragma options info=suboption[:suboption ...]** or **#pragma options noinfo** forms of this compiler option to temporarily enable or disable messages in one or more specific sections of program code.

Available forms of the **-qinfo** option are:

- | | |
|-----------|--|
| all | Turns on all diagnostic messages for all groups. |
| noall | The -qinfo and -qinfo=all forms of the option have the same effect. |
| private | Turns off all diagnostic messages for specific portions of your program. |
| reduction | Lists shared variables made private to a parallel loop. |
| | Lists all variables that are recognized as reduction variables inside a parallel loop. |

group Turns on or off specific groups of messages, where *group* can be one or more of:

<i>group</i>	Type of messages returned or suppressed
c99 noc99	C code that may behave differently between C89 and C99 language levels.
cmp nocmp	Possible redundancies in unsigned comparisons
cnd nocnd	Possible redundancies or problems in conditional expressions
cns nocns	Operations involving constants
cnv nocnv	Conversions
dcl nodcl	Consistency of declarations
eff noeff	Statements and pragmas with no effect
enu noenu	Consistency of enum variables
ext noext	Unused external definitions
gen nogen	General diagnostic messages
gnr nognr	Generation of temporary variables
got nogot	Use of goto statements
ini noini	Possible problems with initialization
inl noinl	Functions not inlined
lan nolan	Language level effects
obs noobs	Obsolete features
ord noord	Unspecified order of evaluation
par nopar	Unused parameters
por nopor	Nonportable language constructs
ppc noppc	Possible problems with using the preprocessor
ppt noppt	Trace of preprocessor actions
pro nopro	Missing function prototypes
rea norea	Code that cannot be reached
ret noret	Consistency of return statements
trd notrd	Possible truncation or loss of data or precision
tru notru	Variable names truncated by the compiler
trx notrx	Hexadecimal floating point constants rounding
uni nouni	Uninitialized variables
upg noupg	Generates messages describing new behaviors of the current compiler release as compared to the previous release.
use nouse	Unused auto and static variables
zea nozea	Zero-extent arrays.

Example

To compile `myprogram.c` to produce informational message about all items except conversions and unreachable statements, enter:

```
xlc myprogram.c -qinfo=all -qinfo=nocnv:norea
```

Related Concepts

“Reduction Operations in Parallelized Loops” on page 14

“Shared and Private Variables in a Parallel Environment” on page 15

Related References

“Compiler Command Line Options” on page 51

“halt” on page 132

“suppress” on page 243

“#pragma info” on page 287

“#pragma options” on page 299

initauto

Purpose

Initializes automatic variables to the two-digit hexadecimal byte value *hex_value*.

Syntax

►► -q noinitauto
initauto=*hex_value* ►►

See also “#pragma options” on page 299.

Notes

The option generates extra code to initialize the value of automatic variables. It reduces the runtime performance of the program and should only be used for debugging.

There is no default setting for the initial value of **-qinitauto**. You must set an explicit value (for example, **-qinitauto=FA**).

Example

To compile `myprogram.c` so that automatic variables are initialized to hex value FF (decimal 255), enter:

```
xlc myprogram.c -qinitauto=FF
```

Related References

“Compiler Command Line Options” on page 51

inlglue

Purpose

Generates fast external linkage by inlining the pointer glue code necessary to make a call to an external function or a call through a function pointer.

Syntax

►► -q noinlglue
inlglue ◄◄

See also “#pragma options” on page 299.

Notes

Glue code, generated by the linker, is used for passing control between two external functions, or when you call functions through a pointer. Therefore the **-qinlglue** option only affects function calls through pointers or calls to an external compilation unit. For calls to an external function, you should specify that the function is imported by using, for example, the **-qprocimported** option.

The inlining of glue code can cause the size of code to grow. This can be overridden by specifying the **-qcompact** option, thereby disabling the **-qinlglue** option.

Related References

“Compiler Command Line Options” on page 51

“compact” on page 91

“proclocal, procimported, procunknown” on page 211

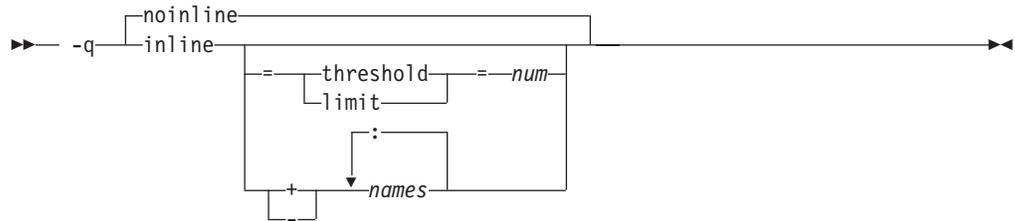
“#pragma options” on page 299

inline

Purpose

Attempts to inline functions instead of generating calls to those functions. Inlining is performed if possible but, depending on which optimizations are performed, some functions might not be inlined.

Syntax



The following **-qinline** options apply:

- qinline** The compiler attempts to inline all appropriate functions with 20 executable source statements or fewer, subject to any other settings of the suboptions to the **-qinline** option. If **-qinline** is specified last, all functions are inlined.
- qinline=threshold=num** Sets a size limit on the functions to be inlined. The number of executable statements must be less than or equal to *num* for the function to be inlined. *num* must be a positive integer. The default value is 20. Specifying a threshold value of 0 causes no functions to be inlined except those functions marked with the `__inline`, `_Inline`, or `_inline` keywords.

The *num* value applies to logical C statements. Declarations are not counted, as you can see in the example below:

```
increment()
{
    int a, b, i;
    for (i=0; i<10; i++) /* statement 1 */
    {
        a=i;           /* statement 2 */
        b=i;           /* statement 3 */
    }
}
```

- qinline-names** The compiler does not inline functions listed by *names*. Separate each *name* with a colon (:). All other appropriate functions are inlined. The option implies **-qinline**.

For example:

```
-qinline-salary:taxes:expenses:benefits
```

causes all functions except those named **salary**, **taxes**, **expenses**, or **benefits** to be inlined if possible.

A warning message is issued for functions that are not defined in the source file.

<code>-qinline+names</code>	<p>Attempts to inline the functions listed by <i>names</i> and any other appropriate functions. Each <i>name</i> must be separated by a colon (:). The option implies -qinline.</p> <p>For example,</p> <pre style="margin-left: 40px;">-qinline+food:clothes:vacation</pre> <p>causes all functions named food, clothes, or vacation to be inlined if possible, along with any other functions eligible for inlining.</p> <p>A warning message is issued for functions that are not defined in the source file or that are defined but cannot be inlined.</p> <p>This suboption overrides any setting of the <i>threshold</i> value. You can use a threshold value of zero along with -qinline+names to inline specific functions. For example:</p> <pre style="margin-left: 40px;">-qinline=threshold=0</pre> <p>followed by:</p> <pre style="margin-left: 40px;">-qinline+salary:taxes:benefits</pre> <p>causes <i>only</i> the functions named salary, taxes, or benefits to be inlined, if possible, and no others.</p>
<code>-qinline=limit=num</code>	<p>Specifies the maximum size (in bytes of generated code) to which a function can grow due to inlining. This limit does not affect the inlining of user specified functions.</p>
<code>-qnoinline</code>	<p>Does not inline any functions. If -qnoinline is specified last, no functions are inlined.</p>

Default

The default is to treat inline specifications as a hint to the compiler, and the result depends on other options that you select:

- If you optimize your program using one of the **-O** compiler options, the compiler attempts to inline all functions declared as inline. Otherwise, the compiler attempts to inline only some of the simpler functions declared as inline.

Notes

The **-qinline** option is functionally equivalent to the **-Q** option.

If you specify the **-g** option (to generate debug information), inlining may be affected. See the information for the **-g** compiler option.

Because inlining does not always improve run time, you should test the effects of this option on your code. Do not attempt to inline recursive or mutually recursive functions.

Normally, application performance is optimized if you request optimization (**-O** option), and compiler performance is optimized if you do not request optimization.

To maximize inlining, specify optimization (**-O**) and also specify the appropriate **-qinline** options.

The XL C Enterprise Edition (**inline**, **_inline**, **_Inline**, and **__inline**) C language keywords override all **-qinline** options except **-qnoinline**. The compiler will try to inline functions marked with these keywords regardless of other **-qinline** option settings.

The inline, _Inline, _inline, and __inline Function Specifiers: The compiler provides a set of keywords that you can use to specify functions that you want the compiler to inline. The keywords, which are functionally equivalent, are:

- **inline** (C99)
- **_Inline** (C only)
- **_inline** (C only)
- **__inline** (C)

For example:

```
    _Inline int catherine(int a);
```

suggests to the compiler that function `catherine` be inlined, meaning that code is generated for the function, rather than a function call.

Using the inline specifiers with `data` or to declare the `main()` function generates an error.

By default, function inlining is turned off, and functions qualified with inline specifiers are treated simply as extern functions. To turn on function inlining, specify either the **-qinline** or **-Q** compiler options. Inlining is also turned on if you turn optimization on with the **-O** or **-qoptimize** compiler option.

Recursive functions (functions that call themselves) are inlined for the first occurrence only. The call to the function from within itself is not inlined.

You can also use the **-qinline** or **-Q** compiler options to automatically inline all functions smaller than a specified size. For best performance, however, use the inline keywords to choose the functions you want to inline rather than using automatic inlining.

An inline function can be declared and defined simultaneously. If it is declared with one of the inline keywords, it can be defined without any inline keywords. The following code fragment shows an inline function definition. Note that the inline keywords are specified for both a function definition and function declaration.

```
    __inline int add(int i, int j) { return i + j; }

    inline double fahr(double t);
```

Note: The use of the inline specifier does not change the meaning of the function, but inline expansion of a function may not preserve the order of evaluation of the actual arguments.

Example

To compile `myprogram.c` so that no functions are inlined, enter:

```
xlc myprogram.c -O -qnoinline
```

To compile `myprogram.c` so that the compiler attempts to inline functions of fewer than 12 lines, enter:

```
xlc myprogram.c -O -qinline=12
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

“O, optimize” on page 194

“Q” on page 214

ipa

Purpose

Turns on or customizes a class of optimizations known as interprocedural analysis (IPA).

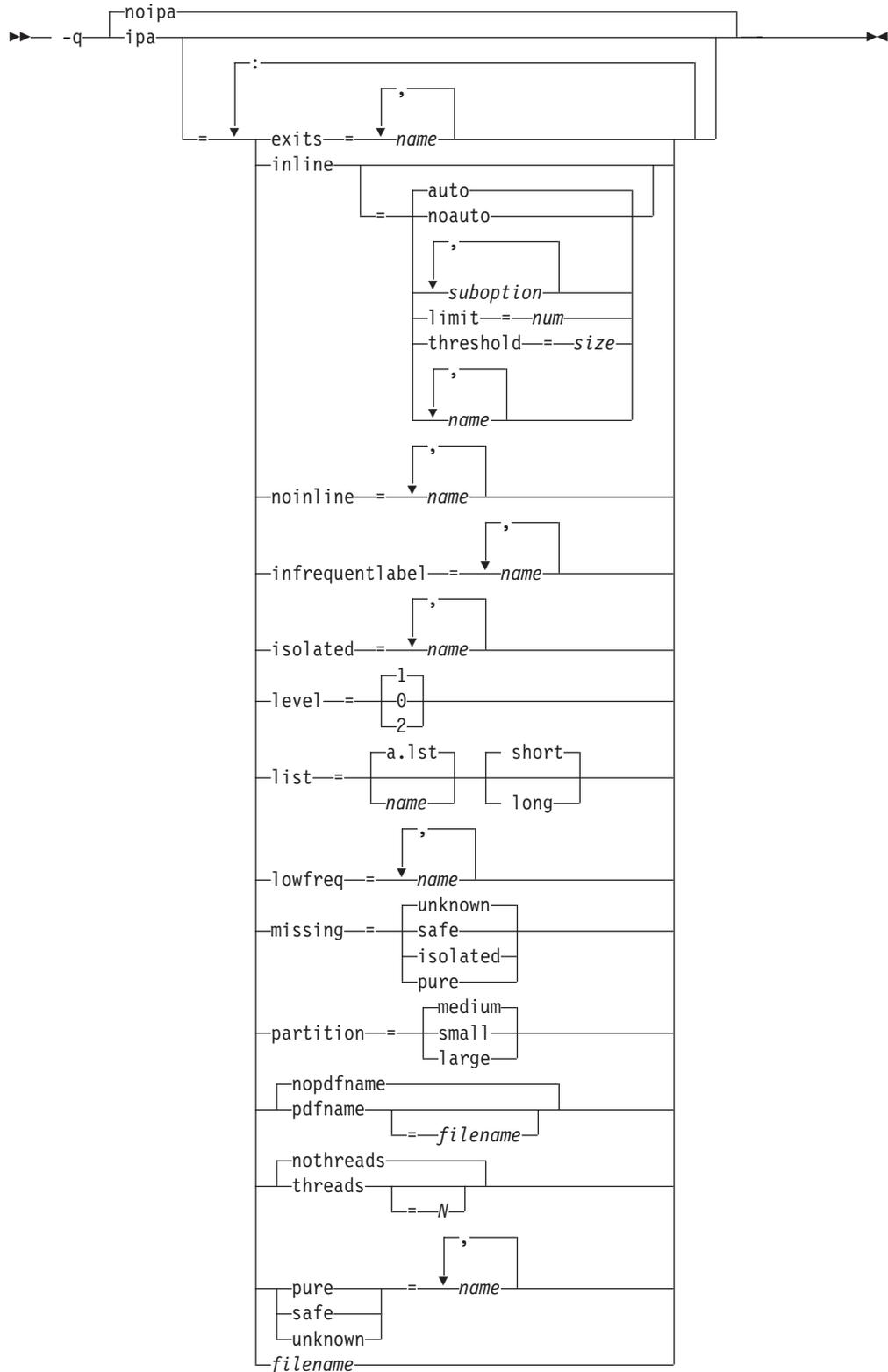
Compile-time syntax



where:

-qipa Compile-time Options	Description
-qipa	Activates interprocedural analysis with the following -qipa <i>suboption</i> defaults: <ul style="list-style-type: none">• inline=auto• level=1• missing=unknown• partition=medium
-qipa=object -qipa=noobject	Specifies whether to include standard object code in the object files. Specifying the noobject suboption can substantially reduce overall compile time by not generating object code during the first IPA phase. If the -S compiler option is specified with noobject , noobject is ignored. If compilation and linking are performed in the same step, and neither the -S nor any listing option is specified, -qipa=noobject is implied by default. If any object file used in linking with -qipa was created with the -qipa=noobject option, any file containing an entry point (the main program for an executable program, or an exported function for a library) must be compiled with -qipa .

Link-time syntax



where:

-qipa Link-time Options	Description
-qnoipa	Deactivates interprocedural analysis.
-qipa	Activates interprocedural analysis with the following -qipa <i>suboption</i> defaults: <ul style="list-style-type: none"> • inline=auto • level=1 • missing=unknown • partition=medium

Suboptions can also include one or more of the forms shown below. Separate multiple suboptions with commas.

Link-time Suboptions	Description
exits= <i>name</i> {, <i>name</i> }	Specifies names of functions which represent program exits. Program exits are calls which can never return and can never call any procedure which has been compiled with IPA pass 1.
inline=auto inline=noauto	Enables or disables automatic inlining only. The compiler still accepts user-specified functions as candidates for inlining.
inline[= <i>suboption</i>]	Same as specifying the -qinline compiler option, with <i>suboption</i> being any valid -qinline suboption.
inline=limit= <i>num</i>	Changes the size limits that the -Q option uses to determine how much inline expansion to do. This established limit is the size below which the calling procedure must remain. <i>num</i> is the optimizer's approximation of the number of bytes of code that will be generated. Larger values for this number allow the compiler to inline larger subprograms, more subprogram calls, or both. This argument is implemented only when inline=auto is on.
inline=threshold= <i>size</i>	Specifies the upper size limit of functions to be inlined, where <i>size</i> is a value as defined under inline=limit . This argument is implemented only when inline=auto is on.
inline= <i>name</i> {, <i>name</i> }	Specifies a comma-separated list of functions to try to inline, where functions are identified by <i>name</i> .
noinline= <i>name</i> {, <i>name</i> }	Specifies a comma-separated list of functions that must not be inlined, where functions are identified by <i>name</i> .
infrequentlabel= <i>name</i> {, <i>name</i> }	Specifies a list of user-defined labels that will be infrequently branched-to during regular execution of the program.
isolated= <i>name</i> {, <i>name</i> }	Specifies a list of <i>isolated</i> functions that are not compiled with IPA. Neither isolated functions nor functions within their call chain can refer to global variables.

Link-time Suboptions	Description
level=0 level=1 level=2	Specifies the optimization level for interprocedural analysis. The default level is 1. Valid levels are as follows: <ul style="list-style-type: none"> • Level 0 - Does only minimal interprocedural analysis and optimization. • Level 1 - Turns on inlining, limited alias analysis, and limited call-site tailoring. • Level 2 - Performs full interprocedural data flow and alias analysis.
list list=[<i>name</i>] [short long]	Specifies that a listing file be generated during the link phase. The listing file contains information about transformations and analyses performed by IPA, as well as an optional object listing generated by the back end for each partition. This option can also be used to specify the name of the listing file. <p>If listings have been requested (using either the -qlist or -qipa=list options), and <i>name</i> is not specified, the listing file name defaults to a.lst.</p> <p>The long and short suboptions can be used to request more or less information in the listing file. The short suboption, which is the default, generates the Object File Map, Source File Map and Global Symbols Map sections of the listing. The long suboption causes the generation of all of the sections generated through the short suboption, as well as the Object Resolution Warnings, Object Reference Map, Inliner Report and Partition Map sections.</p>
lowfreq= <i>name</i> {, <i>name</i> }	Specifies names of functions which are likely to be called infrequently. These will typically be error handling, trace, or initialization functions. The compiler may be able to make other parts of the program run faster by doing less optimization for calls to these functions.
missing= <i>attribute</i>	Specifies the interprocedural behavior of procedures that are not compiled with -qipa and are not explicitly named in an unknown , safe , isolated , or pure suboption. <p>The following attributes may be used to refine this information:</p> <ul style="list-style-type: none"> • safe - Functions which do not indirectly call a visible (not missing) function either through direct call or through a function pointer. • isolated - Functions which do not directly reference global variables accessible to visible functions. Functions bound from shared libraries are assumed to be <i>isolated</i>. • pure - Functions which are <i>safe</i> and <i>isolated</i> and which do not indirectly alter storage accessible to visible functions. <i>pure</i> functions also have no observable internal state. • unknown - The default setting. This option greatly restricts the amount of interprocedural optimization for calls to <i>unknown</i> functions. Specifies that the missing functions are not known to be <i>safe</i>, <i>isolated</i>, or <i>pure</i>.

Link-time Suboptions	Description
partition=small partition=medium partition=large	Specifies the size of each program partition created by IPA during pass 2.
nopdfname pdfname pdfname= <i>filename</i>	Specifies the name of the profile data file containing the PDF profiling information. If you do not specify <i>filename</i> , the default file name is ._pdf . The profile is placed in the current working directory or in the directory named by the PDFDIR environment variable. This lets you do simultaneous runs of multiple executables using the same PDFDIR , which can be useful when tuning with PDF on dynamic libraries.
nothreads threads threads= <i>N</i>	Specifies the number of threads the compiler assigns to code generation. Specifying nothreads is equivalent to running one serial process. This is the default. Specifying threads allows the compiler to determine how many threads to use, depending on the number of processors available. Specifying threads=<i>N</i> instructs the program to use <i>N</i> threads. Though <i>N</i> can be any integer value in the range of 1 to MAXINT, <i>N</i> is effectively limited to the number of processors available on your system.
pure= <i>name</i> {, <i>name</i> }	Specifies a list of <i>pure</i> functions that are not compiled with -qipa . Any function specified as <i>pure</i> must be <i>isolated</i> and <i>safe</i> , and must not alter the internal state nor have side-effects, defined as potentially altering any data visible to the caller.
safe= <i>name</i> {, <i>name</i> }	Specifies a list of <i>safe</i> functions that are not compiled with -qipa and do not call any other part of the program. Safe functions can modify global variables, but may not call functions compiled with -qipa .
unknown= <i>name</i> {, <i>name</i> }	Specifies a list of <i>unknown</i> functions that are not compiled with -qipa . Any function specified as <i>unknown</i> can make calls to other parts of the program compiled with -qipa , and modify global variables and dummy arguments.

Link-time Suboptions	Description
<i>filename</i>	<p>Gives the name of a file which contains suboption information in a special format.</p> <p>The file format is the following:</p> <pre># ... comment attribute{, attribute} = name{, name} missing = attribute{, attribute} exits = name{, name} lowfreq = name{, name} inline [= auto = noauto] inline = name{, name} [from name{, name}] inline-threshold = unsigned_int inline-limit = unsigned_int list [= file-name short long] noinline noinline = name{, name} [from name{, name}] level = 0 1 2 prof [= file-name] noprof partition = small medium large unsigned_int</pre> <p>where <i>attribute</i> is one of:</p> <ul style="list-style-type: none"> • exits • lowfreq • unknown • safe • isolated • pure

Notes

This option turns on or customizes a class of optimizations known as interprocedural analysis (IPA).

- IPA can significantly increase compilation time, even with the **-qipa=noobject** option, so using IPA should be limited to the final performance tuning stage of development.
- Specify the **-qipa** option on both the compile and link steps of the entire application, or as much of it as possible. You should at least compile the file containing **main**, or at least one of the entry points if compiling a library.
- While IPA's interprocedural optimizations can significantly improve performance of a program, they can also cause previously incorrect but functioning programs to fail. Listed below are some programming practices that can work by accident without aggressive optimization, but are exposed with IPA:
 1. Relying on the allocation order or location of automatic variables. For example, taking the address of an automatic variable and then later comparing it with the address of another local to determine the growth direction of a stack. The C language does not guarantee where an automatic variable is allocated, or its position relative to other automatics. Do not compile such a function with IPA (and expect it to work).
 2. Accessing an either invalid pointer or beyond an array's bounds. IPA can reorganize global data structures. A wayward pointer which may have previously modified unused memory may now trample upon user allocated storage.
- Ensure you have sufficient resources to compile with IPA. IPA can generate significantly larger object files than traditional compilations. As a result, the temporary storage location used to hold these intermediate files (by convention

`/tmp`) is sometimes too small. If a large application is being compiled, consider redirecting temporary storage with the `TMPDIR` environment variable.

- Ensure there is enough virtual memory space to run IPA. For example, 200MB may be enough for small programs, but for large complex applications, several gigabytes of virtual memory may be required for successful IPA linking. Otherwise the operating system might kill IPA with a signal 9, which cannot be trapped, and IPA will be unable to clean up its temporary files.
- You can link objects created with different releases of the compiler, but you must ensure that you use a linker that is at least at the same release level as the newer of the compilers used to create the objects being linked.
- Some symbols which are clearly referenced or set in the source code may be optimized away by IPA, and may be lost to `debug`, `nm`, or `dump` outputs. Using IPA together with the `-g` compiler will usually result in non-steppable output.

Regular expression syntax can be used when specifying a *name* for the following suboptions.

- `exits`
- `inline, noinline`
- `isolated`
- `lowfreq`
- `pure`
- `safe`
- `unknown`

Syntax rules for specifying regular expressions are described below:

Expression	Description
<i>string</i>	Matches any of the characters specified in <i>string</i> . For example, <i>test</i> will match <i>testimony</i> , <i>latest</i> , and <i>intestine</i> .
<i>^string</i>	Matches the pattern specified by <i>string</i> only if it occurs at the beginning of a line.
<i>string\$</i>	Matches the pattern specified by <i>string</i> only if it occurs at the end of a line.
<i>string</i>	The period (<i>.</i>) matches any single character. For example, <i>t.st</i> will match <i>test</i> , <i>tast</i> , <i>tZst</i> , and <i>t1st</i> .
<i>string\special_char</i>	The backslash (<i>\</i>) can be used to escape special characters. For example, assume that you want to find lines ending with a period. Simply specifying the expression <i>.\$</i> would show all lines that had at least one character of any kind in it. Specifying <i>\.\$</i> escapes the period (<i>.</i>), and treats it as an ordinary character for matching purposes.
[<i>string</i>]	Matches any of the characters specified in <i>string</i> . For example, <i>t[a-g123]st</i> matches <i>tast</i> and <i>test</i> , but not <i>t-st</i> or <i>tAst</i> .
[<i>^string</i>]	Does not match any of the characters specified in <i>string</i> . For example, <i>t[^a-zA-Z]st</i> matches <i>t1st</i> , <i>t-st</i> , and <i>t,st</i> but not <i>test</i> or <i>tYst</i> .
<i>string*</i>	Matches zero or more occurrences of the pattern specified by <i>string</i> . For example, <i>te*st</i> will match <i>tst</i> , <i>test</i> , and <i>teeeeeest</i> .
<i>string+</i>	Matches one or more occurrences of the pattern specified by <i>string</i> . For example, <i>t(es)+t</i> matches <i>test</i> , <i>tesest</i> , but not <i>tt</i> .
<i>string?</i>	Matches zero or one occurrences of the pattern specified by <i>string</i> . For example, <i>te?st</i> matches either <i>tst</i> or <i>test</i> .
<i>string{m,n}</i>	Matches between <i>m</i> and <i>n</i> occurrence(s) of the pattern specified by <i>string</i> . For example, <i>a{2}</i> matches <i>aa</i> , and <i>b{1,4}</i> matches <i>b</i> , <i>bb</i> , <i>bbb</i> , and <i>bbbb</i> .
<i>string1 string2</i>	Matches the pattern specified by either <i>string1</i> or <i>string2</i> . For example, <i>s o</i> matches both characters <i>s</i> and <i>o</i> .

The necessary steps to use IPA are:

1. Do preliminary performance analysis and tuning before compiling with the **-qipa** option, because the IPA analysis uses a two-pass mechanism that increases compile and link time. You can reduce some compile and link overhead by using the **-qipa=noobject** option.
2. Specify the **-qipa** option on both the compile and the link steps of the entire application, or as much of it as possible. Use suboptions to indicate assumptions to be made about parts of the program *not* compiled with **-qipa**. During compilation, the compiler stores interprocedural analysis information in the **.o** file. During linking, the **-qipa** option causes a complete recompilation of the entire application.

Note: If a Severe error occurs during compilation, **-qipa** returns RC=1 and terminates. Performance analysis also terminates.

Example

To compile a set of files with interprocedural analysis, enter:

```
xlc -c -O3 *.c -qipa
xlc -o product *.o -qipa
```

Here is how you might compile the same set of files, improving the optimization of the second compilation, and the speed of the first compile step. Assume that there exists two functions, *trace_error* and *debug_dump*, which are rarely executed.

```
xlc -c -O3 *.c -qipa=noobject
xlc -c *.o -qipa=lowfreq=trace_error,debug_dump
```

Related References

“Compiler Command Line Options” on page 51

“inline” on page 147

“libansi” on page 173

“list” on page 175

“pdf1, pdf2” on page 203

“S” on page 225

isolated_call

Purpose

Specifies functions in the source file that have no side effects.

Syntax

```
►► -q-isolated_call== : function_name ◄◄
```

where:

function_name Is the name of a function that does not have side effects, except changing the value of a variable pointed to by a pointer or reference parameter, or does not rely on functions or processes that have side effects.

Side effects are any changes in the state of the runtime environment. Examples of such changes are accessing a volatile object, modifying an external object, modifying a file, or calling another function that does any of these things. Functions with no side effects cause no changes to external and static variables.

function_name can be a list of functions separated by colons (:).

See also “#pragma isolated_call” on page 290 and “#pragma options” on page 299.

Notes

Marking a function as isolated can improve the runtime performance of optimized code by indicating the following to the optimizer:

- external and static variables are not changed by the called function
- calls to the function with loop-invariant parameters may be moved out of loops
- multiple calls to the function with the same parameter may be merged into one call
- calls to the function may be discarded if the result value is not needed

The **#pragma options isolated_call** directive must be specified at the top of the file, before the first C statement. You can use the **#pragma isolated_call** directive at any point in your source file.

Example

To compile myprogram.c, specifying that the functions myfunction(int) and classfunction(double) do not have side effects, enter:

```
xlc myprogram.c -qisolated_call=myfunction:classfunction
```

Related References

“Compiler Command Line Options” on page 51

“#pragma isolated_call” on page 290

“#pragma options” on page 299

keepparam

Purpose

Ensures that function parameters are stored on the stack even if the application is optimized.

Syntax



Notes

A function usually stores its incoming parameters on the stack at the entry point. However, when you compile code with optimization options enabled, the compiler may remove these parameters from the stack if it sees an optimizing advantage in doing so.

Specifying `-qkeepparam` ensures that the parameters are stored on the stack even when optimizing. This compiler option ensures that the values of incoming parameters are available to tools, such as debuggers, by preserving those values on the stack. However, doing so may negatively affect application performance.

Related References

“Compiler Command Line Options” on page 51

keyword

Purpose

This option controls whether the specified name is treated as a keyword or as an identifier whenever it appears in your program source.

Syntax

►► -q keyword
nokeyword = *keyword_name* ►►

Notes

By default all the built-in keywords defined in the C language standard are reserved as keywords. You cannot add keywords to the language with this option. However, you can use **-qnokeyword=keyword_name** to disable certain built-in keywords, and use **-qkeyword=keyword_name** to reinstate those keywords.

This option can be used with the following C built-in keywords:

- asm
- `__complex__` (C99)
- `__imag__` (C99)
- inline
- `__real__` (C99)
- restrict
- typeof

Example

You can reinstate restrict with the following invocation:

```
xlc -qkeyword=restrict
```

Related References

“Compiler Command Line Options” on page 51

L

Purpose

Searches the path directory for library files specified by the *-lkey* option.

Syntax

▶▶ — *-L*—*directory*—————▶▶

Default

The default is to search only the standard directories.

Notes

If the LIBPATH environment variable is set, the compiler will search for libraries first in directory paths specified by LIBPATH, and then in directory paths specified by the *-L* compiler option.

If the *-Ldirectory* option is specified both in the configuration file and on the command line, search paths specified in the configuration file are the first to be searched.

Example

To compile *myprogram.c* so that the directory */usr/tmp/old* and all other directories specified by the *-l* option are searched for the library *libspfiles.a*, enter:

```
xlc myprogram.c -lspfiles -L/usr/tmp/old
```

Related References

“Compiler Command Line Options” on page 51

“l” on page 164

Purpose

Searches the specified library file, *libkey.so*, and then *libkey.a* for dynamic linking, or just *libkey.a* for static linking.

Syntax

▶— *-lkey* —▶

Default

The compiler default is to search only some of the compiler run-time libraries. The default configuration file specifies the default library names to search for with the **-l** compiler option, and the default search path for libraries with the **-L** compiler option.

Notes

You must also provide additional search path information for libraries not located in the default search path. The search path can be modified with the **-Ldirectory** or **-Z** options. See **-B**, **-brtl**, and **-bstatic,-bdynamic** for information on specifying the types of libraries that are searched (for static or dynamic linking).

The C runtime libraries are automatically added.

The **-l** option is cumulative. Subsequent appearances of the **-l** option on the command line do not replace, but add to, the list of libraries specified by earlier occurrences of **-l**. Libraries are searched in the order in which they appear on the command line, so the order in which you specify libraries can affect symbol resolution in your application.

For more information, refer to the **ld** documentation for your operating system.

Example

To compile *myprogram.c* and link it with library *mylibrary* (*libmylibrary.a*) found in the */usr/mylibdir* directory, enter:

```
xlc myprogram.c -lmylibrary -L/usr/mylibdir
```

Related Tasks

“Specify Compiler Options in a Configuration File” on page 32

Related References

“Compiler Command Line Options” on page 51

“B” on page 78

“b” on page 79

“brtl” on page 82

“L” on page 163

“Z” on page 270

langlvl

Purpose

Selects the language level and language options for the compilation.

Syntax



where values for *suboption* are described below in the **Notes** section.

See also “#pragma langlvl” on page 292 and “#pragma options” on page 299.

Default

The default language level varies according to the command you use to invoke the compiler:

Invocation	Default language level
<code>xlC</code>	<code>extc89</code>
<code>cc</code>	<code>extended</code>
<code>c89</code>	<code>stdc89</code>
<code>c99</code>	<code>stdc99</code>

Notes

You can also use either of the following pragma directives to specify the language level in your C language source program:

```
#pragma options langlvl=suboption
#pragma langlvl(suboption)
```

The **pragma** directive must appear before any noncommentary lines in the source code.

For C programs, you can use the following **-qlanglvl** suboptions for *suboption*:

<code>classic</code>	Allows the compilation of non-stdc89 programs, and conforms closely to the K&R level preprocessor. This language level is not supported by the AIX v5.1 system header files, such as <code>math.h</code> . If you must use the AIX v5.1 system header files, consider compiling your program to the stdc89 or extended language levels.
<code>extended</code>	Provides compatibility with the RT compiler and classic . This language level is based on C89.
<code>saa</code>	Compilation conforms to the current SAA [®] C CPI language definition. This is currently SAA C Level 2.
<code>saal2</code>	Compilation conforms to the SAA C Level 2 CPI language definition, with some exceptions.
<code>stdc89</code>	Compilation conforms to the ANSI C89 standard, also known as ISO C90.
<code>stdc99</code>	Compilation conforms to the ISO C99 standard. Note: Not all operating system releases support the header files and runtime library required by C99.
<code>extc89</code>	Compilation conforms to the ANSI C89 standard, and accepts implementation-specific language extensions.

extc99	<p>Compilation conforms to the ISO C99 standard, and accepts implementation-specific language extensions.</p> <p>Note: Not all operating system releases support the header files and runtime library required by C99.</p>
[no]c99complex	<p>This suboption instructs the compiler to recognize C99 complex data types and related keywords.</p> <p>Note: Support for complex data types may vary among different compilers, creating potential portability issues. The compiler will issue a portability warning message if you specify this compiler option together with -qinfo=por.</p>
[no]c99complexheader	<p>This suboption instructs the compiler to use the C99 complex.h header file.</p>
[no]compatzea	<p>This option can be used for all C language levels. The default is -qlanglvl=nocompatzea. Specifying -qlanglvl=compatzea has effect only if -qlanglvl=zeroextarray is also in effect.</p> <p>Prior to XL C v7.0 for AIX, zero extent arrays on AIX had an underlying dimension of 1. Applying the sizeof operator to a zero extent array would return 1 times the size of its element type.</p> <p>Starting with XL C v7.0 for AIX, the underlying dimension of zero extent arrays will change to support an underlying dimension of 0. This behavior is consistent with the behavior found in the Linux and Mac OS X versions of the XL C/C++ compiler.</p> <p>This change can influence the behavior of existing code. The -qlanglvl=compatzea option is a compatibility option that lets you continue to use a zero extent array with dimension of 1 in AIX applications compiled with XL C v7.0 for AIX.</p>
[no]ucs	<p>Under language levels stdc99 and extc99, the default is -qlanglvl=ucs.</p> <p>This option controls whether Unicode characters are allowed in identifiers, string literals and character literals in program source code.</p> <p>The Unicode character set is supported by the C standard. This character set contains the full set of letters, digits and other characters used by a wide range of languages, including all North American and Western European languages. Unicode characters can be 16 or 32 bits. The ASCII one-byte characters are a subset of the Unicode character set.</p> <p>When this option is set to yes, you can insert Unicode characters in your source files either directly or using a notation that is similar to escape sequences. Because many Unicode characters cannot be displayed on the screen or entered from the keyboard, the latter approach is usually preferred. Notation forms for Unicode characters are <code>\uhhhh</code> for 16-bit characters, or <code>\Uhhhhhhhhh</code> for 32-bit characters, where <i>h</i> represents a hexadecimal digit. Short identifiers of characters are specified by ISO/IEC 10646.</p>

[no]zeroextarray This option enables support for flexible array members and zero extent arrays. The default is **-qlanglvl=nozeroextarray**.

Though this suboption is supported for all C language levels, the compiler will issue warning messages when compiling with flexible array members to a language level of stdc89. The compiler will also issue severe error messages when compiling with zero extent arrays to a language level of stdc89 or stdc99. You can suppress these messages by preceding the definition for the flexible array member or zero extent array with the **__extension__** keyword.

The following **-qlanglvl** suboptions are accepted but ignored by the C compiler. Use **-qlanglvl=extended**, **-qlanglvl=extc99**, or **-qlanglvl=extc89** to enable the functions that these suboptions imply. For other values of **-qlanglvl**, the functions implied by these suboptions are disabled.

[no]gnu_assert	GNU C portability option.
[no]gnu_explicitregvar	GNU C portability option.
[no]gnu_include_next	GNU C portability option.
[no]gnu_locallabel	GNU C portability option.
[no]gnu_warning	GNU C portability option.

Exceptions to the **stdc89** mode addressed by **classic** are as follows:

Tokenization Tokens introduced by macro expansion may be combined with adjacent tokens in some cases. Historically, this was an artifact of the text-based implementations of older preprocessors, and because, in older implementations, the preprocessor was a separate program whose output was passed on to the compiler.

For similar reasons, tokens separated only by a comment may also be combined to form a single token. Here is a summary of how tokenization of a program compiled in **classic** mode is performed:

1. At a given point in the source file, the next token is the longest sequence of characters that can possibly form a token. For example, `i++++j` is tokenized as `i ++ ++ + j` even though `i ++ + ++ j` may have resulted in a correct program.
2. If the token formed is an identifier and a macro name, the macro is replaced by the text of the tokens specified on its **#define** directive. Each parameter is replaced by the text of the corresponding argument. Comments are removed from both the arguments and the macro text.
3. Scanning is resumed at the first step from the point at which the macro was replaced, as if it were part of the original program.
4. When the entire program has been preprocessed, the result is scanned again by the compiler as in the first step. The second and third steps do not apply here since there will be no macros to replace. Constructs generated by the first three steps that resemble preprocessing directives are not processed as such.

It is in the third and fourth steps that the text of adjacent but previously separate tokens may be combined to form new tokens.

The `\` character for line continuation is accepted only in string and character literals and on preprocessing directives.

Constructs such as:

```
#if 0
  "unterminated
#endif
#define US "Unterminating string
char *s = US terminated now"
```

will not generate diagnostic messages, since the first is an unterminated literal in a **FALSE** block, and the second is completed after macro expansion. However:

```
char *s = US;
```

will generate a diagnostic message since the string literal in **US** is not completed before the end of the line.

Empty character literals are allowed. The value of the literal is zero.

Preprocessing directives

The # token must appear in the first column of the line. The token immediately following # is available for macro expansion. The line can be continued with \ only if the name of the directive and, in the following example, the (has been seen:

```
#define f(a,b) a+b
f\
(1,2)      /* accepted */
#define f(a,b) a+b
f\
1,2)      /* not accepted */
```

The rules concerning \ apply whether or not the directive is valid. For example,

```
#\
define M 1  /* not allowed */
#def\
ine M 1     /* not allowed */
#define\
M 1        /* allowed */
#dfine\
M 1        /* equivalent to #dfine M 1, even
            though #dfine is not valid */
```

Following are the preprocessor directive differences between **classic** mode and **stdc89** mode. Directives not listed here behave similarly in both modes.

#ifndef/#ifndef

When the first token is not an identifier, no diagnostic message is generated, and the condition is FALSE.

#else When there are extra tokens, no diagnostic message is generated.

#endif When there are extra tokens, no diagnostic message is generated.

#include

The < and > are separate tokens. The header is formed by combining the spelling of the < and > with the tokens between them. Therefore /* and // are recognized as comments (and are always stripped), and the " and ' do begin literals within the < and >. (Remember that in C programs, C++-style comments // are recognized when **-qcpluscmt** is specified.)

#line The spelling of all tokens which are not part of the line number form the new file name. These tokens need not be string literals.

#error Not recognized in **classic** mode.

#define

A valid macro parameter list consists of zero or more identifiers each separated by commas. The commas are ignored and the parameter list is constructed as if they were not specified. The parameter names need not be unique. If there is a conflict, the last name specified is recognized.

For an invalid parameter list, a warning is issued. If a macro name is redefined with a new definition, a warning will be issued and the new definition used.

#undef When there are extra tokens, no diagnostic message is generated.

Macro expansion

- When the number of arguments on a macro invocation does not match the number of parameters, a warning is issued.
- If the (token is present after the macro name of a function-like macro, it is treated as too few arguments (as above) and a warning is issued.
- Parameters are replaced in string literals and character literals.
- Examples:

```
#define M()    1
#define N(a)  (a)
#define O(a,b) ((a) + (b))

M(); /* no error */
N(); /* empty argument */
O(); /* empty first argument
      and too few arguments */
```

Text Output No text is generated to replace comments.

Related References

“Compiler Command Line Options” on page 51
“bitfields” on page 80
“chars” on page 88
“digraph” on page 101
“flag” on page 117
“info” on page 142
“inline” on page 147
“M” on page 180
“ro” on page 221
“suppress” on page 243
“#pragma langlvl” on page 292
“#pragma options” on page 299

See also the *IBM C Language Extensions* section of the *C/C++ Language Reference*.

largepage

Purpose

Instructs the compiler to exploit large page heaps available on POWER4 and POWER5 systems running AIX v5.1D or later.

Syntax

►► -q no largepage
largepage ◀◀

Notes

Compiling with **-qlargepage** can result in improved program performance. This option has effect only on POWER4 and POWER5 systems running AIX v5.1D or later.

This option is only valid when used together with IPA (**-qipa**, **-O4**, **-O5** compiler options).

Example

To compile myprogram.c to use large page heaps, enter:

```
xlc myprogram.c -qlargepage
```

Related References

“Compiler Command Line Options” on page 51

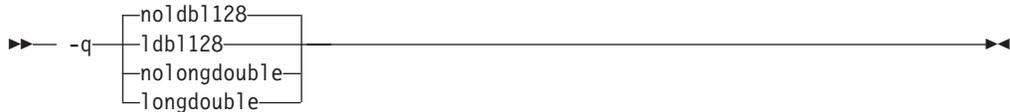
“ipa” on page 151

ldbl128, longdouble

Purpose

Increases the size of **long double** type from 64 bits to 128 bits.

Syntax



See also “#pragma options” on page 299.

Notes

The `-qlongdouble` option is the same as the `-qldbl128` option.

Separate libraries are provided that support 128-bit **long double** types. These libraries will be automatically linked if you use any of the invocation commands with the **128** suffix (`xlc128`, `cc128`, `xlc128_r`, or `cc128_r`). You can also manually link to the 128-bit versions of the libraries using the `-lkey` option, as shown in the following table:

Default (64-bit) long double		128-bit long double	
Library	Form of the <code>-lkey</code> option	Library	Form of the <code>-lkey</code> option
libC.a	-lC	libC128.a	-lC128
libC_r.a	-lC_r	libC128_r.a	-lC128_r

Linking without the 128-bit versions of the libraries when your program uses 128-bit **long doubles** (for example, if you specify `-qldbl128` alone) may produce unpredictable results.

The `-qldbl128` option defines `__LONGDOUBLE128`.

The `#pragma options` directive must appear before the first C statement in the source file, and the option applies to the entire file.

Example

To compile `myprogram.c` so that **long double** types are 128 bits, enter:

```
xlc myprogram.c -qldbl128 -lC128
```

or:

```
xlc128 myprogram.c
```

Related References

“Compiler Command Line Options” on page 51

“l” on page 164

“#pragma options” on page 299

libansi

Purpose

Assumes that all functions with the name of an ANSI C library function are in fact the system functions.

Syntax

►► -q no libansi libansi ◀◀

See also “#pragma options” on page 299.

Notes

This will allow the optimizer to generate better code because it will know about the behavior of a given function, such as whether or not it has any side effects.

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

linedebug

Purpose

Generates line number and source file name information for the debugger.

Syntax



Notes

This option produces minimal debugging information, so the resulting object size is smaller than that produced if the **-g** debugging option is specified. You can use the debugger to step through the source code, but you will not be able to see or query variable information. The traceback table, if generated, will include line numbers.

Avoid using this option with **-O** (optimization) option. The information produced may be incomplete or misleading.

If you specify the **-qlinedebug** option, the inlining option defaults to **-Q!** (no functions are inlined).

The **-g** option overrides the **-qlinedebug** option. If you specify **-g -qno linedebug** on the command line, **-qno linedebug** is ignored and the following warning is issued:

```
1506-... (W) Option -qno linedebug is incompatible with option -g and is ignored
```

Example

To compile `myprogram.c` to produce an executable program `testing` so you can step through it with a debugger, enter:

```
xlc myprogram.c -o testing -qlinedebug
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

“O, optimize” on page 194

“Q” on page 214

“#pragma options” on page 299

list

Purpose

Produces a compiler listing that includes an object listing.

Syntax

►► -q no list
list ◄◄

See also “#pragma options” on page 299.

Notes

The `-qnoprint` compiler option overrides this option.

Example

To compile `myprogram.c` and produce an object listing, enter:

```
xlc myprogram.c -qlist
```

Related References

“Compiler Command Line Options” on page 51

“print” on page 210

“#pragma options” on page 299

listopt

Purpose

Produces a compiler listing that displays all options in effect at time of compiler invocation.

Syntax

►► -q no listopt
listopt ◀◀

Notes

The listing will show options in effect as set by the compiler defaults, default configuration file, and command line settings. Option settings caused by **#pragma** statements in the program source are not shown in the compiler listing.

Specifying **-qnoprint** overrides this compiler option.

Example

To compile myprogram.c to produce a compiler listing that shows all options in effect, enter:

```
xlc myprogram.c -qlistopt
```

Related References

“Compiler Command Line Options” on page 51

“print” on page 210

“Resolving Conflicting Compiler Options” on page 35

longlit

Purpose

Makes unaffixed literals into the long type in 64-bit mode.

Syntax



Notes

The following table shows the implicit types for constants in 64-bit mode when compiling in the **stdc89**, **extc89**, or **extended** language level:

	default 64-bit mode	64-bit mode with qlonglit
unaffixed decimal	signed int signed long unsigned long	signed long unsigned long
unaffixed octal or hex	signed int unsigned int signed long unsigned long	signed long unsigned long
affixed by u/U	unsigned int unsigned long	unsigned long
affixed by l/L	signed long unsigned long	signed long unsigned long
affixed by ul/UL	unsigned long	unsigned long

The following table shows the implicit types for constants in 64-bit mode when compiling in the **stdc99**, **extc99**, or **extended** language level:

	Decimal Constant	-qlonglit effect on Decimal Constant
unaffixed	int long int	long int
u or U	unsigned int unsigned long int	unsigned long int
l or L	long int	long int
Both u or U, and l or L	unsigned long int	unsigned long int
ll or LL	long long int	long long int
Both u or U, and ll or LL	unsigned long long int	unsigned long long int

	Octal or Hexadecimal Constant	-qlonglit effect on Octal or Hexadecimal Constant
unaffixed	int unsigned int long int unsigned long int	long int unsigned long int
u or U	unsigned int unsigned long int	unsigned long int

	Octal or Hexadecimal Constant	-qlonglit effect on Octal or Hexadecimal Constant
l or L	long int unsigned long int	long int unsigned long int
Both u or U, and l or L	unsigned long int	unsigned long int
ll or LL	long long int unsigned long long int	long long int unsigned long long int
Both u or U, and ll or LL	unsigned long long int	unsigned long long int

Related References

“Compiler Command Line Options” on page 51

“langlvl” on page 165

longlong

Purpose

Allows **long long** integer types in your program.

Syntax

►► -q longlong
no longlong ◀◀

Default

The default with **xlc** and **cc** is **-qlonglong**, which defines **_LONG_LONG** (**long long** types will work in programs). The default with **c89** is **-qnolonglong** (**long long** types are not supported).

Notes

This option cannot be specified when the selected language level is **stdc99** or **extc99**. It is used to control the long long support that is provided as an extension to the C89 standard. This extension is slightly different from the long long support that is part of the C99 standard.

Examples

1. To compile `myprogram.c` so that **long long ints** are not allowed, enter:

```
xlc myprogram.c -qnolonglong
```

2. AIX v4.2 and later provides support for files greater than 2 gigabytes in size so you can store large quantities of data in a single file. To allow Large File manipulation in your application, compile with the **-D_LARGE_FILES** and **-qlonglong** compiler options. For example:

```
xlc myprogram.c -D_LARGE_FILES -qlonglong
```

Related References

“Compiler Command Line Options” on page 51

M

Purpose

Creates an output file that contains targets suitable for inclusion in a description file for the **make** command.

Syntax

►— -M —►

Notes

The **-M** option is functionally identical to the **-qmakedep** option.

.u files are not **make** files; **.u** files must be edited before they can be used with the **make** command. For more information on this command, see your operating system documentation.

The output file contains a line for the input file and an entry for each include file. It has the general form:

```
file_name.o:file_name.c
file_name.o:include_file_name
```

Include files are listed according to the search order rules for the **#include** preprocessor directive, described in *Directory Search Sequence for Include Files Using Relative Path Names*. If the include file is not found, it is not added to the **.u** file.

Files with no include statements produce output files containing one line that lists only the input file name.

Examples

If you do not specify the **-o** option, the output file generated by the **-M** option is created in the current directory. It has a **.u** suffix. For example, the command:

```
xlc -M person_years.c
```

produces the output file **person_years.u**.

A **.u** file is created for every input file with a **.c** or **.i** suffix. Otherwise, output **.u** files are not created for any other files.

For example, the command:

```
xlc -M conversion.c filter.c /lib/libm.a
```

produces two output files, **conversion.u** and **filter.u**, and an executable file as well. No **.u** file is created for the library.

If the current directory is not writable, no **.u** file is created. If you specify **-o file_name** along with **-M**, the **.u** file is placed in the directory implied by **-o file_name**. For example, for the following invocation:

```
xlc -M -c t.c -o /tmp/t.o
```

places the **.u** output file in **/tmp/t.u**.

Related Tasks

“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

Related References

“Compiler Command Line Options” on page 51

“makedep” on page 186

“o” on page 198

“sourcetype” on page 235

ma

Purpose

Substitutes inline code for calls to built-in function **alloca**.

Syntax

▶▶ — `-ma` —▶▶

Notes

If `#pragma alloca` is unspecified, or if you do not use `-ma`, **alloca** is treated as a user-defined identifier rather than as a built-in function.

Example

To compile `myprogram.c` so that calls to the function **alloca** are treated as inline, enter:

```
xlc myprogram.c -ma
```

Related References

“Compiler Command Line Options” on page 51

“`alloca`” on page 68

“`#pragma alloca`” on page 274

macpstr

Purpose

Converts Pascal string literals into null-terminated strings where the first byte contains the length of the string.

Syntax

► — -q — nomacpstr
macpstr —►

See also “#pragma options” on page 299.

Notes

A Pascal string literal always contains the characters “\p”. The characters \p in the middle of a string do not form a Pascal string literal; the characters must be *immediately preceded* by the “ (double quote) character.

The final length of the Pascal string literal can be no longer than 255 bytes (the maximum length that can fit in a byte).

For example, when the **-qmacpstr** option is in effect, the compiler converts:

```
"\pABC"
```

to:

```
'\03' , 'A' , 'B' , 'C' , '\0'
```

The compiler ignores the **-qmacpstr** option when the **-qmbcs** or **-qdbcs** option is active because Pascal-string-literal processing is only valid for one-byte characters.

The **#pragma options** keyword **macpstr** is only valid at the top of a source file before any C source statements. If you attempt to use it in the middle of a source file, it is ignored and the compiler issues an error message.

String Literal Processing: The following describes how Pascal string literals are processed.

- Concatenating a Pascal string literal to a normal string gives a non-Pascal string. For example, concatenating the strings:

```
"ABC" "\pDEF"
```

gives:

```
"ABCpDEF"
```

- A Pascal string literal cannot be concatenated with a **wide** string literal.
- The compiler truncates a Pascal string literal that is longer than 255 bytes (excluding the length byte and the terminating NULL) to 255 characters.
- The compiler ignores the **-qmacpstr** option if **-qmbcs** or **-qdbcs** is used, and issues a warning message.
- Because there is no Pascal-string-literal processing of wide strings, using the escape sequence \p in a wide string literal with the **-qmacpstr** option, generates a warning message and the escape sequence is ignored.
- The Pascal string literal is *not* a basic type different from other C string literals. After the processing of the Pascal string literal is complete, the resulting string is

treated the same as all other strings. If the program passes a C string to a function that expects a Pascal string, or vice versa, the behavior is undefined.

- Concatenating two Pascal string literals, for example, `strcat()`, does not result in a Pascal string literal. However, as described above, two adjacent Pascal string literals can be concatenated to form one Pascal string literal in which the first byte is the length of the new string literal. For example, concatenating the strings:

```
"\p ABC" "\p DEF"
```

or

```
"\p ABC" "DEF"
```

results in:

```
"\06ABCDEF"
```

- Modifying any byte of the Pascal string literal after the processing has been completed does not alter the original length value in the first byte. For example, in the string `"\06ABCDEF"`, substituting a null character for one of the existing characters in the middle of the string does not change the value of the first byte of the string, which contains the length of the string.
- No errors or warnings are issued when the bytes of the processed Pascal string literal are modified.
- Entering the characters:

```
'\p' , 'A' , 'B' , 'C' , '\0'
```

into a character array does not form a Pascal string literal.

Example

To compile `mypascal.c` and convert string literals into null-terminated strings, enter:

```
xlc mypascal.c -qmacpstr
```

Related References

“Compiler Command Line Options” on page 51

“mbs, dbs” on page 191

“#pragma options” on page 299

maf

Purpose

Specifies whether floating-point multiply-add instructions are to be generated. This option affects the precision of floating-point intermediate results.

Syntax

►► — -q —  —►►

See also “#pragma options” on page 299.

Notes

This option is obsolete. Use -qfloat=maf in your new applications.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“#pragma options” on page 299

makedep

Purpose

Creates an output file that contains targets suitable for inclusion in a description file for the **make** command.

Syntax

►— -q—makedep—◄

Notes

The **-qmakedep** option is functionally identical to the **-M** option.

.u files are not **make** files; **.u** files must be edited before they can be used with the **make** command. For more information on this command, see your operating system documentation.

If you do not specify the **-o** option, the output file generated by the **-qmakedep** option is created in the current directory. It has a **.u** suffix. For example, the command:

```
xlc -qmakedep person_years.c
```

produces the output file **person_years.u**.

A **.u** file is created for every input file with a **.c** or **.i** suffix. Otherwise, output **.u** files are not created for any other files.

For example, the command:

```
xlc -qmakedep conversion.c filter.c /lib/libm.a
```

produces two output files, **conversion.u** and **filter.u** (and an executable file as well). No **.u** file is created for the library.

If the current directory is not writable, no **.u** file is created. If you specify **-o file_name** along with **-qmakedep**, the **.u** file is placed in the directory implied by **-ofile_name**. For example, for the following invocation:

```
xlc -qmakedep -c t.c -o /tmp/t.o
```

places the **.u** output file in **/tmp/t.u**.

The output file contains a line for the input file and an entry for each include file. It has the general form:

```
file_name.o:include_file_name  
file_name.o:file_name.c
```

Include files are listed according to the search order rules for the **#include** preprocessor directive, described in “Directory Search Sequence for Include Files Using Relative Path Names” on page 37. If the include file is not found, it is not added to the **.u** file.

Files with no include statements produce output files containing one line that lists only the input file name.

Related References

“Compiler Command Line Options” on page 51

“M” on page 180

“o” on page 198

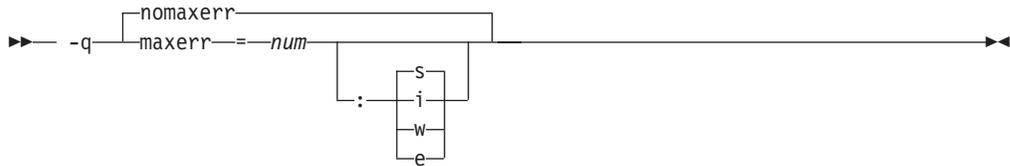
“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

maxerr

Purpose

Instructs the compiler to halt compilation when *num* errors of a specified severity level or higher is reached.

Syntax



where *num* must be an integer. Choices for severity level can be one of the following:

<i>sev_level</i>	Description
i	Informational
w	Warning
e	Error
s	Severe error

Notes

If a severity level is not specified, the current value of the **-qhalt** option is used.

If the **-qmaxerr** option is specified more than once, the **-qmaxerr** option specified last determines the action of the option. If both the **-qmaxerr** and **-qhalt** options are specified, the **-qmaxerr** or **-qhalt** option specified last determines the severity level used by the **-qmaxerr** option.

An unrecoverable error occurs when the number of errors reached the limit specified. The error message issued is similar to:

```
1506-672 (U) The number of errors has reached the limit of ...
```

If **-qnomaxerr** is specified, the entire source file is compiled regardless of how many errors are encountered.

Diagnostic messages may be controlled by the **-qflag** option.

Examples

1. To stop compilation of `myprogram.c` when 10 warnings are encountered, enter the command:

```
xlc myprogram.c -qmaxerr=10:w
```
2. To stop compilation of `myprogram.c` when 5 severe errors are encountered, assuming that the current **-qhalt** option value is **S** (severe), enter the command:

```
xlc myprogram.c -qmaxerr=5
```
3. To stop compilation of `myprogram.c` when 3 informational messages are encountered, enter the command:

```
xlc myprogram.c -qmaxerr=3:i
```

or:

```
xlc myprogram.c -qmaxerr=3 -qhalt=i
```

Related References

“Compiler Command Line Options” on page 51

“flag” on page 117

“halt” on page 132

“Message Severity Levels and Compiler Response” on page 355

maxmem

Purpose

Limits the amount of memory used by the optimizer for local tables of specific, memory-intensive optimizations. The memory size limit is specified in kilobytes.

Syntax

► `-qmaxmem=size` ◄

Defaults

- With `-O2` optimization in effect, `maxmem=8192`.
- With `-O3` or greater optimization in effect, `maxmem=-1`.

Notes

- A *size* value of `-1` permits each optimization to take as much memory as it needs without checking for limits. Depending on the source file being compiled, the size of subprograms in the source, the machine configuration, and the workload on the system, this might exceed available system resources.
- The limit set by `-qmaxmem` is the amount of memory for specific optimizations, and not for the compiler as a whole. Tables required during the entire compilation process are not affected by or included in this limit.
- Setting a large limit has no negative effect on the compilation of source files when the compiler needs less memory.
- Limiting the scope of optimization does not necessarily mean that the resulting program will be slower, only that the compiler may finish before finding all opportunities to increase performance.
- Increasing the limit does not necessarily mean that the resulting program will be faster, only that the compiler is better able to find opportunities to increase performance if they exist.

Depending on the source file being compiled, the size of the subprograms in the source, the machine configuration, and the workload on the system, setting the limit too high might lead to page-space exhaustion. In particular, specifying `-qmaxmem=-1` allows the compiler to try and use an infinite amount of storage, which in the worst case can exhaust the resources of even the most well-equipped machine.

Example

To compile `myprogram.c` so that the memory specified for local table is **16384** kilobytes, enter:

```
xlc myprogram.c -qmaxmem=16384
```

Related References

“Compiler Command Line Options” on page 51

mbcs, dbcs

Purpose

Use the **-qmbcs** option if your program contains multibyte characters. The **-qmbcs** option is equivalent to **-qdbcs**.

Syntax



See also “#pragma options” on page 299.

Notes

Multibyte characters are used in certain languages such as Chinese, Japanese, and Korean.

Multibyte characters are also permitted in comments, if you specify the **-qmbcs** or **-qdbcs** compiler option.

If a source file contains multibyte character literals and the default **-qnombs** or **-qnodbcs** compiler option is in effect, the compiler will treat all literals as single-byte literals.

Example

To compile myprogram.c if it contains multibyte characters, enter:

```
xlc myprogram.c -qmbcs
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

Appendix C, “National Languages Support in XL C Enterprise Edition,” on page 385

mkshrojb

Purpose

Creates a shared object from generated object files.

Syntax

►— -q—mkshrojb—►

Notes

This option, together with the related options described below, should be used instead of the **makeC++SharedLib** command to create a shared object.

The following table shows equivalent options between **makeC++SharedLib** and **-qmkshrojb**:

makeC++SharedLib options	-qmkshrojb and related options
-p <i>nnn</i>	-qmkshrojb= <i>nnn</i>
-e <i>file_name</i>	-qexpfile= <i>file_name</i>
-E <i>export_file</i>	-bE: <i>export_file</i>
-I <i>import_file</i>	-bI: <i>import_file</i>
-x	-qnolib
-x 32	-q32
-x 64	-q64
-n <i>entry_point</i>	-e <i>entry_point</i>

The compiler will automatically export all global symbols from the shared object unless one explicitly specifies which symbols to export with the **-bE**:, **-bexport**: or **-bexpall** options.

Specifying **-qmkshrojb** implies **-qpic**.

Also, the following related options can be used with the **-qmkshrojb** compiler option:

-o <i>shared_file</i>	Is the name of the file that will hold the shared file information. The default is <code>shr.o</code> .
-qexpfile = <i>filename</i>	Saves all exported symbols in <i>filename</i> . This option is ignored unless x1C automatically creates the export list.
-e <i>name</i>	Sets the entry name for the shared executable to <i>name</i> . The default is <code>-eentry</code> .

If you use **-qmkshrojb** to create a shared library, the compiler will:

1. If you don't specify **-bE**:, **-bexport**:, **-bexpall** or **-bnoexpall**, create an export list containing all global symbols using the `CreateExportList` command, described below. You can specify another script with the **-tE**/**-B** or **-qpath=E**: options.
2. Save the export list if either the `CreateExportList` utility or the **-qmkshrojb** and **-qexpfile** compiler options were used to create an export list.
3. Call the linker with the appropriate options and object files to build a shared object.

For more information about the CreateExportList utility and creating shared objects, see *Creating a Library* in the Programming Guide.

Related References

“Compiler Command Line Options” on page 51

“32, 64” on page 60

“b” on page 79

“e” on page 107

“expfile” on page 112

“o” on page 198

“path” on page 202

“pic” on page 208

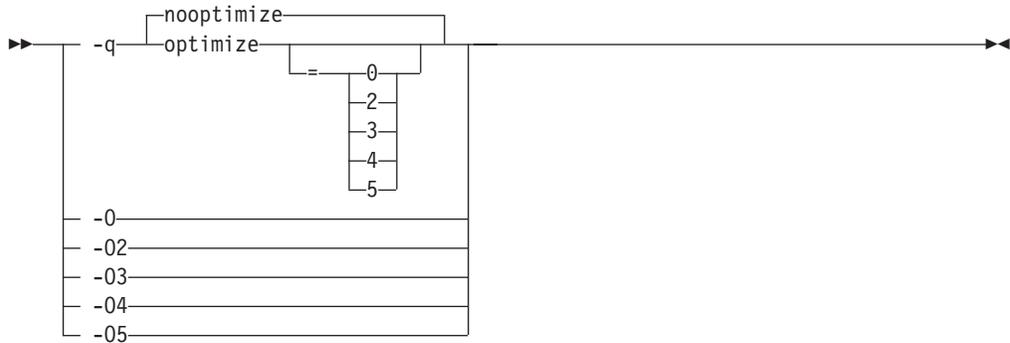
See also the *Creating a Library* section of the *XL C Programming Guide*.

O, optimize

Purpose

Optimizes code at a choice of levels during compilation.

Syntax



where optimization settings are:

<code>-qNOOPTimize</code> <code>-qOPTimize=0</code>	<p>Performs only quick local optimizations such as constant folding and elimination of local common subexpressions.</p> <p>This setting implies <code>-qstrict_induction</code> unless <code>-qnostrict_induction</code> is explicitly specified.</p>
<code>-O</code> <code>-qOPTimize</code>	<p>Performs optimizations that the compiler developers considered the best combination for compilation speed and runtime performance. The optimizations may change from product release to release. If you need a specific level of optimization, specify the appropriate numeric value.</p> <p>This setting implies <code>-qstrict</code> and <code>-qstrict_induction</code>, unless explicitly negated by <code>-qnostrict_induction</code> or <code>-qnostrict</code>.</p>
<code>-O2</code> <code>-qOPTimize=2</code>	<p>Same as <code>-O</code>.</p>
<code>-O3</code> <code>-qOPTimize=3</code>	<p>Performs additional optimizations that are memory intensive, compile-time intensive, or both. They are recommended when the desire for runtime improvement outweighs the concern for minimizing compilation resources.</p> <p><code>-O3</code> applies the <code>-O2</code> level of optimization, but with unbounded time and memory limits. <code>-O3</code> also performs higher and more aggressive optimizations that have the potential to slightly alter the semantics of your program. The compiler guards against these optimizations at <code>-O2</code>.</p> <p>Use the <code>-qstrict</code> option with <code>-O3</code> to turn off the aggressive optimizations that might change the semantics of a program. Specifying <code>-qstrict</code> together with <code>-O3</code> invokes all the optimizations performed at <code>-O2</code> as well as further loop optimizations. The <code>-qstrict</code> compiler option must appear after the <code>-O3</code> option, otherwise it is ignored.</p>

-O3
-qOPTimize=3
(continued)

The aggressive optimizations performed when you specify **-O3** are:

1. Aggressive code motion, and scheduling on computations that have the potential to raise an exception, are allowed.

Loads and floating-point computations fall into this category. This optimization is aggressive because it may place such instructions onto execution paths where they *will* be executed when they *may* not have been according to the actual semantics of the program.

For example, a loop-invariant floating-point computation that is found on some, but not all, paths through a loop will not be moved at **-O2** because the computation may cause an exception. At **-O3**, the compiler will move it because it is not certain to cause an exception. The same is true for motion of loads. Although a load through a pointer is never moved, loads off the static or stack base register are considered movable at **-O3**. Loads in general are not considered to be absolutely safe at **-O2** because a program can contain a declaration of a static array `a` of 10 elements and load `a[60000000003]`, which could cause a segmentation violation.

The same concepts apply to scheduling.

Example:

In the following example, at **-O2**, the computation of `b+c` is not moved out of the loop for two reasons:

- It is considered dangerous because it is a floating-point operation
- `t` does not occur on every path through the loop

At **-O3**, the code is moved.

```
...
int i ;
float a[100], b, c ;
for (i = 0 ; i < 100 ; i++)
{
    if (a[i] < a[i+1])
        a[i] = b + c ;
}
...
```

2. Conformance to IEEE rules are relaxed.

With **-O2** certain optimizations are not performed because they may produce an incorrect sign in cases with a zero result, and because they remove an arithmetic operation that may cause some type of floating-point exception.

For example, `X + 0.0` is not folded to `X` because, under IEEE rules, `-0.0 + 0.0 = 0.0`, which is `-X`. In some other cases, some optimizations may perform optimizations that yield a zero result with the wrong sign. For example, `X - Y * Z` may result in a `-0.0` where the original computation would produce `0.0`.

In most cases the difference in the results is not important to an application and **-O3** allows these optimizations.

3. Floating-point expressions may be rewritten.

Computations such as `a*b*c` may be rewritten as `a*c*b` if, for example, an opportunity exists to get a common subexpression by such rearrangement. Replacing a divide with a multiply by the reciprocal is another example of reassociating floating-point computations.

<p><code>-O3</code>, <code>-qOPTimize=3</code> (continued)</p>	<p>Notes</p> <ul style="list-style-type: none"> • <code>-qfloat=fltint:rsqrt</code> is set by default with <code>-O3</code>. • <code>-qmaxmem=1</code> is set by default with <code>-O3</code>, allowing the compiler to use as much memory as necessary when performing optimizations. • Built-in functions do not change <code>errno</code> at <code>-O3</code>. • Aggressive optimizations do <i>not</i> include the following floating-point suboptions: <code>-qfloat=hsflt</code>, <code>hssngl</code>, and <code>-qfloat=rndsngl</code>, or anything else that affects the precision mode of a program. • Integer divide instructions are considered too dangerous to optimize even at <code>-O3</code>. • The default <code>-qmaxmem</code> value is <code>-1</code> at <code>-O3</code>. • Refer to <code>-qflttrap</code> to see the behavior of the compiler when you specify <code>optimize</code> options with the <code>flttrap</code> option. • You can use the <code>-qstrict</code> and <code>-qstrict_induction</code> compiler options to turn off effects of <code>-O3</code> that might change the semantics of a program. To have effect, the <code>-qstrict</code> compiler option must appear after the <code>-O3</code> option on the command line. • The <code>-O3</code> compiler option followed by the <code>-O</code> option leaves <code>-qignerrno</code> on.
<p><code>-O4</code> <code>-qOPTimize=4</code></p>	<p>This option is the same as <code>-O3</code>, except that it also:</p> <ul style="list-style-type: none"> • Sets the <code>-qarch</code> and <code>-qtune</code> options to the architecture of the compiling machine • Sets the <code>-qcache</code> option most appropriate to the characteristics of the compiling machine • Sets the <code>-qhot</code> option • Sets the <code>-qipa</code> option <p>Note: Later settings of <code>-O</code>, <code>-qcache</code>, <code>-qhot</code>, <code>-qipa</code>, <code>-qarch</code>, and <code>-qtune</code> options will override the settings implied by the <code>-O4</code> option.</p>
<p><code>-O5</code> <code>-qOPTimize=5</code></p>	<p>This option is the same as <code>-O4</code>, except that it:</p> <ul style="list-style-type: none"> • Sets the <code>-qipa=level=2</code> option to perform full interprocedural data flow and alias analysis. <p>Note: Later settings of <code>-O</code>, <code>-qcache</code>, <code>-qipa</code>, <code>-qarch</code>, and <code>-qtune</code> options will override the settings implied by the <code>-O5</code> option.</p>

Notes

You can abbreviate `-qoptimize...` to `-qopt...`. For example, `-qnoopt` is equivalent to `-qnooptimize`.

Increasing the level of optimization may or may not result in additional performance improvements, depending on whether additional analysis detects further opportunities for optimization.

Compilations with optimizations may require more time and machine resources than other compilations.

Optimization can cause statements to be moved or deleted, and generally should not be specified along with the `-g` flag for debugging programs. The debugging information produced may not be accurate.

Example

To compile `myprogram.c` for maximum optimization, enter:

```
xlc myprogram.c -O3
```

Related References

"Compiler Command Line Options" on page 51

"arch" on page 70

"cache" on page 86

"float" on page 118

"flttrap" on page 123

"g" on page 130

"hot" on page 134

"ignerrno" on page 140

"ignprag" on page 141

"ipa" on page 151

"langlvl" on page 165

"maxmem" on page 190

"strict" on page 241

"strict_induction" on page 242

"tune" on page 253

See also the *Getting Started with Optimization* section in *Getting Started with* .

See also .

O

Purpose

Specifies an output location for the object, assembler, or executable files created by the compiler. When the **-o** option is used during compiler invocation, *filespec* can be the name of either a file or a directory. When the **-o** option is used during direct linkage-editor invocation, *filespec* can only be the name of a file.

Syntax

►— **-o** *filespec* —►

Notes

When **-o** is specified as part of a compiler invocation, *filespec* can be the relative or absolute path name of either a directory or a file.

1. If *filespec* is the name of a directory, files created by the compiler are placed into that directory.
2. If a directory with the name *filespec* does not exist, the **-o** option specifies that the name of the file produced by the compiler will be *filespec*. For example, the compiler invocation:

```
xlc test.c -c -o new.o
```

produces the object file **new.o** instead of **test.o**, and

```
xlc test.c -o new
```

produces the object file **new** instead of **a.out**, provided there is no directory also named **new**. Otherwise, the default object name **a.out** is used and placed in the **new** directory.

A *filespec* with a C source file suffix (**.c**, **.cpp**, or **.i**), such as *myprog.c* or *myprog.i*, results in an error and neither the compiler nor the linkage editor is invoked.

If you use **-c** and **-o** together and the *filespec* does not specify a directory, you can only compile one source file at a time. In this case, if more than one source file name is listed in the compiler invocation, the compiler issues a warning message and ignores **-o**.

The **-E**, **-P**, and **-qsyntaxonly** options override the **-ofilename** option.

Example

To compile *myprogram.c* so that the resulting file is called **myaccount**, assuming that no directory with name **myaccount** exists, enter:

```
xlc myprogram.c -o myaccount
```

If the directory **myaccount** does exist, the compiler produces the executable file **a.out** and places it in the **myaccount** directory.

Related References

“Compiler Command Line Options” on page 51

“c” on page 84

“E” on page 105

“P” on page 199

“syntaxonly” on page 245

P

Purpose

Preprocesses the C source files named in the compiler invocation and creates an output preprocessed source file, *file_name.i* for each input source file *file_name.c*.

Syntax

▶— -P —▶

Notes

The **-P** option retains all white space including line-feed characters, with the following exceptions:

- All comments are reduced to a single space (unless **-C** is specified).
- Line feeds at the end of preprocessing directives are not retained.
- White space surrounding arguments to function-style macros is not retained.

#line directives are not issued.

The **-P** option cannot accept a preprocessed source file, such as *file_name.i* as input. The compiler will issue an error message.

Source files with unrecognized filename suffixes are treated and preprocessed as C files, and no error message is generated.

In extended mode, the preprocessor interprets the backslash character when it is followed by a new-line character as line-continuation in:

- macro replacement text
- macro arguments
- comments that are on the same line as a preprocessor directive.

Line continuations elsewhere are processed in **ANSI** mode only.

The **-P** option is overridden by the **-E** option. The **-P** option overrides the **-c**, **-o**, and **-qsyntaxonly** option. The **-C** option may be used in conjunction with both the **-E** and **-P** options.

The default is to compile and link-edit C source files to produce an executable file.

Related References

“Compiler Command Line Options” on page 51

“C” on page 83

“c” on page 84

“E” on page 105

“o” on page 198

“syntaxonly” on page 245

p

Purpose

Sets up the object files produced by the compiler for profiling.

Syntax

►► -p ◀◀

Notes

If the `-qtbtable` option is not set, the `-p` option will generate full traceback tables.

When compiling and linking in separate steps, the `-p` option must be specified in both steps.

Example

To compile `myprogram.c` so that it can be used with your operating system's `prof` command, enter:

```
xlc myprogram.c -p
```

Related References

"Compiler Command Line Options" on page 51

"tbtable" on page 248

Also, on the Web see:

`prof` Command section in *Commands Reference, Volume 4: n through r* for information about profiling.

pascal

Purpose

The default **-qpascal** instructs the compiler to ignore the word **pascal** when it appears in type specifiers and function declarations.

Syntax

►► -q nopascal pascal ◄◄

Notes

Specifying the **-qpascal** compiler option instructs the compiler to recognize and accept the keyword **pascal** in type specifiers and function declarations.

This option can be used to improve compatibility of IBM XL C Enterprise Edition programs on some other systems.

Related References

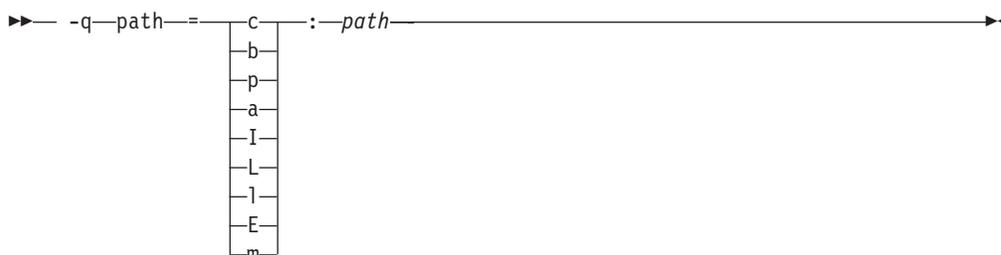
“Compiler Command Line Options” on page 51

path

Purpose

Constructs alternate program names for compiler components. The program and directory *path* specified by this option is used in place of the regular program.

Syntax



where program names are:

Program	Description
c	Compiler front end
b	Compiler back end
p	Compiler preprocessor
a	Assembler
I	Interprocedural Analysis tool - compile phase
L	Interprocedural Analysis tool - link phase
l	Linkage editor
E	CreateExportList utility
m	Linkage helper (munch utility)

Notes

The **-qpath** option overrides the **-Fconfig_file**, **-t**, and **-B** options.

Examples

To compile myprogram.c using a substitute **xlc** compiler in **/lib/tmp/mine/** enter:

```
xlc myprogram.c -qpath=c:/lib/tmp/mine/
```

To compile myprogram.c using a substitute linkage editor in **/lib/tmp/mine/**, enter:

```
xlc myprogram.c -qpath=l:/lib/tmp/mine/
```

Related References

“Compiler Command Line Options” on page 51

“B” on page 78

“F” on page 114

“t” on page 246

pdf1, pdf2

Purpose

Tunes optimizations through *profile-directed feedback* (PDF), where results from sample program execution are used to improve optimization near conditional branches and in frequently executed code sections.

Syntax



Notes

To use PDF, follow these steps:

1. Compile some or all of the source files in a program with the **-qpdf1** option. You need to specify at least the **-O2** optimizing option and you also need to link with at least **-O2** in effect as well. Pay special attention to the compiler options that you use to compile the files, because you will need to use the same options later.

In a large application, concentrate on those areas of the code that can benefit most from optimization. You do not need to compile all of the application's code with the **-qpdf1** option.

2. Run the program all the way through using a typical data set. The program records profiling information when it finishes. You can run the program multiple times with different data sets, and the profiling information is accumulated to provide an accurate count of how often branches are taken and blocks of code are executed.

Important: Use data that is representative of the data that will be used during a normal run of your finished program.

3. Relink your program using the same compiler options as before, but change **-qpdf1** to **-qpdf2**. Remember that **-L**, **-l**, and some others are linker options, and you can change them at this point. In this second compilation, the accumulated profiling information is used to fine-tune the optimizations. The resulting program contains no profiling overhead and runs at full speed.

For best performance, use the **-O3**, **-O4**, or **-O5** option with all compilations when you use PDF.

The profile is placed in the current working directory or in the directory that the PDFDIR environment variable names, if that variable is set.

To avoid wasting compilation and execution time, make sure that the PDFDIR environment variable is set to an absolute path. Otherwise, you might run the application from the wrong directory, and it will not be able to locate the profile data files. When that happens, the program may not be optimized correctly or may be stopped by a segmentation fault. A segmentation fault might also happen if you change the value of the PDFDIR variable and execute the application before finishing the PDF process.

Because this option requires compiling the entire application twice, it is intended to be used after other debugging and tuning is finished, as one of the last steps before putting the application into production.

Restrictions

- PDF optimizations require at least the **-O2** optimization level.
- You must compile the main program with PDF for profiling information to be collected at run time.
- Do not compile or run two different applications that use the same PDFDIR directory at the same time, unless you have used the **-qipa=pdfname** suboption to distinguish the sets of profiling information.
- You must use the same set of compiler options at all compilation steps for a particular program. Otherwise, PDF cannot optimize your program correctly and may even slow it down. All compiler settings must be the same, including any supplied by configuration files.
- Avoid mixing PDF files created by the current version level of XL C Enterprise Edition with PDF files created by other version levels of the compiler.
- If **-qipa** is not invoked either directly or through other options, **-qpdf1** and **-qpdf2** will invoke the **-qipa=level=0** option.
- If you compile a program with **-qpdf1**, remember that it will generate profiling information when it runs, which involves some performance overhead. This overhead goes away when you recompile with **-qpdf2** or with no PDF at all.

The following utility programs, found in **/usr/xlopt/bin**, are available for managing the **PDFDIR** directory:

cleanpdf cleanpdf [*pathname*]

Removes all profiling information from the *pathname* directory; or if *pathname* is not specified, from the **PDFDIR** directory; or if **PDFDIR** is not set, from the current directory. Removing profiling information reduces run-time overhead if you change the program and then go through the PDF process again.

Run **cleanpdf** only when you are finished with the PDF process for a particular application. Otherwise, if you want to resume using PDF with that application, you will need to recompile all of the files again with **-qpdf1**.

mergepdf mergepdf [-r *scaling*] input {[-r *scaling*] input} ... -o *output* [-n] [-v]

Merges two or more PDF records into a single PDF output record.

-r *scaling* Specifies the scaling ratio for the PDF record file. This value must be greater than zero and can be either an integer or floating point value. If not specified, a ratio of 1.0 is assumed.

record Specifies the name of a PDF input record file, or a directory that contains PDF record files.

-o *output* Specifies the name of the PDF output record file, or a directory to which the merged output will be written.

-n If specified, PDF record files are not normalized. If not specified, **mergepdf** normalizes records based on an internally-calculated ratio before applying any user-defined scaling factor.

-v Specifies verbose mode, and causes internal and user-specified scaling ratios to be displayed to the screen.

resetpdf resetpdf [*pathname*]

Same as `cleanpdf` [*pathname*], described above.

showpdf showpdf

Displays the call and block counts for all procedures executed in a program run. To use this command, you must first compile your application specifying both `-qpdf1` and `-qshowpdf` compiler options on the command line.

Examples

Here is a simple example:

```
/* Set the PDFDIR variable. */
export PDFDIR=$HOME/project_dir

/* Compile all files with -qpdf1. */
xlc -qpdf1 -O3 file1.c file2.c file3.c

/* Run with one set of input data. */
a.out <sample.data

/* Recompile all files with -qpdf2. */
xlc -qpdf2 -O3 file1.c file2.c file3.c

/* The program should now run faster than
   without PDF if the sample data is typical. */
```

Here is a more elaborate example.

```
/* Set the PDFDIR variable. */
export PDFDIR=$HOME/project_dir

/* Compile most of the files with -qpdf1. */
xlc -qpdf1 -O3 -c file1.c file2.c file3.c

/* This file is not so important to optimize.
xlc -c file4.c

/* Non-PDF object files such as file4.o can be linked in. */
xlc -qpdf1 file1.o file2.o file3.o file4.o

/* Run several times with different input data. */
a.out <polar_orbit.data
a.out <elliptical_orbit.data
a.out <geosynchronous_orbit.data

/* No need to recompile the source of non-PDF object files (file4.c). */
xlc -qpdf2 -O3 file1.c file2.c file3.c

/* Link all the object files into the final application. */
xlc file1.o file2.o file3.o file4.o
```

Related References

“Compiler Command Line Options” on page 51

“ipa” on page 151

“O, optimize” on page 194

“showpdf” on page 229

pg

Purpose

Sets up the object files for profiling, but provides more information than is provided by the **-p** option.

If the **-qtbtable** option is not set, the **-pg** option will generate full traceback tables.

Syntax

▶▶ — `-pg` —▶▶

Example

To compile `myprogram.c` for use with your operating system's **gprof** command, enter:

```
xlc myprogram.c -pg
```

Remember to compile *and* link with the **-pg** option. For example:

```
xlc myprogram.c -pg -c  
xlc myprogram.o -pg -o program
```

Related References

"Compiler Command Line Options" on page 51

"tbtable" on page 248

Also, on the Web see:

`gprof` Command section in *Commands Reference, Volume 2: d through h* for information about profiling.

phsinfo

Purpose

Reports the time taken in each compilation phase. Phase information is sent to standard output.

Syntax

►► -q nophsinfo
phsinfo ◀◀

Notes

The output takes the form *number1/number2* for each phase where *number1* represents the CPU time used by the compiler and *number2* represents the total of the compiler time and the time that the CPU spends handling system calls.

Example

To compile `myprogram.c` and report the time taken for each phase of the compilation, enter:

```
xlc myprogram.c -qphsinfo
```

The output will look similar to:

```
C Init      - Phase Ends;  0.010/  0.040
IL Gen     - Phase Ends;  0.040/  0.070
W-TRANS    - Phase Ends;  0.000/  0.010
OPTIMIZ    - Phase Ends;  0.000/  0.000
REGALLO    - Phase Ends;  0.000/  0.000
AS         - Phase Ends;  0.000/  0.000
```

Compiling the same program with **-O4** gives:

```
C Init      - Phase Ends;  0.010/  0.040
IL Gen     - Phase Ends;  0.060/  0.070
IPA        - Phase Ends;  0.060/  0.070
IPA        - Phase Ends;  0.070/  0.110
W-TRANS    - Phase Ends;  0.060/  0.180
OPTIMIZ    - Phase Ends;  0.010/  0.010
REGALLO    - Phase Ends;  0.010/  0.020
AS         - Phase Ends;  0.000/  0.000
```

Related References

“Compiler Command Line Options” on page 51

pic

Purpose

Instructs the compiler to generate Position-Independent Code suitable for use in shared libraries.

Syntax



where

- | | |
|-------|--|
| small | Instructs the compiler to assume that the size of the Table of Contents is no larger than 64 Kb. |
| large | Allows the Table of Contents to be larger than 64 Kb in size, allowing more addresses to be stored in the table. Code generated with this option is usually larger than that generated with -qpic=small . |

Notes

If **-qpic** is specified without any suboptions, **-qpic=small** is assumed.

The **-qpic** option is implied if either the **-G** or **-qmkshrobj** compiler option is specified.

Specifying **-qpic=large** has the same effect as passing **-bbigtoc** to **ld**.

Example

To compile a shared library **libmylib.so**, use the following command:

```
xlc mylib.c -qpic=small -Wl, -shared, -soname="libmylib.so.1" -o libmylib.so.1
```

Refer to the **ld** command in your operating system documentation for more information about the **-shared** and **-soname** options.

Related References

“Compiler Command Line Options” on page 51

“32, 64” on page 60

“G” on page 129

“mkshrobj” on page 192

prefetch

Purpose

Enables generation of prefetching instructions such as `dcbt` and `dcbz` in compiled code.

Syntax

►► — `-q` prefetch
noprefetch —►►

Notes

By default, the compiler may insert prefetch instructions in compiled code. The `-qnoprefetch` option lets you disable this feature.

The `-qnoprefetch` option will not prevent built-in functions such as `__prefetch_by_stream()` from generating prefetch instructions.

Related References

“Compiler Command Line Options” on page 51

print

Purpose

Enables or suppresses listings. Specifying **-qnoprint** overrides all listing-producing options, regardless of where they are specified, to suppress listings.

Syntax



Notes

The default of **-qprint** enables listings if they are requested by other compiler options. These options are:

- `-qattr`
- `-qlist`
- `-qlistopt`
- `-qsource`
- `-qxref`

Example

To compile `myprogram.c` and suppress all listings, even if some files have **#pragma options source** and similar directives, enter:

```
xlc myprogram.c -qnoprint
```

Related References

“Compiler Command Line Options” on page 51

“attr” on page 77

“list” on page 175

“listopt” on page 176

“source” on page 234

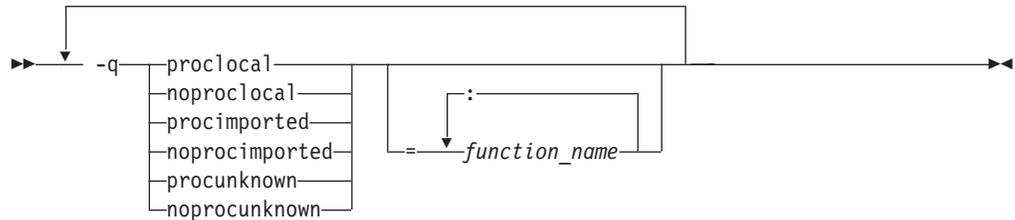
“xref” on page 268

proclocal, procimported, procunknown

Purpose

Marks functions as local, imported, or unknown.

Syntax



See also “#pragma options” on page 299.

Default

The default is to assume that all functions whose definition is in the current compilation unit are local **proclocal**, and that all other functions are unknown **procunknown**. If any functions that are marked as local resolve to shared library functions, the linkage editor will detect the error and issue warnings.

Notes

Available suboptions are:

Local functions	Local functions are statically bound with the functions that call them. Specifying -qproclocal changes the default to assume that all functions are local. -qproclocal=names marks the named functions as local, where <i>names</i> is a list of function identifiers separated by colons (:). The default is not changed. Smaller, faster code is generated for calls to functions marked as local.
Imported functions	Imported functions are dynamically bound with a shared portion of a library. -qprocimported changes the default to assume that all functions are imported. Specifying -qprocimported=names marks the named functions as imported, where <i>names</i> is a list of function identifiers separated by colons (:). The default is not changed. Code generated for calls to functions marked as imported may be larger, but is faster than the default code sequence generated for functions marked as unknown. If marked functions resolve to statically bound objects, the generated code may be larger and run more slowly than the default code sequence generated for unknown functions.
Unknown functions	Unknown functions are resolved to either statically or dynamically bound objects during link-editing. Specifying -qprocunknown changes the default to assume that all functions are unknown. -qprocunknown=names marks the named functions as unknown, where <i>names</i> is a list of function identifiers separated by colons (:). The default is not changed.

Conflicts among the procedure-marking options are resolved in the following manner:

Options that list function names	The last explicit specification for a particular function name is used.
Options that change the default	This form does not specify a name list. The last option specified is the default for functions not explicitly listed in the name-list form.

Examples

1. To compile myprogram.c along with the archive library **oldprogs.a** so that:

- functions **fun** and **sun** are specified as **local**,
- functions **moon** and **stars** are specified as **imported**, and,
- function **venus** is specified as **unknown**,

enter:

```
xlc myprogram.c oldprogs.a -qprolocal=fun(int):sun()
-qprocimported=moon():stars(float) -qprocunknown=venus()
```

2. The following example shows typical error messages that result when a function marked as local instead resolves to a shared library function.

```
int main(void)
{
    printf("Just in function fool()\n");
    printf("Just in function fool()\n");
}
```

```
ld: 0711-768 WARNING: Object t.o, section 1, function .printf:
The branch at address 0x18 is not followed by a recognized no-op or
TOC-reload instruction. The unrecognized instruction is 0x83E1004C.
```

An executable file is produced, but it will not run. The error message indicates that a call to **printf** in object file **t.o** caused the problem. When you have confirmed that the called routine should be imported from a shared object, recompile the source file that caused the warning and explicitly mark **printf** as imported. For example:

```
xlc -c -qprocimported=printf t.c
```

Related References

“Compiler Command Line Options” on page 51

proto

Purpose

If this option is set, the compiler assumes that all functions are prototyped.

Syntax

►► — -q — no proto —►►

Notes

This option asserts that procedure call points agree with their declarations even if the procedure has not been prototyped.

Callers can pass floating-point arguments in floating-point registers only and not in General-Purpose Registers (GPRs). The compiler assumes that the arguments on procedure calls are the same types as the corresponding parameters of the procedure definition.

The compiler will issue warnings for functions that do not have prototypes.

Example

To compile `my_c_program.c` to assume that all functions are prototyped, enter:

```
xlc my_c_program.c -qproto
```

Related References

“Compiler Command Line Options” on page 51

Q

Purpose

This option specifies which specific functions the compiler should attempt to inline.

Syntax



In the C language, the following **-Q** options apply:

- Q** Attempts to inline all appropriate functions with 20 executable source statements or fewer, subject to the setting of any of the suboptions to the **-Q** option. If **-Q** is specified last, all functions are inlined.
- Q!** Does not inline any functions. If **-Q!** is specified last, no functions are inlined.
- Q-names** Does not inline functions listed by *names*. Separate each function name in *names* with a colon (:). All other appropriate functions are inlined. The option implies **-Q**.

For example:

```
-Q-salary:taxes:expenses:benefits
```

causes all functions except those named **salary**, **taxes**, **expenses**, or **benefits** to be inlined if possible.

A warning message is issued for functions that are not defined in the source file.

- Q+names** Attempts to inline the functions listed by *names* and any other appropriate functions. Each function name in *names* must be separated by a colon (:). The option implies **-Q**.

For example,

```
-Q+food:clothes:vacation
```

causes all functions named **food**, **clothes**, or **vacation** to be inlined if possible, along with any other functions eligible for inlining.

A warning message is issued for functions that are not defined in the source file or that are defined but cannot be inlined.

This suboption overrides any setting of the *threshold* value. You can use a threshold value of zero along with **-Q+names** to inline specific functions. For example:

```
-Q=0
```

followed by:

```
-Q+salary:taxes:benefits
```

causes *only* the functions named **salary**, **taxes**, or **benefits** to be inlined, if possible, and no others.

`-Q=threshold` Sets a size limit on the functions to be inlined. The number of executable statements must be less than or equal to *threshold* for the function to be inlined. *threshold* must be a positive integer. The default value is 20. Specifying a threshold value of 0 causes no functions to be inlined except those functions marked with the `__inline`, `_Inline`, or `_inline` keywords.

The *threshold* value applies to logical C statements. Declarations are not counted, as you can see in the example below:

```
increment()
{
    int a, b, i;
    for (i=0; i<10; i++) /* statement 1 */
    {
        a=i;             /* statement 2 */
        b=i;             /* statement 3 */
    }
}
```

Default

The default is to treat inline specifications as a hint to the compiler. Whether or not inlining occurs may also be dependent on other options that you select:

- If you optimize your programs, (specify the `-O` option) the compiler attempts to inline the functions declared as inline.

Notes

The `-Q` option is functionally equivalent to the `-qinline` option.

If you specify the `-g` option (to generate debug information), inlining may be affected. See the information for the `-g` compiler option.

Because inlining does not always improve run time, you should test the effects of this option on your code.

Do not attempt to inline recursive or mutually recursive functions.

Normally, application performance is optimized if you request optimization (`-O` option), and compiler performance is optimized if you do not request optimization.

The `inline`, `_inline`, `_Inline`, and `__inline` language keywords override all `-Q` options except `-Q!`. The compiler will try to inline functions marked with these keywords regardless of other `-Q` option settings.

To maximize inlining, specify optimization (`-O`) and also specify the appropriate `-Q` options for the C language.

Examples

To compile the program `myprogram.c` so that no functions are inlined, enter:

```
xlc myprogram.c -O -Q!
```

To compile the program `my_c_program.c` so that the compiler attempts to inline functions of fewer than 12 lines, enter:

```
xlc my_c_program.c -O -Q=12
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

“inline” on page 147

“O, optimize” on page 194

“Q” on page 214

“The inline, _Inline, _inline, and __inline Function Specifiers” on page 149

r

Purpose

Produces a relocatable object. This permits the output file to be produced even though it contains unresolved symbols.

Syntax

▶— -r —▶

Notes

A file produced with this flag is expected to be used as a file parameter in another call to `xlc`.

Example

To compile `myprogram.c` and `myprog2.c` into a single object file `mytest.o`, enter:

```
xlc myprogram.c myprog2.c -r -o mytest.o
```

Related References

“Compiler Command Line Options” on page 51

report

Purpose

Instructs the compiler to produce transformation reports that show how program loops are parallelized and/or optimized. The transformation reports are included as part of the compiler listing.

Syntax

►► -q noreport
report _____►►

Notes

This option has no effect unless **-qhot** or **-qsmp** are also in effect.

Specifying **-qreport** together with **-qhot** instructs the compiler to produce a pseudo-C code listing and summary showing how loops are transformed. You can use this information to tune the performance of loops in your program.

Specifying **-qreport** together with **-qsmp** instructs the compiler to also produce a report showing how the program deals with data and automatic parallelization of loops in your program. You can use this information to determine how loops in your program are or are not parallelized.

The pseudo-C code listing is not intended to be compilable. Do not include any of the pseudo-C code in your program, and do not explicitly call any of the internal routines whose names may appear in the pseudo-C code listing.

Example

To compile `myprogram.c` so the compiler listing includes a report showing how loops are optimized, enter:

```
xlc -qhot -O3 -qreport myprogram.c
```

To compile `myprogram.c` so the compiler listing also includes a report showing how parallelized loops are transformed, enter:

```
xlc -qsmp -O3 -qreport myprogram.c
```

Related References

“Compiler Command Line Options” on page 51

“hot” on page 134

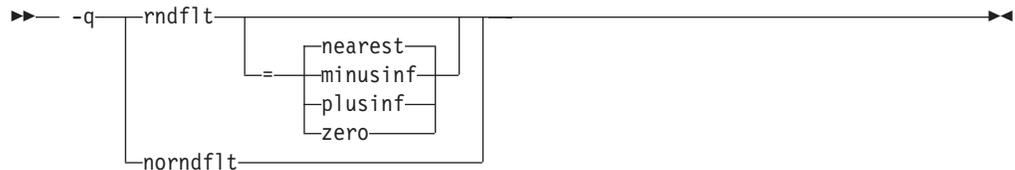
“smp” on page 232

rndflt

Purpose

This option controls the compile-time rounding mode of constant floating point expressions. It does not affect run-time rounding.

Syntax



where available rounding options are:

Option	Effect
nearest	Round to nearest representable number. This is the default.
minusinf	Round toward minus infinity.
plusinf	Round toward plus infinity.
zero	Round toward zero.

Notes

Compile-time floating-point arithmetic can have two effects on program results:

- In specific cases, the result of a computation at compile time might differ slightly from the result that would have been calculated at run time. The reason is that more rounding operations occur at compile time. For example, where a multiply-add floating point operation might be used at run time, separate multiply and add operations might be used at compile time, producing a slightly different result.
- Computations that produce exceptions can be folded to the IEEE result that would have been produced by default in a run-time operation. This would prevent an exception from occurring at run time. The `-qfltrp` option can be used to generate instructions that detect and trap floating-point exceptions.

In general, code that affects the rounding mode at run time should be compiled with the option that matches that rounding mode. For example, when the following program is compiled, the expression `1.0/3.0` is folded at compile time into a double-precision result:

```
main()
{
    float x, y;
    int i;
    x = 1.0/3.0;
    i = *(int *)&x;
    printf("1/3 = %.8x\n", i);
    x = 1.0;
    y = 3.0;
    x = x/y;
    i = *(int *)&x;
    printf("1/3 = %.8x\n", i);
}
```

This result is then converted to single precision and stored in `float x`.

The **-qfloat=nofold** option can be specified to suppress all compile-time folding of floating-point computations. For example, the following code fragment may be evaluated either at compile time or at run time, depending on the setting of **-qfloat** and other options:

```
x = 1.0;  
y = 3.0;  
x = x/y;
```

The **-qrndflt** option only affects compile-time rounding of floating-point computations. If this code is evaluated at run time, the default run-time rounding of “round to nearest” is still in effect and takes precedence over the compile-time rounding mode.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“flttrap” on page 123

“y” on page 269

rrm

Purpose

Prevents floating-point optimizations that are incompatible with run-time rounding to plus and minus infinity modes.

Syntax

►► -q norm
rrm ◀◀

See also “#pragma options” on page 299.

Notes

*This option is obsolete. Use **-qfloat=rrm** in your new applications.*

This option informs the compiler that, at run time, the floating-point rounding mode may change or that the mode is not set to **-yn** (rounding to the nearest representable number.)

-qrrm must also be specified if the Floating Point Status and Control register is changed at run time.

The default, **-qnorm**, generates code that is compatible with run-time rounding modes **nearest** and **zero**. For a list of rounding mode options, see the **-y** compiler option.

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“#pragma options” on page 299

S

Purpose

Generates an assembler language file (.s) for each source file. The resulting .s files can be assembled to produce object .o files or an executable file (a.out).

Syntax

►— -S —————►

Notes

You can invoke the assembler with the `xlc` command. For example,

```
xlc myprogram.s
```

will invoke the assembler, and if successful, the loader to create an executable file, **a.out**.

If you specify `-S` with `-E` or `-P`, `-E` or `-P` takes precedence. Order of precedence holds regardless of the order in which they were specified on the command line.

You can use the `-o` option to specify the name of the file produced only if no more than one source file is supplied. For example, the following is *not* valid:

```
xlc myprogram1.c myprogram2.c -o -S
```

Restrictions

The generated assembler files do not include all the data that is included in a .o file by the `-g` or `-qipa` options.

Examples

1. To compile `myprogram.c` to produce an assembler language file **myprogram.s**, enter:

```
xlc myprogram.c -S
```

2. To assemble this program to produce an object file **myprogram.o**, enter:

```
xlc myprogram.s -c
```

3. To compile `myprogram.c` to produce an assembler language file **asmprogram.s**, enter:

```
xlc myprogram.c -S -o asmprogram.s
```

Related References

“Compiler Command Line Options” on page 51

“E” on page 105

“g” on page 130

“ipa” on page 151

“o” on page 198

“P” on page 199

“thtable” on page 248

Also, on the Web see:

- AIX 5L for POWER-based Systems: Assembler Language Reference
- Files Reference

S

Purpose

This option strips the symbol table, line number information, and relocation information from the output file. Specifying **-s** saves space, but limits the usefulness of traditional debug programs when you are generating debug information using options such as **-g**.

Syntax

▶▶ — **-s** —————▶▶

Notes

Using the strip command has the same effect.

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

saveopt

Purpose

Saves the command line compiler options into an object file.

Syntax

►► -q nosaveopt
saveopt _____►►

Notes

This option lets you save command line compiler options into the object file you are compiling. The option has effect only when compiling to an object (.o) file.

The string is saved in the following format:

►► @(#)opt-*B* f
c
C -*B*-*stanza*-*B*-*options* _____►►

where:

- B* Indicates a space.
- f** Signifies a Fortran language compilation.
- c** Signifies a C language compilation.
- C** Signifies a C++ language compilation.
- stanza* Specifies the driver used for the compilation, for example, c89 or xlC.
- options* The list of command line options specified on the command line, with individual options separated by spaces.

Related References

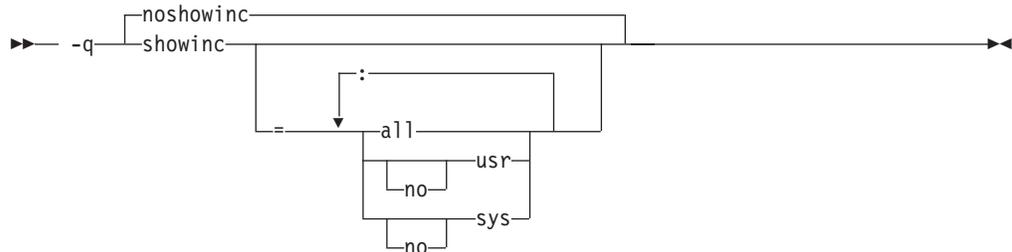
“Compiler Command Line Options” on page 51

showinc

Purpose

Used with **-qsource** to selectively show user header files (includes using " ") or system header files (includes using < >) in the program source listing.

Syntax



where options are:

- | | |
|-----------|---|
| noshowinc | Do not show user or system include files in the program source listing.
This is the same as specifying -qshowinc=nousr:nosys . |
| showinc | Show both user and system include files in the program source listing.
This is the same as specifying -qshowinc=usr:sys or -qshowinc=all . |
| all | Show both user and system include files in the program source listing.
This is the same as specifying -qshowinc or -qshowinc=usr:sys . |
| usr | Show user include files in the program source listing. |
| sys | Show system include files in the program source listing. |

See also “#pragma options” on page 299.

Notes

This option has effect only when the **-qlist** or **-qsource** compiler options are in effect.

Example

To compile myprogram.c so that all included files appear in the source listing, enter:

```
xlc myprogram.c -qsource -qshowinc
```

Related References

“Compiler Command Line Options” on page 51

“source” on page 234

“#pragma options” on page 299

showpdf

Purpose

Used together with **-qpdf1** to add additional call and block count profiling information to an executable.

Syntax

►► -q no**showpdf** **showpdf** ◀◀

Notes

This option has effect only when specified together with the **-qpdf1** compiler option.

When specified together with **-qpdf1**, the compiler will insert additional profiling information into the compiled application to collect call and block counts for all procedures in the application. Running the compiled application will record the call and block counts to the file **._pdf**.

After you run your application with training data, you can retrieve the contents of the **._pdf** file with the **showpdf** utility. This utility is described in the **-qpdf** pages.

Example

The example assumes the following source for program file `hello.c`:

```
#include <stdio.h>

void HelloWorld()
{
    printf("Hello World");
}

main()
{
    HelloWorld();
}
```

Compile the source with:

```
xlc -qpdf1 -qshowpdf hello.c
```

Run the resulting program executable:

```
a.out
```

Run the **showpdf** utility to display the call and block counts for the executable:

```
showpdf
```

Something similar to the following will be returned by the **showpdf** utility:

```
HelloWorld(4): 1 (hello.c)

Call Counters:
 5 | 1 printf(6)

Call coverage = 100% ( 1/1 )

Block Counters:
 3-5 | 1
 6 |
 6 | 1
```

Block coverage = 100% (2/2)

main(5): 1 (hello.c)

Call Counters:

10 | 1 HelloWorld(4)

Call coverage = 100% (1/1)

Block Counters:

8-11 | 1
11 |

Block coverage = 100% (1/1)

Total Call coverage = 100% (2/2)

Total Block coverage = 100% (3/3)

Related References

“Compiler Command Line Options” on page 51

“pdf1, pdf2” on page 203

smallstack

Purpose

Instructs the compiler to reduce the size of the stack frame.

Syntax

►► -q nosmallstack
smallstack ◀◀

Notes

AIX limits the stack size to 256 MB. Programs that allocate large amounts of data to the stack, such as threaded programs, may result in stack overflows. This option can reduce the size of the stack frame to help avoid overflows.

This option is only valid when used together with IPA (**-qipa**, **-O4**, **-O5** compiler options).

Specifying this option may adversely affect program performance.

Example

To compile myprogram.c to use a small stack frame, enter:

```
xlc myprogram.c -qipa -qsmallstack
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

“ipa” on page 151

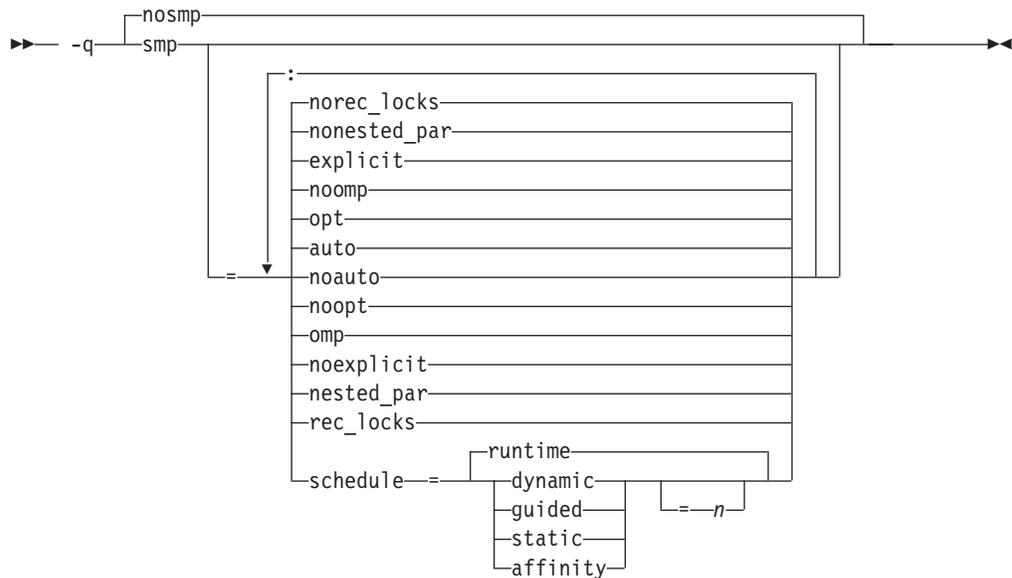
“O, optimize” on page 194

smp

Purpose

Enables parallelization of program code.

Syntax



where:

- | | |
|-------------------|--|
| auto | Enables automatic parallelization and optimization of program code. |
| noauto | Disables automatic parallelization of program code. Program code explicitly parallelized with SMP or OpenMP pragma statements is optimized. |
| opt | Enables automatic parallelization and optimization of program code. |
| noopt | Enables automatic parallelization, but disables optimization of parallelized program code. Use this setting when debugging parallelized program code. |
| omp | Enables strict compliance to the OpenMP 2.0 standard. Automatic parallelization is disabled. Parallelized program code is optimized. Only OpenMP parallelization pragmas are recognized. |
| noomp | Enables automatic parallelization and optimization of program code. |
| explicit | Enables pragmas controlling explicit parallelization of loops. |
| noexplicit | Disables pragmas controlling explicit parallelization of loops. |

nested_par	If specified, nested parallel constructs are not serialized. nested_par does not provide true nested parallelism because it does not cause new team of threads to be created for nested parallel regions. Instead, threads that are currently available are re-used.
	This option should be used with caution. Depending on the number of threads available and the amount of work in an outer loop, inner loops could be executed sequentially even if this option is in effect. Parallelization overhead may not necessarily be offset by program performance gains.
nonested_par	Disables parallelization of nested parallel constructs.
rec_locks	If specified, recursive locks are used, and nested critical sections will not cause a deadlock.
norec_locks	If specified, recursive locks are not used.
schedule=sched_type[=n]	Specifies what kind of scheduling algorithms and chunk size (<i>n</i>) are used for loops to which no other scheduling algorithm has been explicitly assigned in the source code. If <i>sched_type</i> is not specified, schedule=runtime is assumed for the default setting.

Notes

- **-qsmp** must be used only with thread-safe compiler mode invocations such as **xlcr**. These invocations ensure that the **pthread**, **xlsmp**, and thread-safe versions of all default run-time libraries are linked to the resulting executable.
- The **-qnosmp** default option setting specifies that no code should be generated for parallelization directives, though syntax checking will still be performed. Use **-qignprag=omp:ibm** to completely ignore parallelization directives.
- Specifying **-qsmp** without suboptions is equivalent to specifying **-qsmp=auto:explicit:noomp:norec_locks:nonested_par:schedule=runtime** or **-qsmp=opt:explicit:noomp:norec_locks:nonested_par:schedule=runtime**.
- Specifying **-qsmp** implicitly sets **-O2**. The **-qsmp** option overrides **-qnooptimize**, but does not override **-O3**, **-O4**, or **-O5**. When debugging parallelized program code, you can disable optimization in parallelized program code by specifying **qsmp=noopt**.
- Specifying **-qsmp** defines the **_IBMSMP** preprocessing macro.

Related Concepts

“Program Parallelization” on page 11

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

“Control Parallel Processing with Pragmas” on page 39

Related References

“Compiler Command Line Options” on page 51

“O, optimize” on page 194

“threaded” on page 249

“Pragmas to Control Parallel Processing” on page 321

“#pragma ibm schedule” on page 329

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

source

Purpose

Produces a compiler listing and includes source code.

Syntax

►► — -q — nosource
source —————►►

See also “#pragma options” on page 299.

Notes

The **-qnoprint** option overrides this option.

Parts of the source can be selectively printed by using pairs of **#pragma options source** and **#pragma options nosource** preprocessor directives throughout your source program. The source following **#pragma options source** and preceding **#pragma options nosource** is printed.

Examples

To compile `myprogram.c` to produce a compiler listing that includes the source for `myprogram.c`, enter:

```
xlc myprogram.c -qsource
```

Do not use the **-qsource** compiler option if you want the compiler listing to show only selected parts of your program source. The following code causes the only the source found between the **#pragma options source** and **#pragma options nosource** directives to be included in the compiler listing:

```
#pragma options source
. . .
/* Source code to be included in the compiler listing
   is bracketed by #pragma options directives.
*/
. . .
#pragma options nosource
```

Related References

“Compiler Command Line Options” on page 51

“print” on page 210

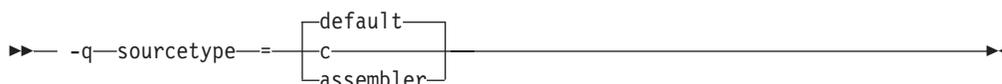
“#pragma options” on page 299

sourcetype

Purpose

Instructs the compiler to treat all source files as if they are the source type specified by this option, regardless of actual source filename suffix.

Syntax



where:

default	The compiler assumes that the programming language of a source file will be implied by its filename suffix.
c	The compiler compiles all source files following this option as if they are C language source files.
assembler	The compiler compiles all source files following this option as if they are Assembler language source files.

Notes

The **-qsourcetype** option instructs the compiler to not rely on the filename suffix, and to instead assume a source type as specified by the option.

Ordinarily, the compiler uses the filename suffix of source files specified on the command line to determine the type of the source file. For example, a `.c` suffix normally implies C source code, a `.C` suffix normally implies C++ source code, and the compiler will treat them as follows:

hello.c	The file is compiled as a C file.
hello.C	The file is compiled as a C++ file, assuming a C++ compiler is available.

This applies whether the file system is case-sensitive or not. However, in a case-insensitive file system, the above two compilations refer to the same physical file. That is, the compiler still recognizes the case difference of the filename argument on the command line and determines the source type accordingly, but will ignore the case when retrieving the file from the file system.

Examples

To treat the source file **hello.C** as being a C language source file, enter:

```
xlc -qsourcetype=c hello.C
```

Related References

“Compiler Command Line Options” on page 51

spill

Purpose

Specifies the register allocation spill area as being *size* bytes.

Syntax

►► -q-spill=size ◀◀

See also “#pragma options” on page 299.

Notes

If your program is very complex, or if there are too many computations to hold in registers at one time and your program needs temporary storage, you might need to increase this area. Do not enlarge the spill area unless the compiler issues a message requesting a larger spill area. In case of a conflict, the largest spill area specified is used.

Example

If you received a warning message when compiling myprogram.c and want to compile it specifying a spill area of **900** entries, enter:

```
xlc myprogram.c -qspill=900
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

spnans

Purpose

Generates extra instructions to detect signalling NaN on conversion from single precision to double precision. The **-qnospnans** option specifies that this conversion need not be detected.

Syntax

►► -q nospnans
spnans _____ ►►

See “#pragma options” on page 299.

Notes

*This option is obsolete. Use **-qfloat=nans** in your new applications.*

The **-qhsflt** option overrides the **-qspnans** option

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“hsflt” on page 136

“#pragma options” on page 299

srcmsg

Purpose

Adds the corresponding source code lines to the diagnostic messages in the **stderr** file.

Syntax

►► -q nosrcmsg
srcmsg ◀◀

See also “#pragma options” on page 299.

Notes

The compiler reconstructs the source line or partial source line to which the diagnostic message refers and displays it before the diagnostic message. A pointer to the column position of the error may also be displayed. Specifying **-qnosrcmsg** suppresses the generation of both the source line and the finger line, and the error message simply shows the file, line and column where the error occurred.

The reconstructed source line represents the line as it appears after macro expansion. At times, the line may be only partially reconstructed. The characters “...” at the start or end of the displayed line indicate that some of the source line has not been displayed.

The default (**-qnosrcmsg**) displays concise messages that can be parsed. Instead of giving the source line and pointers for each error, a single line is displayed, showing the name of the source file with the error, the line and character column position of the error, and the message itself.

Example

To compile `myprogram.c` so that the source line is displayed along with the diagnostic message when an error occurs, enter:

```
xlc myprogram.c -qsrcmsg
```

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

statsym

Purpose

Adds user-defined, nonexternal names that have a persistent storage class, such as initialized and uninitialized static variables, to the name list (the symbol table of `xcoff` objects).

Syntax

►► — -q — nostatsym
statsym —————►►

Default

The default is to not add static variables to the symbol table. However, static functions are added to the symbol table.

Example

To compile `myprogram.c` so that static symbols are added to the symbol table, enter:

```
xlc myprogram.c -qstatsym
```

Related References

“Compiler Command Line Options” on page 51

stdinc

Purpose

Specifies which directories are used for files included by the **#include** *<file_name>* and **#include** *"file_name"* directives. The **-qnostdinc** option excludes the standard include directories from the search path.

Syntax

►► -q stdinc
nostdinc ◀◀

See also “#pragma options” on page 299.

Notes

If you specify **-qnostdinc**, the compiler will not search the default search path directories unless you explicitly add it with the **-I***directory* option.

If a full (absolute) path name is specified, this option has no effect on that path name. It will still have an effect on all relative path names.

-qnostdinc is independent of **-qidirfirst**. (**-qidirfirst** searches the directory specified with **-I** *directory* before searching the directory where the current source file resides.

The search order for files is described in *Directory Search Sequence for Include Files Using Relative Path Names*.

The last valid **#pragma options [NO]STDINC** remains in effect until replaced by a subsequent **#pragma options [NO]STDINC**.

Example

To compile `myprogram.c` so that the directory `/tmp/myfiles` is searched for a file included in `myprogram.c` with the **#include** `"myinc.h"` directive, enter:

```
xlc myprogram.c -qnostdinc -I/tmp/myfiles
```

Related References

“Compiler Command Line Options” on page 51

“I” on page 138

“idirfirst” on page 139

“#pragma options” on page 299

“Directory Search Sequence for Include Files Using Relative Path Names” on page 37

strict

Purpose

Turns off the aggressive optimizations that have the potential to alter the semantics of your program.

Syntax



See also “#pragma options” on page 299.

Default

- **-qnostrict** with optimization levels of **-O3** or higher.
- **-qstrict** otherwise.

Notes

-qstrict turns off the following optimizations:

- Performing code motion and scheduling on computations such as loads and floating-point computations that may trigger an exception.
- Relaxing conformance to IEEE rules.
- Reassociating floating-point expressions.

This option is only valid with **-O2** or higher optimization levels.

-qstrict sets **-qfloat=noftint:norsqrt**.

-qnostrict sets **-qfloat=fltint:rsqrt**.

You can use **-qfloat=fltint** and **-qfloat=rsqrt** to override the **-qstrict** settings.

For example:

- Using **-O3 -qstrict -qfloat=fltint** means that **-qfloat=fltint** is in effect, but there are no other aggressive optimizations.
- Using **-O3 -qnostrict -qfloat=norsqrt** means that the compiler performs all aggressive optimizations except **-qfloat=rsqrt**.

If there is a conflict between the options set with **-qnostrict** and **-qfloat=options**, the last option specified is recognized.

Example

To compile `myprogram.c` so that the aggressive optimizations of **-O3** are turned off, range checking is turned off (**-qfloat=fltint**), and division by the result of a square root is replaced by multiplying by the reciprocal (**-qfloat=rsqrt**), enter:

```
xlc myprogram.c -O3 -qstrict -qfloat=fltint:rsqrt
```

Related References

“Compiler Command Line Options” on page 51

“float” on page 118

“O, optimize” on page 194

“#pragma options” on page 299

strict_induction

Purpose

Disables loop induction variable optimizations that have the potential to alter the semantics of your program. Such optimizations can change the result of a program if truncation or sign extension of a loop induction variable should occur as a result of variable overflow or wrap-around.

Syntax

►► — -q —┐nostrict_induction┐—————►►
 └strict_induction┘

Default

- **-qnostrict_induction** with optimization levels 3 or higher.
- **-qstrict_induction** otherwise.

Notes

Use of this option is generally not recommended because it can cause considerable performance degradation. If your program is not sensitive to induction variable overflow or wrap-around, you should consider using **-qnostrict_induction** in conjunction with the **-O2** optimization option.

Related References

“Compiler Command Line Options” on page 51

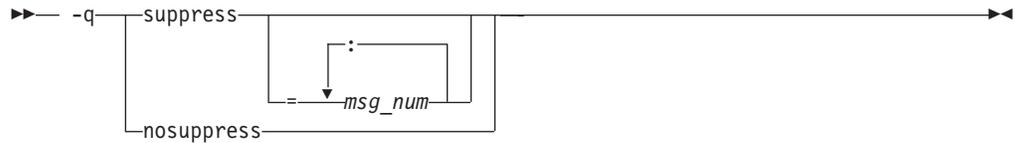
“O, optimize” on page 194

suppress

Purpose

Prevents the specified compiler or driver informational or warning messages from being displayed or added to the listings.

Syntax



Notes

This option suppresses compiler messages only, and has no effect on linker or operating system messages.

To suppress IPA messages, enter **-qsuppress** before **-qipa** on the command line.

Compiler messages that cause compilation to stop, such as (S) and (U) level messages, cannot be suppressed.

The **-qnosuppress** compiler option cancels previous settings of **-qsuppress**.

Example

If your program normally results in the following output:

```
"myprogram.c", line 1.1:1506-224 (I) Incorrect #pragma ignored
```

you can suppress the message by compiling with:

```
xlc myprogram.c -qsuppress
```

Related References

"Compiler Command Line Options" on page 51

"halt" on page 132

"ipa" on page 151

symtab

Purpose

Settings for this option determine what information appears in the symbol table.

Syntax

► — -q—symtab—=—┌unref┐
 └static┘

where:

unref Specifies that all **typedef** declarations, **struct**, **union**, and **enum** type definitions are included for processing by the IBM Distributed Debugger.

Use this option with the **-g** option to produce additional debugging information for use with the debugger.

When you specify the **-g** option, debugging information is included in the object file. To minimize the size of object and executable files, the compiler only includes information for symbols that are referenced. Debugging information is not produced for unreferenced arrays, pointers, or file-scope variables unless **-qsymtab=unref** is specified.

Using **-qsymtab=unref** may make your object and executable files larger.

static Adds user-defined, nonexternal names that have a persistent storage class, such as initialized and uninitialized static variables, to the name list (the symbol table of **xcoff** objects).

The default is to not add static variables to the symbol table.

Examples

To compile `myprogram.c` so that static symbols are added to the symbol table, enter:

```
xlc myprogram.c -qsymtab=static
```

To include all symbols in `myprogram.c` in the symbols table for use with a debugger, enter:

```
xlc myprogram.c -g -qsymtab=unref
```

Related References

“Compiler Command Line Options” on page 51

“g” on page 130

syntaxonly

Purpose

Causes the compiler to perform syntax checking without generating an object file.

Syntax

►— -q—syntaxonly—►

Notes

The **-P**, **-E**, and **-C** options override the **-qsyntaxonly** option, which in turn overrides the **-c** and **-o** options.

The **-qsyntaxonly** option suppresses only the generation of an object file. All other files (listings, etc) are still produced if their corresponding options are set.

Examples

To check the syntax of myprogram.c without generating an object file, enter:

```
xlc myprogram.c -qsyntaxonly
```

or

```
xlc myprogram.c -o testing -qsyntaxonly
```

Note that in the second example, the **-qsyntaxonly** option overrides the **-o** option so no object file is produced.

Related References

“Compiler Command Line Options” on page 51

“C” on page 83

“c” on page 84

“E” on page 105

“o” on page 198

“P” on page 199

tabsize

Purpose

Changes the length of tabs as perceived by the compiler.

Syntax

►► — `-q-tabsize= n` —►►

where n is the number of character spaces representing a tab in your source program.

Notes

This option only affects error messages that specify the column number at which an error occurred. For example, the compiler will consider tabs as having a width of one character if you specify `-qtabsize=1`. In this case, you can consider one character position (where each character and each tab equals one position, regardless of tab length) as being equivalent to one character column.

Related References

“Compiler Command Line Options” on page 51

threaded

Purpose

Indicates to the compiler that the program uses multiple threads. Always use this option when compiling or linking multi-threaded applications. This option ensures that all optimizations are thread-safe.

Syntax

►► -q {nothreaded | threaded} ◀◀

Default

The default is **-qthreaded** when compiling with **_r** invocation modes, and **-qnothreaded** when compiling with other invocation modes.

Notes

This option applies to both compile and linkage editor operations.

To maintain thread safety, a file compiled with the **-qthreaded** option, whether explicitly by option selection or implicitly by choice of **_r** compiler invocation mode, must also be linked with the **-qthreaded** option.

This option does not make code thread-safe, but it will ensure that code already thread-safe will remain so after compile and linking.

Related References

“Compiler Command Line Options” on page 51

“smp” on page 232

tocdata

Purpose

Marks data as local.

Syntax

►► -q notocdata
tocdata ◀◀

Notes

Local variables are statically bound to the functions that use them. **-qtocdata** instructs the compiler to assume that all variables are local.

If an imported variable is assumed to be local, performance may decrease. Imported variables are dynamically bound to a shared portion of a library. **-qnotocdata** instructs the compiler to assume that all variables are imported.

Conflicts among the data-marking options are resolved in the following manner:

Options that list variable names	The last explicit specification for a particular variable name is used.
Options that change the default	This form does not specify a name list. The last option specified is the default for variables not explicitly listed in the name-list form.

Related References

“Compiler Command Line Options” on page 51

tocmerge

Purpose

Enables TOC merging to reduce TOC pointer loads and improves the scheduling of external loads.

Syntax

►► — -q — notocmerge
tocmerge —————►►

Notes

If **-qtocmerge** specified, the compiler reads from the file specified in the **-bImportfile** linker option. If **-qtocmerge** is specified but no import filename is provided, the option is ignored and a warning message is issued.

Related References

“Compiler Command Line Options” on page 51

trigraph

Purpose

Instructs the compiler to recognize trigraph key combinations used to represent characters not found on some keyboards.

Syntax



Defaults

The default setting for **-qtrigraph** varies according to the command used to invoke the compiler. Defaults are:

- **-qnotrigraph** for **xl**, **xl_r**, **cc**, and **cc_r**
- **-qtrigraph** for **c89**, **c89_r**, **c99**, and **c99_r**.

Notes

A trigraph is a combination of three-key character combinations that let you produce a character that is not available on all keyboards.

The trigraph key combinations are:

Key Combination	Character Produced
??=	#
??([
??)]
??/	\
??'	^
??<	{
??>	}
??!	
??-	~

The default **-qtrigraph** setting can be overridden by explicitly setting the **-q[no]trigraph** option on the command line.

An explicit **-q[no]trigraph** specification on the command line takes precedence over the **-q[no]trigraph** setting normally associated with a given **-q[lang]vl** compiler option, regardless of where the **-q[no]trigraph** specification appears on the command line.

Example

To disable trigraph character sequences when compiling your program, enter:

```
xl c myprogram.c -qnotrigraph
```

Related References

“Compiler Command Line Options” on page 51

“digraph” on page 101

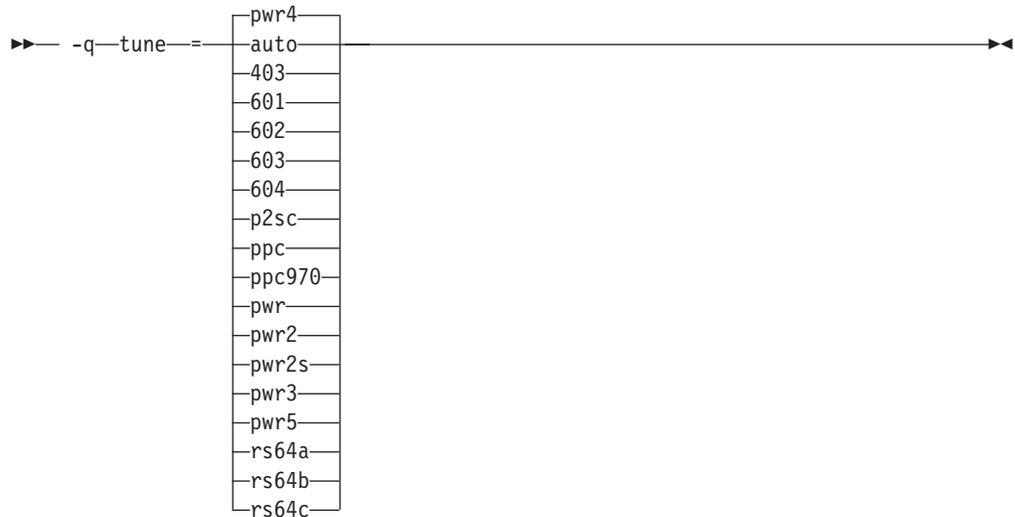
“langlvl” on page 165

tune

Purpose

Specifies the architecture system for which the executable program is optimized.

Syntax



where architecture suboptions are:

403	Produces object code optimized for the PowerPC 403 processor.
601	Produces object code optimized for the PowerPC 601 processor.
602	Produces object code optimized for the PowerPC 602 processor.
603	Produces object code optimized for the PowerPC 603 processor.
604	Produces object code optimized for the PowerPC 604 processor.
auto	Produces object code optimized for the platform on which it is compiled.
p2sc	Produces object code optimized for the PowerPC P2SC processor.
pwr	Produces object code optimized for the POWER hardware platforms.
pwr2	Produces object code optimized for the POWER2 hardware platforms.
pwr2s	Produces object code optimized for the POWER2 hardware platforms, avoiding certain quadruple-precision instructions that would slow program performance.
pwr3	Produces object code optimized for the POWER3 hardware platforms.
pwr4	Produces object code optimized for the POWER4 hardware platforms.
pwr5	Produces object code optimized for the POWER5 hardware platforms.
pwr x	Produces object code optimized for the POWER2 hardware platforms (same as <code>-qtune=pwr2</code>).
rs64a	Produces object code optimized for the RS64I processor.
rs64b	Produces object code optimized for the RS64II processor.
rs64c	Produces object code optimized for the RS64III processor.
ppc970	Produces object code optimized for the PowerPC 970 processor.

See also “#pragma options” on page 299.

Default

The default setting of the `-qtune` option depends on the setting of the `-qarch` option.

- If `-qtune` is specified without `-qarch`, the compiler uses `-qarch=com`.

- If **-qarch** is specified without **-qtune**, the compiler uses the default tuning option for the specified architecture. Listings will show only: TUNE=DEFAULT

Default **-qtune** settings for specific **-qarch** settings are described in “Acceptable Compiler Mode and Processor Architecture Combinations” on page 350.

Notes

You can use **-qtune=***suboption* with **-qarch=***suboption*.

- **-qarch=***suboption* specifies the architecture for which the instructions are to be generated, and,
- **-qtune=***suboption* specifies the target platform for which the code is optimized.

Example

To specify that the executable program testing compiled from myprogram.c is to be optimized for a POWER hardware platform, enter:

```
xlc -o testing myprogram.c -qtune=pwr
```

Related Tasks

“Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33

Related References

“Compiler Command Line Options” on page 51

“arch” on page 70

“Acceptable Compiler Mode and Processor Architecture Combinations” on page 350

U

Purpose

Undefines the identifier *name* defined by the compiler or by the **-Dname** option.

Syntax

▶▶ — **-U***name* —▶▶

Notes

The **-Uname** option is *not* equivalent to the **#undef** preprocessor directive. It *cannot* undefine names defined in the source by the **#define** preprocessor directive. It can only undefine names defined by the compiler or by the **-Dname** option.

The identifier name can also be undefined in your source program using the **#undef** preprocessor directive.

The **-Uname** option has a higher precedence than the **-Dname** option.

Example

Assume that your operating system defines the name **__unix**, but you do not want your compilation to enter code segments conditional on that name being defined. Compile `myprogram.c` so that the definition of the name **__unix** is nullified by entering:

```
xlc myprogram.c -U__unix
```

Related References

“Compiler Command Line Options” on page 51

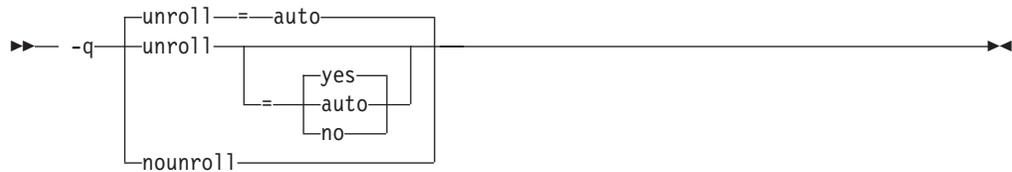
“D” on page 96

unroll

Purpose

Unrolls inner loops in the program. This can help improve program performance.

Syntax



where:

<code>-qunroll=auto</code>	Leaves the decision to unroll loops to the compiler. This is the compiler default.
<code>-qunroll</code> or <code>-qunroll=yes</code>	Suggests to the compiler that it unroll loops.
<code>-qnounroll</code> or <code>-qunroll=no</code>	Instructs the compiler to not unroll loops.

See also “`#pragma unroll`” on page 315 and “`#pragma options`” on page 299.

Notes

The compiler default for this option, unless explicitly specified otherwise on the command line, is **`-qunroll=auto`**.

Specifying **`-qunroll`** without any suboptions is equivalent to specifying **`-qunroll=yes`**.

When **`-qunroll`**, **`-qunroll=yes`**, or **`-qunroll=auto`** is specified, the bodies of inner loops will be unrolled, or duplicated, by the optimizer. The optimizer determines and applies the best unrolling factor for each loop. In some cases, the loop control may be modified to avoid unnecessary branching.

To see if the **`unroll`** option improves performance of a particular application, you should first compile the program with usual options, then run it with a representative workload. You should then recompile with command line **`-qunroll`** option and/or the **`unroll`** pragmas enabled, then rerun the program under the same conditions to see if performance improves.

You can use the **`#pragma unroll`** directive to gain more control over unrolling. Setting this pragma overrides the **`-qunroll`** compiler option setting.

Examples

1. In the following examples, unrolling is disabled:

```
xlc -qnounroll file.c
```

```
xlc -qunroll=no file.c
```

2. In the following examples, unrolling is enabled:

```
xlc -qunroll file.c
```

```
xlc -qunroll=yes file.c
```

```
xlc -qunroll=auto file.c
```

3. See “#pragma unroll” on page 315 for examples of how program code is unrolled by the compiler.

Related References

“Compiler Command Line Options” on page 51

“#pragma options” on page 299

“#pragma unroll” on page 315

unwind

Purpose

Informs the compiler that the stack can be unwound while a routine in the compilation is active.

Syntax

►► -q unwind
noundwind ◀◀

Notes

Specifying **-qnoundwind** can improve optimization of non-volatile register saves and restores.

Related References

“Compiler Command Line Options” on page 51

upconv

Purpose

Preserves the **unsigned** specification when performing integral promotions.

Syntax

►► -q noupconv
upconv ◀◀

See also “#pragma options” on page 299.

Notes

The **-qupconv** option promotes any **unsigned** type smaller than an **int** to an **unsigned int** instead of to an **int**.

Unsignedness preservation is provided for compatibility with older dialects of C. The ANSI C standard requires value preservation as opposed to unsignedness preservation.

Default

The default is **-qnoupconv**, except when **-qlanglvl=extc89**, in which case the default is **-qupconv**. The compiler does not preserve the **unsigned** specification.

The default compiler action is for integral promotions to convert a **char**, **short int**, **int** bit field or their **signed** or **unsigned** types, or an **enumeration** type to an **int**. Otherwise, the type is converted to an **unsigned int**.

Example

To compile myprogram.c so that all **unsigned** types smaller than **int** are converted to **unsigned int**, enter:

```
xlc myprogram.c -qupconv
```

The following short listing demonstrates the effect of **-qupconv**:

```
#include <stdio.h>
int main(void) {
    unsigned char zero = 0;
    if (-1 < zero)
        printf("Value-preserving rules in effect\n");
    else
        printf("Unsignedness-preserving rules in effect\n");
    return 0;
}
```

Related References

“Compiler Command Line Options” on page 51

“langlvl” on page 165

utf

Purpose

Enables recognition of UTF literal syntax.

Syntax

►► -q noutf
utf _____ ◀◀

Notes

The compiler uses **iconv** to convert the source file to Unicode. If the source file cannot be converted, the compiler will ignore the **-qutf** option and issue a warning.

Related References

“Compiler Command Line Options” on page 51

V

Purpose

Instructs the compiler to report information on the progress of the compilation, names the programs being invoked within the compiler and the options being specified to each program. Information is displayed in a space-separated list.

Syntax

▶▶ — -V —————▶▶

Notes

The `-V` option is overridden by the `-#` option.

Example

To compile `myprogram.c` so you can watch the progress of the compilation and see messages that describe the progress of the compilation, the programs being invoked, and the options being specified, enter:

```
xlc myprogram.c -V
```

Related References

“Compiler Command Line Options” on page 51

“# (pound sign)” on page 59

“v” on page 262

V

Purpose

Instructs the compiler to report information on the progress of the compilation, names the programs being invoked within the compiler and the options being specified to each program. Information is displayed in a comma-separated list.

Syntax

▶▶ -v —————▶▶

Notes

The `-v` option is overridden by the `-#` option.

Example

To compile `myprogram.c` so you can watch the progress of the compilation and see messages that describe the progress of the compilation, the programs being invoked, and the options being specified, enter:

```
xlc myprogram.c -v
```

Related References

“Compiler Command Line Options” on page 51

“# (pound sign)” on page 59

“V” on page 261

W

Purpose

Passes the listed option to a designated compiler program.

Syntax



where programs are:

program	Description
a	Assembler
b	Compiler back end
c	Compiler front end
I	Interprocedural Analysis tool - compile phase
L	Interprocedural Analysis tool - link phase
l	linkage editor
p	compiler preprocessor

Notes

When used in the configuration file, the **-W** option accepts the escape sequence backslash comma (\,) to represent a comma in the parameter string.

Examples

1. To compile myprogram.c so that the *option* **-pg** is passed to the linkage editor (**l**) and the assembler (**a**), enter:

```
xlc myprogram.c -Wl,-pg -Wa,-pg
```

2. In a configuration file, use the \, sequence to represent the comma (,).

```
-Wl\,-pg,-Wa\,-pg
```

Related References

“Compiler Command Line Options” on page 51

W

Purpose

Requests that warnings and lower-level messages be suppressed. Specifying this option is equivalent to specifying **-qflag=e:e**.

Syntax

► — -w — ◄

Notes

Informational and warning messages that supply additional information to a severe error are not disabled by this option. For example, a severe error caused by problems with overload resolution will also produce information messages. These informational messages are not disabled with **-w** option:

```
void func(int a){}
void func(int a, int b){}
int main(void)
{
  func(1,2,3);
  return 0;
}
```

```
"x.cpp", line 6.4: 1540-0218 (S) The call does not match any parameter list for "func".
"x.cpp", line 1.6: 1540-1283 (I) "func(int)" is not a viable candidate.
"x.cpp", line 6.4: 1540-0215 (I) The wrong number of arguments have been specified for "func(int)".
"x.cpp", line 2.6: 1540-1283 (I) "func(int, int)" is not a viable candidate.
"x.cpp", line 6.4: 1540-0215 (I) The wrong number of arguments have been specified for "func(int, int)".
```

Example

To compile myprogram.c so that no warning messages are displayed, enter:

```
xlc myprogram.c -w
```

Related References

“Compiler Command Line Options” on page 51

“flag” on page 117

weaksymbol

Purpose

Instructs the compiler to generate weak symbols.

Syntax

►► -q noweaksymbol
weaksymbol ◀◀

Notes

When **-qweaksymbol** is in effect, the compiler generates weak symbols for the following:

- Inline functions with external linkage.
- Identifiers specified as weak with **#pragma weak** or functions specified as weak with **__attribute__((weak))**.

Related References

“Compiler Command Line Options” on page 51

“#pragma weak” on page 319

xcall

Purpose

Generates code to treat static routines within a compilation unit as if they were external routines.

Syntax

►► — -q — $\begin{array}{l} \text{noxcall} \\ \text{xcall} \end{array}$ —————►►

Notes

-qxcall generates slower code than -qnoxcall.

Example

To compile myprogram.c so that all static routines are compiled as external routines, enter:

```
xlc myprogram.c -qxcall
```

Related References

“Compiler Command Line Options” on page 51

xref

Purpose

Produces a compiler listing that includes a cross-reference listing of all identifiers.

Syntax



where:

noxref	Do not report identifiers in the program.
xref	Reports only those identifiers that are used.
xref=full	Reports all identifiers in the program.

See also “#pragma options” on page 299.

Notes

The **-qnoprint** option overrides this option.

Any function defined with the **#pragma mc_func** *function_name* directive is listed as being defined on the line of the **#pragma** directive.

Example

To compile `myprogram.c` and produce a cross-reference listing of all identifiers whether they are used or not, enter:

```
xlc myprogram.c -qxref=full -qattr
```

A typical cross-reference listing has the form:

Identifier name	Description of the item
<code>xy</code>	<code>auto int in function adder</code>
	<code>0-59Y 0-36.12Z 0-48.12Z</code>
	Function invocation
	Column number
	Line number
	File
	Function definition

Related References

“Compiler Command Line Options” on page 51

“attr” on page 77

“print” on page 210

“#pragma mc_func” on page 297

“#pragma options” on page 299

y

Purpose

Specifies the compile-time rounding mode of constant floating-point expressions.

Syntax



where suboptions are:

n	Round to the nearest representable number. This is the default.
m	Round toward minus infinity.
p	Round toward plus infinity.
z	Round toward zero.

Example

To compile `myprogram.c` so that constant floating-point expressions are rounded toward zero at compile time, enter:

```
xlc myprogram.c -yz
```

Related References

“Compiler Command Line Options” on page 51

“`rndflt`” on page 219

Z

Purpose

This linker option specifies a prefix for the library search path.

Syntax

▶— *-Z—string—* ▶

Notes

This linker option is useful when developing a new version of a library. Usually you use it to build on one level of an operating system but run on a different level, so that you can search a different path on the development platform than on the target platform. This is possible because the prefix is not stored in the executable.

If you use this option more than once, the strings are appended to each other in the order specified and then they are added to the beginning of the library search paths.

Related References

“Compiler Command Line Options” on page 51

General Purpose Pragmas

The pragmas listed below are available for general programming use.

Language Application	#pragma	Description
	#pragma align	Aligns data items within structures.
	#pragma alloca	Provides an inline version of the function alloca(size_t size) .
	#pragma block_loop	Instructs the compiler to create a blocking loop for a specific loop in a loop nest.
	#pragma chars	Sets the sign type of character data.
	#pragma comment	Places a comment into the object file.
	#pragma disjoint	Lists the identifiers that are not aliased to each other within the scope of their use.
	#pragma enum	Specifies the size of enum variables that follow.
	#pragma execution_frequency	Lets you mark program source code that you expect will be either very frequently or very infrequently executed.
	#pragma ibm snapshot	Sets a debugging breakpoint at the point of the pragma, and defines a list of variables to examine when program execution reaches that point.
	#pragma info	Controls the diagnostic messages generated by the info(...) compiler options.
	#pragma isolated_call	Marks a function that does not have or rely on side effects, other than those implied by its parameters.
	#pragma langlvl	Selects the C language level for compilation.
	#pragma leaves	Takes a function name and specifies that the function never returns to the instruction after the function call.
	#pragma loopid	Marks a block with a scope-unique identifier.
	#pragma map	Tells the compiler that all references to an identifier are to be converted to a new name.
	#pragma mc_func	Lets you define a function containing a short sequence of machine instructions.
	#pragma options	Specifies options to the compiler in your source program.
	#pragma option_override	Specifies alternate optimization options for specific functions.
	#pragma pack	Modifies the current alignment mode for members of structures that follow this pragma.
	#pragma reachable	Declares that the point after the call to a routine marked reachable can be the target of a branch from some unknown location.
	#pragma reg_killed_by	Specifies those registers which value will be corrupted by the specified function. It must be used together with #pragma mc_func .

Language Application	#pragma	Description
	#pragma stream_unroll	Breaks a stream contained in a loop into multiple streams.
	#pragma strings	Sets storage type for strings.
	#pragma unroll	Unrolls innermost and outermost loops in the program. This can help improve program performance.
	#pragma unrollandfuse	Instructs the compiler to attempt an unroll and fuse operation on nested for loops. This can help improve program performance.
	#pragma weak	Prevents the link editor from issuing error messages if it does not find a definition for a symbol, or if it encounters a symbol multiply-defined during linking.

Related Concepts

“Program Parallelization” on page 11

Related Tasks

“Specify Compiler Options in Your Program Source Files” on page 31

“Control Parallel Processing with Pragmas” on page 39

Related References

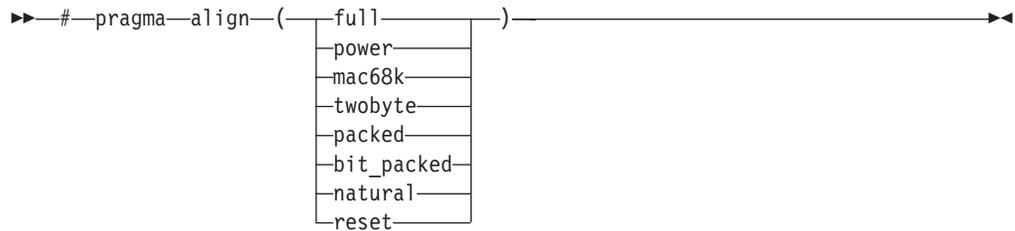
“Pragmas to Control Parallel Processing” on page 321

#pragma align

Description

The **#pragma align** directive specifies how the compiler should align data items within structures.

Syntax



See also “#pragma options” on page 299.

Notes

The **#pragma align=suboption** directive overrides the **-qalign** compiler option setting for a specified section of program source code.

The compiler stacks alignment directives, so you can go back to using a previous alignment directive without knowing what it is by specifying the **#pragma align(reset)** directive.

For example, you can use this option if you have a class declaration within an include file and you do not want the alignment rule specified for the class to apply to the file in which the class is included. You can code **#pragma align(reset)** in a source file to change the alignment option to what it was before the last alignment option was specified. If no previous alignment rule appears in the file, the alignment rule specified in the invocation command is used.

Specifying **#pragma align** has the same effect as specifying **#pragma options align** in your source file. For more information and examples of **#pragma align** and **#pragma options align** usage, see “align” on page 64.

Related References

“General Purpose Pragmas” on page 271

“align” on page 64

See also *Aligning data in aggregates* in the *Programming Guide*.

#pragma alloca

Description

The **#pragma alloca** directive specifies that the compiler should provide an inline version of the function **alloca(size_tsize)**. The function **alloca(size_tsize)** can be used to allocate space for an object. The amount of space allocated is determined by the value of *size*, which is measured in bytes. The allocated space is put on the stack.

Syntax

▶▶ #pragma alloca ◀◀

Notes

You must specify the **#pragma alloca** directive, or either of the **-qalloca** or **-ma** compiler options to have the compiler provide an inline version of **alloca**.

Once specified, **#pragma alloca** applies to the rest of the file and cannot be turned off. If a source file contains any functions that you want compiled without **#pragma alloca**, place these functions in a different file.

Related References

“General Purpose Pragmas” on page 271

“alloca” on page 68

“ma” on page 182

#pragma block_loop

Description

Instructs the compiler to create a blocking loop for a specific loop in a loop nest.

Syntax

```
▶▶ #pragma block_loop ( n, name_list ) ▶▶
```

where:

- n* Is an integer expression composed of variables and literals that specifies the block size of the iteration group.
- name_list* Is a list of zero or more unique, comma-separated identifiers created with the **#pragma loopid** directive. If you do not specify *name_list*, blocking occurs on the first **for** loop following the **#pragma block_loop** directive.

Notes

Blocking of a loop entails dividing a loop's iteration space into parts or blocks. An additional outer loop is then created, known as the blocking loop, which drives the original loop (called the blocked loop) for each part.

For loop blocking to occur, a **#pragma block_loop** directive must precede a **for** loop.

You must not combine **#pragma block_loop** with the **nounroll**, **unroll**, **nounrollandfuse**, **unrollandfuse**, or **stream_unroll** #pragma directives for the same **for** loop.

Important: Applying **#pragma block_loop** to a loop with dependencies or with alternate entry or exit points may produce unexpected results.

Examples

1. Simple blocking loop with no name specified.

```
#pragma block_loop(2)
for (...) { ... }
```

2. Blocking a named loop.

```
#pragma block_loop(5, myloop)
#pragma loopid(myloop)
for (...) { ... }
```

3. Blocking within a blocked loop.

```
#pragma block_loop(2, outerLoop)
for (i = 0; i < 100; i++) {
    #pragma loopid(outerLoop)
    #pragma block_loop(5)
    for (j = 0; j < 100; j++) {
        A[i][j] = i + j;
    }
}
```

Related References

"General Purpose Pragmas" on page 271

"#pragma loopid" on page 295

"#pragma unroll" on page 315
"#pragma unrollandfuse" on page 317
"#pragma stream_unroll" on page 312
"unroll" on page 256

#pragma chars

Description

The **#pragma chars** directive sets the sign type of char objects to be either **signed** or **unsigned**.

Syntax

▶▶ #pragma chars ($\left. \begin{array}{l} \text{unsigned} \\ \text{signed} \end{array} \right\}$) ▶▶

Notes

In order to have effect, this pragma must appear before any source statements.

Once specified, the pragma applies to the entire file and cannot be turned off. If a source file contains any functions that you want to be compiled without **#pragma chars**, place these functions in a different file. If the pragma is specified more than once in the source file, the first one will take precedence.

Note: The default character type behaves like an unsigned char.

Related References

“General Purpose Pragmas” on page 271

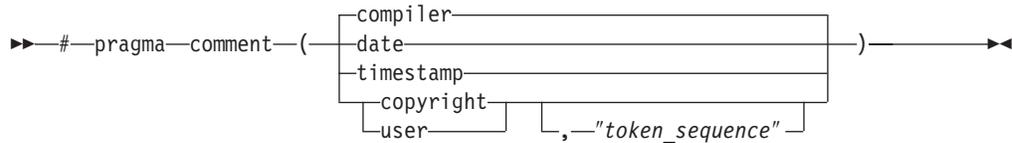
“chars” on page 88

#pragma comment

Description

The **#pragma comment** directive places a comment string into the target or object file.

Syntax



where suboptions do the following:

<code>compiler</code>	The name and version of the compiler is appended to the end of the generated object module.
<code>date</code>	The date and time of compilation is appended to the end of the generated object module.
<code>timestamp</code>	The date and time of the last modification of the source is appended to the end of the generated object module.
<code>copyright</code>	The text specified by the <i>token_sequence</i> is placed by the compiler into the generated object module and is loaded into memory when the program is run.
<code>user</code>	The text specified by the <i>token_sequence</i> is placed by the compiler into the generated object but is <i>not</i> loaded into memory when the program is run.

Example

Assume that following program code is compiled to produce output file **a.out**:

```
#pragma comment(date)
#pragma comment(compiler)
#pragma comment(timestamp)
#pragma comment(copyright,"My copyright")

int main() {
    return 0;
}
```

You can use the operating system **strings** command to look for these and other strings in an object or binary file. Issuing the command:

```
strings a.out
```

will cause the comment information embedded in **a.out** to be displayed, along with any other strings that may be found in **a.out**. For example, assuming the program code shown above:

```
Mon Mar 1 10:28:09 2004
XL C for AIX Compiler Version 7.0
Mon Mar 1 10:28:13 2004
My copyright
```

Related References

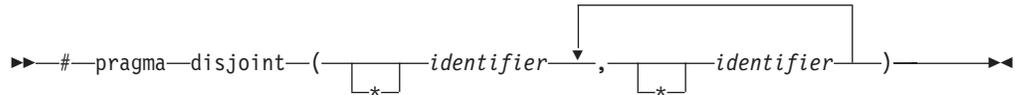
“General Purpose Pragmas” on page 271

#pragma disjoint

Description

The **#pragma disjoint** directive lists the identifiers that are not aliased to each other within the scope of their use.

Syntax



Notes

The directive informs the compiler that none of the identifiers listed shares the same physical storage, which provides more opportunity for optimizations. If any identifiers actually share physical storage, the pragma may cause the program to give incorrect results.

An identifier in the directive must be visible at the point in the program where the pragma appears. The identifiers in the disjoint name list cannot refer to any of the following:

- a member of a structure, or union
- a structure, union, or enumeration tag
- an enumeration constant
- a typedef name
- a label

This pragma can be disabled with the **-qignprag** compiler option.

Example

```
int a, b, *ptr_a, *ptr_b;
#pragma disjoint(*ptr_a, b) // *ptr_a never points to b
#pragma disjoint(*ptr_b, a) // *ptr_b never points to a
void one_function()
{
    b = 6;
    *ptr_a = 7; // Assignment does not alter the value of b
    another_function(b); // Argument "b" has the value 6
}
```

Because external pointer *ptr_a* does not share storage with and never points to the external variable *b*, the assignment of 7 to the object that *ptr_a* points to will not change the value of *b*. Likewise, external pointer *ptr_b* does not share storage with and never points to the external variable *a*. The compiler can assume that the argument of *another_function* has the value 6 and will not reload the variable from memory.

Related References

“General Purpose Pragmas” on page 271

“ignprag” on page 141

“alias” on page 62

#pragma enum

Description

The **#pragma enum** directive specifies the size of enum variables that follow. The size at the left brace of a declaration is the one that affects that declaration, regardless of whether further enum directives occur within the declaration. This pragma pushes a value on a stack each time it is used, with a reset option available to return to the previously pushed value.

Syntax

```
▶▶ #pragma enum ( suboption )
```

where *suboption* is any of the following:

1	The enumeration type is one byte in length, of type char if the range of enumeration values falls within the limits of signed char , and unsigned char otherwise.
2	The enumeration type is two bytes in length, of type short if the range of enumeration values falls within the limits of signed short , and unsigned short otherwise.
4	The enumeration type is four bytes in length, of type int if the range of enumeration values falls within the limits of signed int , and unsigned int otherwise.
8	The enumeration type is eight bytes in length. In 32-bit compilation mode, the enumeration is of type long long if the range of enumeration values falls within the limits of signed long long , and unsigned long long otherwise. In 64-bit compilation mode, the enumeration is of type long if the range of enumeration values falls within the limits of signed long , and unsigned long otherwise.
int	Same as #pragma enum=4 .
intlong	Specifies that enumeration will occupy 8 bytes of storage if the range of values in the enumeration exceeds the limit for int . See the description for #pragma enum=8 . If the range of values in the enumeration does not exceed the limit for int , the enumeration will occupy 4 bytes of storage and is represented by int .
small	The enumeration type is the smallest integral type that can contain all variables. If an 8-byte enum results, the actual enumeration type used is dependent on compilation mode. See the description for #pragma enum=8 .
pop	This suboption resets the enum size setting to its previous #pragma enum setting. If there is no previous setting, the command line setting for -qenum is used.
reset	Same as pop . This option is provided for backwards compatibility.

Notes

Popping on an empty stack generates a warning message and the enum value remains unchanged.

The **#pragma enum** directive overrides the **-qenum** compiler option.

For each **#pragma enum** directive that you put in a source file, it is good practice to have a corresponding **#pragma enum=reset** before the end of that file. This is the only way to prevent one file from potentially changing the **enum** setting of another file that **#includes** it.

The **#pragma options enum=** directive can be used instead of **#pragma enum**. The two pragmas are interchangeable.

A **-qenum=reset** option corresponding to the **#pragma enum=reset** directive does not exist. Attempting to use **-qenum=reset** generates a warning message and the option is ignored.

Examples

1. Usage of the **pop** and **reset** suboptions are shown in the following code segment.

```
#pragma enum(1)
#pragma enum(2)
#pragma enum(4)
#pragma enum(pop) /* will reset enum size to 2 */
#pragma enum(reset) /* will reset enum size to 1 */
#pragma enum(pop) /* will reset enum size to the -qenum setting,
                  assuming -qenum was specified on the command
                  line. If -qenum was not specified on the
                  command line, the compiler default is used. */
```

2. One typical use for the **reset** suboption is to reset the enumeration size set at the end of an include file that specifies an enumeration storage different from the default in the main file. For example, the following include file, `small_enum.h`, declares various minimum-sized enumerations, then resets the specification at the end of the include file to the last value on the option stack:

```
#ifndef small_enum_h
#define small_enum_h 1
/*
 * File small_enum.h
 * This enum must fit within an unsigned char type
 */

#pragma options enum=small
enum e_tag {a, b=255};
enum e_tag u_char_e_var; /* occupies 1 byte of storage */

/* Reset the enumeration size to whatever it was before */
#pragma options enum=reset
#endif
```

The following source file, `int_file.c`, includes `small_enum.h`:

```
/*
 * File int_file.c
 * Defines 4 byte enums
 */
#pragma options enum=int
enum testing {ONE, TWO, THREE};
enum testing test_enum;

/* various minimum-sized enums are declared */
#include "small_enum.h"

/* return to int-sized enums. small_enum.h has reset the
 * enum size
 */
enum sushi {CALIF_ROLL, SALMON_ROLL, TUNA, SQUID, UNI};
enum sushi first_order = UNI;
```

The enumerations **test_enum** and **first_order** both occupy 4 bytes of storage and are of type **int**. The variable **u_char_e_var** defined in `small_enum.h` occupies 1 byte of storage and is represented by an **unsigned char** data type.

3. If the following C fragment is compiled with the **enum=small** option:

```
enum e_tag {a, b, c} e_var;
```

the range of enum constants is 0 through 2. This range falls within all of the ranges described in the table above. Based on priority, the compiler uses predefined type **unsigned char**.

4. If the following C code fragment is compiled with the **enum=small** option:

```
enum e_tag {a=-129, b, c} e_var;
```

the range of enum constants is -129 through -127. This range only falls within the ranges of **short (signed short)** and **int (signed int)**. Because **short (signed short)** is smaller, it will be used to represent the **enum**.

5. If you compile a file `myprogram.c` using the command:

```
xlc myprogram.c -qenum=small
```

assuming file `myprogram.c` does not contain **#pragma options=int** statements, all **enum** variables within your source file will occupy the minimum amount of storage.

6. If you compile a file `yourfile.c` that contains the following lines:

```
enum testing {ONE, TWO, THREE};
enum testing test_enum;

#pragma options enum=small
enum sushi {CALIF_ROLL, SALMON_ROLL, TUNA, SQUID, UNI};
enum sushi first_order = UNI;

#pragma options enum=int
enum music {ROCK, JAZZ, NEW_WAVE, CLASSICAL};
enum music listening_type;
```

using the command:

```
xlc yourfile.c
```

only the enum variable **first_order** will be minimum-sized (that is, enum variable **first_order** will only occupy 1 byte of storage). The other two enum variables **test_enum** and **listening_type** will be of type **int** and occupy 4 bytes of storage.

The following examples show invalid enumerations or usage of **#pragma enum**:

- You cannot change the storage allocation of an enum using a **#pragma enum** within the declaration of an enum. The following code segment generates a warning and the second occurrence of the **enum** option is ignored:

```
#pragma enum=small
enum e_tag {
    a,
    b,
    #pragma enum=int /* error: cannot be within a declaration */
    c
} e_var;
#pragma enum=reset /* second reset isn't required */
```

- The range of **enum** constants must fall within the range permitted by your chosen **enum** suboption. The following code segments contain errors:

```
#pragma enum(small)
enum e_tag {
    a = -1,
    b = 0x8000000000000000, /* error: enum value > MAX long long */
} e_var;
#pragma enum(reset)
```

- The **enum** constant range does not fit within the range of an **unsigned long long** in 32-bit mode, or **unsigned long** in 64-bit mode.

```
#pragma enum(small)
enum e_tag {
    a=0,
    b=0xFFFFFFFFFFFFFFFFLL + 1LL, /* error, larger than max permitted range */
} e_var;
#pragma enum(reset)
```

Related References

“General Purpose Pragmas” on page 271

“enum” on page 108

“#pragma options” on page 299

#pragma execution_frequency

Description

The `#pragma execution_frequency` directive lets you mark program source code that you expect will be either very frequently or very infrequently executed.

Syntax

```
▶ #pragma execution_frequency ( [very_low] | [very_high] ) ▶
```

Notes

Use this pragma to mark program source code that you expect will be executed very frequently or very infrequently. The pragma must be placed within block scope, and acts on the closest point of branching.

The pragma is used as a hint to the optimizer. If optimization is not selected, this pragma has no effect.

Examples

1. This pragma is used in an `if` statement block to mark code that is executed infrequently.

```
int *array = (int *) malloc(10000);

if (array == NULL) {
    /* Block A */
    #pragma execution_frequency(very_low)
    error();
}
```

The code block "Block B" would be marked as infrequently executed and "Block C" is likely to be chosen during branching.

```
if (Foo > 0) {
    #pragma execution_frequency(very_low)
    /* Block B */
    doSomething();
} else {
    /* Block C */
    doAnotherThing();
}
```

2. This pragma is used in a `switch` statement block to mark code that is executed frequently.

```
while (counter > 0) {
    #pragma execution_frequency(very_high)
    doSomething();
} /* This loop is very likely to be executed. */

switch (a) {
    case 1:
        doOneThing();
        break;
    case 2:
        #pragma execution_frequency(very_high)
        doTwoThings();
        break;
    default:
        doNothing();
} /* The second case is frequently chosen. */
```

3. This pragma cannot be used at file scope. It can be placed anywhere within a block scope and it affects the closest branching.

```
int a;
#pragma execution_frequency(very_low)
int b;

int foo(boolean boo) {
    #pragma execution_frequency(very_low)
    char c;

    if (boo) {
        /* Block A */
        doSomething();
        {
            /* Block C */
            doSomethingAgain();
            #pragma execution_frequency(very_low)
            doAnotherThing();
        }
    } else {
        /* Block B */
        doNothing();
    }

    return 0;
}

#pragma execution_frequency(very_low)
```

The first and fourth pragmas are invalid, while the second and third are valid. However, only the third pragma has effect, and it affects whether program execution branches to Block A or Block B during the decision of "if (**boo**)". The second pragma is ignored by the compiler.

Related References

"General Purpose Pragmas" on page 271

#pragma ibm snapshot

Description

The **#pragma ibm snapshot** directive sets a debugging breakpoint at the point of the pragma, and defines a list of variables to examine when program execution reaches that point.

Syntax

```
▶▶ #pragma ibm snapshot ( variable_name ) ▶▶
```

where *variable_name* is a predefined or namespace scope type. Class, structure, or union members cannot be specified.

Notes

Variables specified in **#pragma ibm snapshot** can be observed in the debugger, but should not be modified. Modifying these variables in the debugger may result in unpredictable behavior.

Example

```
#pragma ibm snapshot(a, b, c)
```

Related References

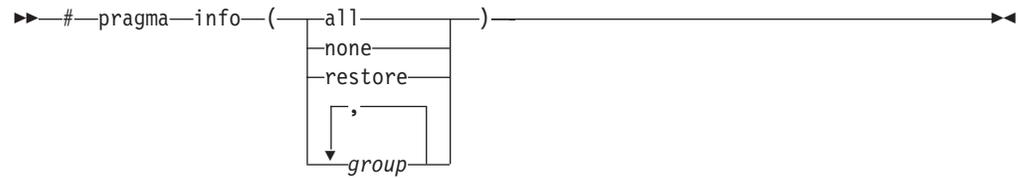
“General Purpose Pragmas” on page 271

#pragma info

Description

The **#pragma info** directive instructs the compiler to produce or suppress specific groups of compiler messages.

Syntax



where:

- `all` Turns on all diagnostic checking.
- `none` Turns off all diagnostic suboptions for specific portions of your program
- `restore` Restores the option that was in effect before the previous **#pragma info** directive.

group Generates or suppresses all messages associated with the specified diagnostic *group*. More than one *group* name in the following list can be specified.

<i>group</i>	Type of messages returned or suppressed
c99 noc99	C code that may behave differently between C89 and C99 language levels.
cmp nocmp	Possible redundancies in unsigned comparisons
cnd nocnd	Possible redundancies or problems in conditional expressions
cns nocns	Operations involving constants
cnv nocnv	Conversions
dc1 nodc1	Consistency of declarations
eff noeff	Statements and pragmas with no effect
enu noenu	Consistency of enum variables
ext noext	Unused external definitions
gen nogen	General diagnostic messages
gnr nognr	Generation of temporary variables
got nogot	Use of goto statements
ini noini	Possible problems with initialization
inl noinl	Functions not inlined
lan nolan	Language level effects
obs noobs	Obsolete features
ord noord	Unspecified order of evaluation
par nopar	Unused parameters
por nopor	Nonportable language constructs
ppc noppc	Possible problems with using the preprocessor
ppt noppt	Trace of preprocessor actions
pro nopro	Missing function prototypes
rea norea	Code that cannot be reached
ret noret	Consistency of return statements
trd notrd	Possible truncation or loss of data or precision
tru notru	Variable names truncated by the compiler
trx notrx	Hexadecimal floating point constants rounding
uni noui	Uninitialized variables
upg noupg	Generates messages describing new behaviors of the current compiler release as compared to the previous release.
use nouse	Unused auto and static variables
zea nozea	Zero-extent arrays.

Notes

You can use the **#pragma info** directive to temporarily override the current **-qinfo** compiler option settings specified on the command line, in the configuration file, or by earlier invocations of the **#pragma info** directive.

Example

For example, in the code segments below, the **#pragma info(eff, nouni)** directive preceding MyFunction1 instructs the compiler to generate messages identifying statements or pragmas with no effect, and to suppress messages identifying uninitialized variables. The **#pragma info(restore)** directive preceding MyFunction2 instructs the compiler to restore the message options that were in effect before the **#pragma info(eff, nouni)** directive was invoked.

```
#pragma info(eff, nouni)
int MyFunction1()
{
    .
    .
    .
}

#pragma info(restore)
int MyFunction2()
{
    .
    .
    .
}
```

Related References

“General Purpose Pragmas” on page 271

“info” on page 142

#pragma isolated_call

Description

The **#pragma isolated_call** directive marks a function that does not have or rely on side effects, other than those implied by its parameters.

Syntax

```
▶▶#—pragma—isolated_call—(—function—)—————▶▶
```

where *function* is a primary expression that can be an identifier, operator function, conversion function, or qualified name. An identifier must be of type function or a typedef of function. If the name refers to an overloaded function, all variants of that function are marked as isolated calls.

Notes

The **-qisolated_call** compiler option has the same effect as this pragma.

The pragma informs the compiler that the function listed does not have or rely on side effects, other than those implied by its parameters. Functions are considered to have or rely on side effects if they:

- Access a volatile object
- Modify an external object
- Modify a static object
- Modify a file
- Access a file that is modified by another process or thread
- Allocate a dynamic object, unless it is released before returning
- Release a dynamic object, unless it was allocated during the same invocation
- Change system state, such as rounding mode or exception handling
- Call a function that does any of the above

Essentially, any change in the state of the runtime environment is considered a side effect. Modifying function arguments passed by pointer or by reference is the only side effect that is allowed. Functions with other side effects can give incorrect results when listed in **#pragma isolated_call** directives.

Marking a function as **isolated_call** indicates to the optimizer that external and static variables cannot be changed by the called function and that pessimistic references to storage can be deleted from the calling function where appropriate. Instructions can be reordered with more freedom, resulting in fewer pipeline delays and faster execution in the processor. Multiple calls to the same function with identical parameters can be combined, calls can be deleted if their results are not needed, and the order of calls can be changed.

The function specified is permitted to examine non-volatile external objects and return a result that depends on the non-volatile state of the runtime environment. The function can also modify the storage pointed to by any pointer arguments passed to the function, that is, calls by reference. Do not specify a function that calls itself or relies on local static storage. Listing such functions in the **#pragma isolated_call** directive can give unpredictable results.

The `-qignprag` compiler option causes aliasing pragmas to be ignored. Use the `-qignprag` compiler option to debug applications containing the `#pragma isolated_call` directive.

Example

The following example shows the use of the `#pragma isolated_call` directive. Because the function `this_function` does not have side effects, a call to it will not change the value of the external variable `a`. The compiler can assume that the argument to `other_function` has the value 6 and will not reload the variable from memory.

```
int a;

// Assumed to have no side effects
int this_function(int);

#pragma isolated_call(this_function)
that_function()
{
    a = 6;
    // Call does not change the value of "a"
    this_function(7);

    // Argument "a" has the value 6
    other_function(a);
}
```

Related References

“General Purpose Pragmas” on page 271

“ignprag” on page 141

“isolated_call” on page 160

#pragma langlvl

Description

The `#pragma langlvl` directive selects the C language level for compilation.

Syntax

```
▶▶ #pragma langlvl ( language ) ▶▶
```

where values for *language* are described below.

For C programs, you can specify one of the following values for *language*:

<code>classic</code>	Allows the compilation of non-stdc89 programs, and conforms closely to the K&R level preprocessor. This language level is not supported by the AIX v5.1 system header files, such as <code>math.h</code> . If you must use the AIX v5.1 system header files, consider compiling your program to the stdc89 or extended language levels.
<code>extended</code>	Provides compatibility with the RT compiler and classic . This language level is based on C89.
<code>saa</code>	Compilation conforms to the current SAA C CPI language definition. This is currently SAA C Level 2.
<code>saa12</code>	Compilation conforms to the SAA C Level 2 CPI language definition, with some exceptions.
<code>stdc89</code>	Compilation conforms to the ANSI C89 standard, also known as ISO C90.
<code>stdc99</code>	Compilation conforms to the ISO C99 standard. Note: Not all operating system releases support the header files and runtime library required by C99.
<code>extc89</code>	Compilation conforms to the ANSI C89 standard, and accepts implementation-specific language extensions.
<code>extc99</code>	Compilation conforms to the ISO C99 standard, and accepts implementation-specific language extensions. Note: Not all operating system releases support the header files and runtime library required by C99.

Default

The default language level varies according to the command you use to invoke the compiler:

Invocation	Default language level
<code>xlc</code>	<code>extc89</code>
<code>cc</code>	<code>extended</code>
<code>c89</code>	<code>stdc89</code>
<code>c99</code>	<code>stdc99</code>

Notes

This pragma can be specified only once in a source file, and it must appear before any noncommentary statements in a source file.

The compiler uses predefined macros in the header files to make declarations and definitions available that define the specified language level.

This directive can dynamically alter preprocessor behavior. As a result, compiling with the `-E` compiler option may produce results different from those produced when not compiling with the `-E` option.

Related References

“General Purpose Pragmas” on page 271

“E” on page 105

“langlvl” on page 165

See also the *IBM C Language Extensions* section of the *C/C++ Language Reference*.

#pragma leaves

Description

The **#pragma leaves** directive takes a function name and specifies that the function never returns to the instruction after the call.

Syntax

Notes

This pragma tells the compiler that *function* never returns to the caller.

The advantage of the pragma is that it allows the compiler to ignore any code that exists after *function*, in turn, the optimizer can generate more efficient code. This pragma is commonly used for custom error-handling functions, in which programs can be terminated if a certain error is encountered. Some functions which also behave similarly are **exit**, **longjmp**, and **terminate**.

Example

```
#pragma leaves(handle_error_and_quit)
void test_value(int value)
{
    if (value == ERROR_VALUE)
    {
        handle_error_and_quit(value);
        TryAgain(); // optimizer ignores this because
                   // never returns to execute it
    }
}
```

Related References

“General Purpose Pragas” on page 271

#pragma loopid

Description

Marks a function block with a scope-unique identifier.

Syntax

```
▶▶ #pragma loopid (—name—) ▶▶
```

where *name* is an identifier that is unique within the function scoping unit.

Notes

The **#pragma loopid** directive must immediately precede a **#pragma block_loop** directive or **for** loop. The specified name can be used by **#pragma block_loop** to control transformations on that loop. It can also be used to provide information on loop transformations through the use of the **-qreport** compiler option.

You must not specify **#pragma loopid** more than once for a given loop.

Related References

“General Purpose Pragmas” on page 271

“report” on page 218

“unroll” on page 256

“#pragma block_loop” on page 275

“#pragma stream_unroll” on page 312

“#pragma unroll” on page 315

“#pragma unrollandfuse” on page 317

#pragma map

Description

The **#pragma map** directive tells the compiler that all references to an identifier are to be converted to “*name*”.

Syntax

```
▶ #pragma map ( identifier , “name” ) ▶
```

where:

<i>identifier</i>	A name of a data object with external linkage.
<i>name</i>	The external name that is to be bound to the given object, function, or operator.

Notes

You should not use **#pragma map** to map functions with built in linkage.

The directive can appear anywhere in the program. The identifiers appearing in the directive, including any type names used in the prototype argument list, are resolved as though the directive had appeared at file scope, independent of its actual point of occurrence.

Example

```
int funcname1()
{
    return 1;
}

#pragma map(func , "funcname1") /* maps func to funcname1 */

int main()
{
    return func(); /* no function prototype needed in C */
}
```

Related References

“General Purpose Pragmas” on page 271

#pragma mc_func

Description

The **#pragma mc_func** directive lets you define a function containing a short sequence of machine instructions.

Syntax

```
▶▶ #pragma mc_func function { instruction_seq } ▶▶
```

where:

<i>function</i>	Should specify a previously-defined function. If the function is not previously-defined, the compiler will treat the pragma as a function definition.
<i>instruction_seq</i>	Is a string containing a sequence of zero or more hexadecimal digits. The number of digits must comprise an integral multiple of 32 bits.

Notes

The **mc_func** pragma lets you embed a short sequence of machine instructions "inline" within your program source code. The pragma instructs the compiler to generate specified instructions in place rather than the usual linkage code. Using this pragma avoids performance penalties associated with making a call to an assembler-coded external function. This pragma is similar in function to the **asm** keyword found in this and other compilers.

The **mc_func** pragma defines a function and should appear in your program source only where functions are ordinarily defined. The function name defined by **#pragma mc_func** should be previously declared or prototyped.

The compiler passes parameters to the function in the same way as any other function. For example, in functions taking integer-type arguments, the first parameter is passed to GPR3, the second to GPR4, and so on. Values returned by the function will be in GPR3 for integer values, and FPR1 for float or double values. See **#pragma reg_killed_by** for a list of volatile registers available on your system.

Code generated from *instruction_seq* may use any and all volatile registers available on your system unless you use **#pragma reg_killed_by** to list a specific register set for use by the function.

Inlining options do not affect functions defined by **#pragma mc_func**. However, you may be able to improve runtime performance of such functions with **#pragma isolated_call**.

Example

In the following example, **#pragma mc_func** is used to define a function called **add_logical**. The function consists of machine instructions to add 2 ints with so-called *end-around carry*; that is, if a carry out results from the add then add the carry to the sum. This is frequently used in checksum computations.

The example also shows the use of `#pragma reg_killed_by` to list a specific set of volatile registers that can be altered by the function defined by `#pragma mc_func`.

```
int add_logical(int, int);
#pragma mc_func add_logical {"7c632014" "7c630194"}
        /* addc      r3 <- r3, r4      */
        /* addze     r3 <- r3, carry bit */

#pragma reg_killed_by add_logical gr3, xer
        /* only gpr3 and the xer are altered by this function */

main() {

    int i,j,k;

    i = 4;
    k = -4;
    j = add_logical(i,k);
    printf("\n\nresult = %d\n\n",j);
}
```

Related References

“General Purpose Pragmas” on page 271

“#pragma isolated_call” on page 290

“#pragma reg_killed_by” on page 310

“asm” on page 74

Valid settings for #pragma options <i>option_keyword</i>	Compiler option equivalent	Description
align= <i>option</i>	-qalign	Specifies what aggregate alignment modes the compiler uses for file compilation.
[no]ansialias	-qalias	Specifies whether type-based aliasing is to be used during optimization.
assert= <i>option</i>	-qalias	Requests the compiler to apply aliasing assertions to your compilation unit.
[no]attr attr=full	-qattr	Produces an attribute listing containing all names.
chars= <i>option</i>	-qchars See also #pragma chars	Instructs the compiler to treat all variables of type char as either signed or unsigned.
[no]check	-qcheck	Generates code which performs certain types of run-time checking.
[no]compact	-qcompact	Reduces code size, where possible, when used with optimization. Code size is reduced at the expense of execution speed.
[no]dbcs	-qmbcs, dbcs	Permits the use of DBCS characters in string literals and comments.
[no]dbxextra	-qdbxextra	Generates symbol table information for unreferenced variables.
[no]digraph	-qdigraph	Allows special digraph and keyword operators.
[no]dollar	-qdollar	Allows the \$ symbol to be used in the names of identifiers.
enum= <i>option</i>	-qenum See also #pragma enum	Specifies the amount of storage occupied by the enumerations.
[no]extchk	-qextchk	Performs external name type-checking and function call checking.
flag= <i>option</i>	-qflag	Specifies the minimum severity level of diagnostic messages to be reported.
float=[no] <i>option</i>	-qfloat	Specifies various floating point options to speed up or improve the accuracy of floating point operations.
[no]flttrap= <i>option</i>	-qflttrap	Generates extra instructions to detect and trap floating point exceptions.
[no]fold	-qfold	Specifies that constant floating point expressions are to be evaluated at compile time.

Valid settings for #pragma options <i>option_keyword</i>	Compiler option equivalent	Description
[no]fullpath	-qfullpath	Specifies the path information stored for files for dbx stabstrings.
[no]funcsect	-qfuncsect	Places instructions for each function in a separate cset.
halt	-qhalt	Stops compiler when errors of the specified severity detected.
[no]idirfirst	-qidirfirst	Specifies search order for user include files.
[no]ignerrno	-qignerrno	Allows the compiler to perform optimizations that assume errno is not modified by system calls.
[no]ignprag	-qignprag	Instructs the compiler to ignore certain pragma statements.
[no]info= <i>option</i>	-qinfo See also #pragma info	Produces informational messages.
initauto= <i>value</i>	-qinitauto	Initializes automatic storage to a specified hexadecimal byte value.
[no]inlglue	-qinlglue	Generates fast external linkage by inlining the pointer glue code necessary to make a call to an external function or a call through a function pointer.
isolated_call= <i>names</i>	-qisolated_call See also #pragma isolated_call	Specifies functions in the source file that have no side effects.
langlvl	-qlanglvl	Specifies different language levels. This directive can dynamically alter preprocessor behavior. As a result, compiling with the -E compiler option may produce results different from those produced when not compiling with the -E option.
[no]dbl128	-qdbl128, longdouble	Increases the size of long double type from 64 bits to 128 bits.
[no]libansi	-qlibansi	Assumes that all functions with the name of an ANSI C library function are in fact the system functions.
[no]list	-qlist	Produces a compiler listing that includes an object listing.
[no]longlong	-qlonglong	Allows long long types in your program.
[no]macpstr	-qmacpstr	Converts Pascal string literals into null-terminated strings where the first byte contains the length of the string.

Valid settings for #pragma options <i>option_keyword</i>	Compiler option equivalent	Description
[no]maf	-qmaf	Specifies whether floating-point multiply-add instructions are to be generated.
[no]maxmem= <i>number</i>	-qmaxmem	Instructs the compiler to halt compilation when a specified number of errors of specified or greater severity is reached.
[no]mbcs	-qmbcs, dbcs	Permits the use of DBCS characters in string literals and comments.
[no]optimize optimize= <i>number</i>	-O, optimize	Specifies the optimization level to apply to a section of program code. Valid settings for <i>number</i> are: <ul style="list-style-type: none"> • 0 - sets level 0 optimization • 2 - sets level 2 optimization • 3 - sets level 3 optimization If no value is specified for <i>number</i> , the compiler assumes level 2 optimization.
[no]proclcal, [no]procimported, [no]procunknown	-qproclcal, procimported, procunknown	Marks functions as local, imported, or unknown.
[no]proto	-qproto	Instructs the compiler to assume that all functions are prototyped.
[no]ro	-qro	Specifies the storage type for string literals.
[no]roconst	-qroconst	Specifies the storage location for constant values.
[no]rptr	-qrptr	Specifies the storage location for constant pointers.
[no]rrm	-qfloat=rrm	Prevents floating-point optimizations that are incompatible with run-time rounding to plus and minus infinity modes.
[no]showinc	-qshowinc	Used in conjunction with -qsource to add all include files to the source listing.
[no]source	-qsource	Produces a source listing.
spill= <i>number</i>	-qspill	Specifies the size of the register allocation spill area.
[no]srcmsg	-qsrcmsg	Adds the corresponding source code lines to the diagnostic messages in the stderr file.
[no]stdinc	-qstdinc	Specifies which files are included with #include <file_name> and #include "file_name" directives.

Valid settings for #pragma options <i>option_keyword</i>	Compiler option equivalent	Description
[no]strict	-qstrict	Turns off aggressive optimizations of the -O3 compiler option that have the potential to alter the semantics of your program.
thtable= <i>option</i>	-qthtable	Changes the length of tabs as perceived by the compiler.
tune= <i>option</i>	-qtune	Specifies the architecture for which the executable program is optimized.
[no]unroll unroll= <i>number</i>	-qunroll	Unrolls inner loops in the program by a specified factor.
[no]upconv	-qupconv	Preserves the unsigned specification when performing integral promotions.
[no]xref	-qxref	Produces a compiler listing that includes a cross-reference listing of all identifiers.

Related References

"General Purpose Pragmas" on page 271

"E" on page 105

#pragma option_override

Description

The **#pragma option_override** directive lets you specify alternate optimization options to apply to specific functions.

Syntax

```
▶▶ #pragma option_override ( fname " option " ) ▶▶
```

Valid settings and syntax for *option*, and their corresponding command line options, are shown below:

Settings and Syntax for #pragma option_override <i>option</i>	Command Line Option	Examples
opt(level,number)	-O, -O2, -O3, -O4, -O5	#pragma option_override (fname, "opt(level, 3)")
opt(registerSpillSize,num)	-qspill=num	#pragma option_override (fname, "opt(registerSpillSize,512)")
opt(size[,yes])	-qcompact	#pragma option_override (fname, "opt(size)") #pragma option_override (fname, "opt(size,yes)")
opt(size,no)	-qnocompact	#pragma option_override (fname, "opt(size,no)")
opt(strict)	-qstrict	#pragma option_override (fname, "opt(strict)")
opt(strict,no)	-qnostrict	#pragma option_override (fname, "opt(strict,no)")

Notes

By default, optimization options specified on the command line apply to the entire source program. However, certain types of runtime errors may occur only when optimization is turned on. This pragma lets you override command line optimization settings for specific functions (*fname*) in your program, which may be useful in identifying and correcting programming errors in those functions.

Per-function optimizations have effect only if optimization is already enabled by compilation option. You can request per-function optimizations at a level less than that applied to the rest of the program being compiled. Selecting options through this pragma affects only the specific optimization option selected, and does not affect the implied settings of related options.

Options are specified in double quotes, so they are not subject to macro expansion. The option specified within quotes must comply with the syntax of the build option.

This pragma affects only functions defined in your compilation unit and can appear anywhere in the compilation unit, for example:

- before or after a compilation unit
- before or after the function definition
- before or after the function declaration
- before or after a function has been referenced
- inside or outside a function definition.

Related References

“General Purpose Pragmas” on page 271

“compact” on page 91

“O, optimize” on page 194

“spill” on page 236

“strict” on page 241

Examples

1. In the code shown below, the structure s_t2 will have its members packed to 1-byte, but structure s_t1 will not be affected. This is because the declaration for s_t1 began before the pragma directive. However, s_t2 is affected because its declaration began after the pragma directive.

```
struct s_t1 {
    char a;
    int b;
    #pragma pack(1)
    struct s_t2 {
        char x;
        int y;
    } S2;
    char c;
    int d;
} S1;
```

2. In the code segment below:
 - a. The members of s_t1 would be aligned with the twobyte alignment mode, s_t2 members would be packed on 1-byte, s_t3 members packed on 2-bytes, and s_t4 packed on 4-bytes.
 - b. The **#pragma options align=reset** directive pops the current alignment mode (which in the above case was the twobyte alignment mode). All **#pragma pack** directives issued while the **#pragma options align=twobyte** directive was in effect are also popped.
 - c. The **#pragma pack(4)** directive encountered is restored, as well as the alignment mode that was in effect before the **#pragma options align=twobyte** directive was popped.

```
#pragma pack(4)
#pragma options align=twobyte
struct s_t1 {
    char a;
    int b;
} S1;

#pragma pack(1)
struct s_t2 {
    char a;
    short b;
} S2;

#pragma pack(2)
struct s_t3 {
    char a;
    double b;
} S3;

#pragma options align=reset
struct s_t4 {
    char a;
    int b;
} S4;
```

3. This example shows how a **#pragma pack** directive can affect the size and mapping of a structure:

```
struct s_t {
    char a;
    int b;
    short c;
    int d;
} S;
```

Default mapping:

sizeof s_t = 16
offsetof a = 0
offsetof b = 4
offsetof c = 8
offsetof d = 12
align of a = 1
align of b = 4
align of c = 2
align of d = 4

With #pragma pack(1):

sizeof s_t = 11
offsetof a = 0
offsetof b = 1
offsetof c = 5
offsetof d = 7
align of a = 1
align of b = 1
align of c = 1
align of d = 1

Related References

“General Purpose Pragmas” on page 271

“align” on page 64

“#pragma options” on page 299

#pragma reachable

Description

The **#pragma reachable** directive declares that the point after the call to a routine, *function*, can be the target of a branch from some unknown location. This pragma should be used in conjunction with `setjmp`.

Syntax

```
▶▶ #pragma reachable ( function ) ◀◀
```

Related References

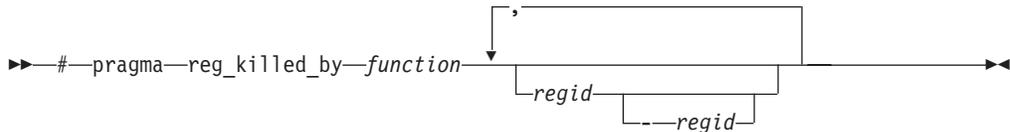
“General Purpose Pragma” on page 271

#pragma reg_killed_by

Description

The `#pragma reg_killed_by` directive specifies a set of volatile registers that may be altered (killed) by the specified function. This pragma can only be used on functions that are defined using `#pragma mc_func`.

Syntax



where:

function The function previously defined using the `#pragma mc_func`.
regid The symbolic name(s) of either a single register or a range of registers to be altered by the named *function*. A range of registers is identified by providing the symbolic names of both starting and ending registers, separated by a dash. If no registers are specified, no registers will be altered by the specified *function*.

The symbolic name is made up of two parts. The first part is the register class name, specified using a sequence of one or more characters in the range of "a" to "z" and/or "A" to "Z".

The second part is a integral number in the range of unsigned int. This number identifies a specific register number within a register class. Some register classes do not require that a register number be specified, and an error will result if you try to do so.

If *regid* is not specified, no volatile registers will be killed by the named *function*.

Registers	
Class and [Register numbers]	Description and usage
ctr	Count register (CTR)
cr[0-7]	Condition register (CR) <ul style="list-style-type: none">Each register in this class is one of the 4-bit fields in the condition register.Of the 8 CR fields, only cr0, cr1, and cr5-cr7 can be specified by <code>#pragma reg_killed_by</code>.
fp[0-31]	Floating point registers (FPR) <ul style="list-style-type: none">Of the 32 machine registers, only fp0-fp13 can be specified by <code>#pragma reg_killed_by</code>.
fs	Floating point status and control register (FPSCR)
lr	Link register (LR)
mq	MQ register (MQ)
gr[0-31]	General purpose registers (GPR) <ul style="list-style-type: none">Of the 32 machine registers, only gr0 and gr3-gr12 can be specified by <code>#pragma reg_killed_by</code>.
vr[0-31]	Vector registers (Altivec processors only)
xer	Fixed point exception (XER)

Notes

Ordinarily, code generated for functions specified by `#pragma mc_func` may alter any or all volatile registers available on your system. You can use `#pragma reg_killed_by` to explicitly list a specific set of volatile registers to be altered by such functions. Registers not in this list will not be altered.

Registers specified by `regid` must meet the following requirements:

- the class name part of the register name must be valid
- the register number is either required or prohibited
- when the register number is required, it must be in the valid range

If any of these requirements are not met, an error is issued and the pragma is ignored.

Example

The following example shows how to use `#pragma reg_killed_by` to list a specific set of volatile registers to be used by the function defined by `#pragma mc_func`.

```
int add_logical(int, int);
#pragma mc_func add_logical {"7c632014" "7c630194"}
    /* addc    r3 <- r3, r4      */
    /* addze   r3 <- r3, carry bit */

#pragma reg_killed_by add_logical gr3, xer
    /* only gpr3 and the xer are altered by this function */

main() {
    int i,j,k;

    i = 4;
    k = -4;
    j = add_logical(i,k);
    printf("\n\nresult = %d\n\n",j);
}
```

Related References

“General Purpose Pragma” on page 271

“#pragma mc_func” on page 297

#pragma stream_unroll

Description

Breaks a stream contained in a **for** loop into multiple streams.

Syntax

```
▶▶ #pragma stream_unroll ( [  $n$  ] ) ▶▶
```

where n is a loop unrolling factor. The value of n is a positive integral constant expression. An unroll factor of 1 disables unrolling. If n is not specified, and if optimization is set at **-O3** or higher, the optimizer determines an appropriate unrolling factor for each nested loop.

Notes

This pragma has effect only on POWER4, POWER5, and PowerPC 970 architectures.

You must specify **-qhot**, **-qipa=level=2**, or **-qsmp** compiler option to enable stream unrolling. An optimization level of **-O4** or higher also allows the compiler to perform stream unrolling.

For stream unrolling to occur, the **#pragma stream_unroll** directive must precede a **for** loop.

You must not specify this directive more than once for a given **for** loop. You should also avoid combining this directive with **block_loop**, **unroll**, **nounroll**, **unrollandfuse**, or **nounrollandfuse** directives for the same **for** loop. No warning messages will be issued, but the compiler may choose to not apply stream_unroll optimization.

Examples

The following is an example of how **#pragma stream_unroll** can increase performance.

```
int i, m, n;
int a[1000][1000];
int b[1000][1000];
int c[1000][1000];

....

#pragma stream_unroll(4)
for (i=1; i<n; i++) {
    a[i] = b[i] * c[i];
}
```

The *unroll factor* of 4 reduces the number of iterations from n to $n/4$, as follows:

```
for (i=1; i<n/4; i++) { /* consider m = n/4 */
    a[i] = b[i] + c[i];
    a[i+m] = b[i+m] + c[i+m];
    a[i+2*m] = b[i+2*m] + c[i+2*m];
    a[i+3*m] = b[i+3*m] + c[i+3*m];
}
```

The increased number of read and store operations are distributed among a number of streams determined by the compiler, reducing computation time and boosting performance.

Related References

“General Purpose Pragmas” on page 271

“unroll” on page 256

“#pragma block_loop” on page 275

“#pragma unroll” on page 315

“#pragma unrollandfuse” on page 317

#pragma strings

Description

The **#pragma strings** directive sets the storage type for string literals, and specifies that the compiler can place strings into read-only memory. Otherwise, the compiler will place strings into read/write memory.

Syntax

```
▶ #pragma strings ( [writeable] | [readonly] ) ▶
```

Notes

Strings are read-only by default.

This pragma must appear before any source statements in order to have effect.

Example

```
#pragma strings(writeable)
```

Related References

“General Purpose Pragas” on page 271

#pragma unroll

Description

The **#pragma unroll** directive is used to unroll the innermost or outermost loops in your program, which can help improve program performance.

Syntax



where n is the loop unrolling factor. The value of n is a positive integral constant expression. An unroll factor of 1 disables unrolling. If n is not specified, and if optimization is set at **-O3** or higher, the optimizer determines an appropriate unrolling factor for each loop.

Notes

The **#pragma unroll** or **#pragma nounroll** directive must appear immediately before either the given **for** loop to be affected, or before any **#pragma block_loop** or SMP and OMP pragmas that apply to the given **for** loop.

You must not specify this directive more than once, or combine this directive with the **block_loop**, **nounrollandfuse**, **unrollandfuse**, or **stream_unroll** directives for the same **for** loop.

The loop structure must meet the following conditions:

- There must be only one loop counter variable, one increment point for that variable, and one termination variable. These cannot be altered at any point in the loop nest.
- Loops cannot have multiple entry and exit points. The loop termination must be the only means to exit the loop.
- Dependencies in the loop must not be "backwards-looking". For example, a statement such as $A[i][j] = A[i - 1][j + 1] + 4$) must not appear within the loop.

Specifying **#pragma nounroll** for a loop instructs the compiler to not unroll that loop. Specifying **#pragma unroll(1)** has the same effect.

To see if the **unroll** option improves performance of a particular application, you should first compile the program with usual options, then run it with a representative workload. You should then recompile with command line **-qunroll** option and/or the **unroll** pragmas enabled, then rerun the program under the same conditions to see if performance improves.

Examples

1. In the following example, loop control is not modified:

```
#pragma unroll(2)
while (*s != 0)
{
    *p++ = *s++;
}
```

Unrolling this by a factor of 2 gives:

```

while (*s)
{
    *p++ = *s++;
    if (*s == 0) break;
    *p++ = *s++;
}

```

2. In this example, loop control is modified:

```

#pragma unroll(3)
for (i=0; i<n; i++) {
    a[i]=b[i] * c[i];
}

```

Unrolling by 3 gives:

```

i=0;
if (i>n-2) goto remainder;
for (; i<n-2; i+=3) {
    a[i]=b[i] * c[i];
    a[i+1]=b[i+1] * c[i+1];
    a[i+2]=b[i+2] * c[i+2];
}
if (i<n) {
    remainder:
    for (; i<n; i++) {
        a[i]=b[i] * c[i];
    }
}

```

Related References

“General Purpose Pragmas” on page 271

“unroll” on page 256

“#pragma block_loop” on page 275

“#pragma loopid” on page 295

“#pragma stream_unroll” on page 312

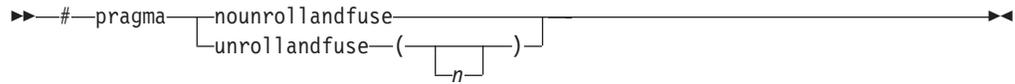
“#pragma unrollandfuse” on page 317

#pragma unrollandfuse

Description

This pragma instructs the compiler to attempt an unroll and fuse operation on nested **for** loops.

Syntax



where n is a loop unrolling factor. The value of n is a positive integral constant expression. An unroll factor of 1 disables unrolling. If n is not specified, and if optimization is set at **-O3** or higher, the optimizer determines an appropriate unrolling factor for each nested loop.

Notes

The **#pragma unrollandfuse** directive applies only to the outer loops of nested **for** loop structures that meet the following conditions:

- There must be only one loop counter variable, one increment point for that variable, and one termination variable. These cannot be altered at any point in the loop nest.
- Loops cannot have multiple entry and exit points. The loop termination must be the only means to exit the loop.
- Dependencies in the loop must not be "backwards-looking". For example, a statement such as $A[i][j] = A[i-1][j+1] + 4$) must not appear within the loop.

For loop unrolling to occur, the **#pragma unrollandfuse** directive must appear immediately before either the given **for** loop to be affected, or before any **#pragma block_loop** or SMP and OMP pragmas that apply to the given **for** loop. You must not specify **#pragma unrollandfuse** for the innermost **for** loop.

You must not specify this directive more than once, or combine this directive with the **block_loop**, **nounrollandfuse**, **nounroll**, **unroll**, or **stream_unroll** directives for the same **for** loop.

Specifying **#pragma nounrollandfuse** instructs the compiler to not unroll that loop.

Examples

1. In the following example, a **#pragma unrollandfuse** directive replicates and fuses the body of the loop. This reduces the number of cache misses for array b .

```
int i, j;
int a[1000][1000];
int b[1000][1000];
int c[1000][1000];

....

#pragma unrollandfuse(2)
for (i=1; i<1000; i++) {
    for (j=1; j<1000; j++) {
        a[j][i] = b[i][j] * c[j][i];
    }
}
```

The `for` loop below shows a possible result of applying the `#pragma unrollandfuse(2)` directive to the loop structure shown above.

```
for (i=1; i<1000; i=i+2) {
    for (j=1; j<1000; j++) {
        a[j][i] = b[i][j] * c[j][i];
        a[j][i+1] = b[i+1][j] * c[j][i+1];
    }
}
```

2. You can also specify multiple `#pragma unrollandfuse` directives in a nested loop structure.

```
int i, j, k;
int a[1000][1000];
int b[1000][1000];
int c[1000][1000];
int d[1000][1000];
int e[1000][1000];

....

#pragma unrollandfuse(4)
for (i=1; i<1000; i++) {
    #pragma unrollandfuse(2)
    for (j=1; j<1000; j++) {
        for (k=1; k<1000; k++) {
            a[j][i] = b[i][j] * c[j][i] + d[j][k] * e[i][k];
        }
    }
}
```

Related References

- “General Purpose Pragmas” on page 271
- “unroll” on page 256
- “#pragma block_loop” on page 275
- “#pragma loopid” on page 295
- “#pragma stream_unroll” on page 312
- “#pragma unroll” on page 315

#pragma weak

Description

The **#pragma weak** directive prevents the link editor from issuing error messages if it does not find a definition for a symbol, or if it encounters a symbol multiply-defined during linking.

Syntax

```
▶ #pragma weak identifier [ = identifier2 ] ▶
```

Notes

While this pragma is intended for use primarily with functions, it will also work for most data objects.

This pragma should not be used with uninitialized global data, or with shared library data objects that are exported to executables.

Two forms of **#pragma weak** can be specified in your program source.

#pragma weak *identifier*

This form of the pragma defines *identifier* as a weak global symbol.

If *Identifier* is defined in the same compilation unit as **#pragma weak** *identifier*, *identifier* is treated as a weak definition. If **#pragma weak** exists in a compilation unit that does not use or declare *identifier*, the pragma is accepted and ignored.

#pragma weak *identifier=identifier2*

This form of the pragma defines *identifier* as a weak global symbol. References to *identifier* will use the value of *identifier2*.

Identifier may or may not be declared in the same compilation unit as the **#pragma weak**, but must never be defined in the compilation unit.

If *identifier* is declared in the compilation unit, *identifier*'s declaration must be compatible to that of *identifier2*. For example, if *identifier2* is a function, *identifier* must have the same return and argument types as *identifier2*.

Identifier2 must be declared in the same compilation unit as **#pragma weak**.

The compiler will ignore **#pragma weak** and issue warning messages if:

- If *identifier2* (if specified) is not defined in the compilation unit.
- If *identifier* is declared but its type is not compatible with that of *identifier2* (if specified).

The compiler will ignore **#pragma weak** and issue a severe error message if the weak *identifier* is defined.

Examples

1. The following is an example of the **#pragma weak** *identifier* form of the pragma:

```

// Begin Compilation Unit 1
#include <stdio.h>
extern int foo;
#pragma weak foo

int main()
{
    int *ptr;
    ptr = &foo;
    if (ptr == 0)
        printf("foo was not defined\n");
    else
        printf("foo was already defined\n");
}
//End Compilation Unit 1

// Begin Compilation Unit 2
int foo = 1;
// End Compilation Unit 2

```

If Compilation Unit 1 is compiled alone to produce an executable, the linker will issue an error message stating that identifier `foo` is not defined. Compilation Unit 1 and Compilation Unit 2 must be combined to produce an executable.

2. The following is an example of the **#pragma weak *identifier=identifier2*** form of the pragma:

```

//Begin Compilation Unit
extern "C" void printf(char *,...);

void fool(void)
{
    printf("Just in function fool()\n");
}

#pragma weak foo__Fv = fool__Fv

int main()
{
    foo();
}
//End Compilation Unit

```

Related References

“General Purpose Pragmas” on page 271

“weaksymbol” on page 266

Pragmas to Control Parallel Processing

The `#pragma` directives on this page give you control over how the compiler handles parallel processing in your program. These pragmas fall into two groups; IBM-specific directives, and directives conforming to the OpenMP Application Program Interface specification.

Use the `-qsmp` compiler option to specify how you want parallel processing handled in your program. You can also instruct the compiler to ignore all parallel processing-related `#pragma` directives by specifying the `-qignprag=ibm:omp` compiler option.

Directives apply only to the statement or statement block immediately following the directive.

IBM Pragma Directives ▶ C	Description
<code>#pragma ibm critical</code>	Instructs the compiler that the statement or statement block immediately following this pragma is a critical section.
<code>#pragma ibm independent_calls</code>	Asserts that specified function calls within the chosen loop have no loop-carried dependencies.
<code>#pragma ibm independent_loop</code>	Asserts that iterations of the chosen loop are independent, and that the loop can therefore be parallelized.
<code>#pragma ibm iterations</code>	Specifies the approximate number of loop iterations for the chosen loop.
<code>#pragma ibm parallel_loop</code>	Explicitly instructs the compiler to parallelize the chosen loop.
<code>#pragma ibm permutation</code>	Asserts that specified arrays in the chosen loop contain no repeated values.
<code>#pragma ibm schedule</code>	Specifies scheduling algorithms for parallel loop execution.
<code>#pragma ibm sequential_loop</code>	Explicitly instructs the compiler to execute the chosen loop sequentially.

OpenMP Pragma Directives ▶ C ▶ C++	Description
<code>#pragma omp atomic</code>	Identifies a specific memory location that must be updated atomically and not be exposed to multiple, simultaneous writing threads.
<code>#pragma omp parallel</code>	Defines a parallel region to be run by multiple threads in parallel. With specific exceptions, all other OpenMP directives work within parallelized regions defined by this directive.
<code>#pragma omp for</code>	Work-sharing construct identifying an iterative for-loop whose iterations should be run in parallel.
<code>#pragma omp parallel for</code>	Shortcut combination of omp parallel and omp for pragma directives, used to define a parallel region containing a single for directive.
<code>#pragma omp ordered</code>	Work-sharing construct identifying a structured block of code that must be executed in sequential order.

OpenMP Pragma Directives ▶ C ▶ C++	Description
#pragma omp section, #pragma omp sections	Work-sharing construct identifying a non-iterative section of code containing one or more subsections of code that should be run in parallel.
#pragma omp parallel sections	Shortcut combination of omp parallel and omp sections pragma directives, used to define a parallel region containing a single sections directive.
#pragma omp single	Work-sharing construct identifying a section of code that must be run by a single available thread.
#pragma omp master	Synchronization construct identifying a section of code that must be run only by the master thread.
#pragma omp critical	Synchronization construct identifying a statement block that must be executed by a single thread at a time.
#pragma omp barrier	Synchronizes all the threads in a parallel region.
#pragma omp flush	Synchronization construct identifying a point at which the compiler ensures that all threads in a parallel region have the same view of specified objects in memory.
#pragma omp threadprivate	Defines the scope of selected file-scope data variables as being private to a thread, but file-scope visible within that thread.

Related Concepts

“Program Parallelization” on page 11

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

“Control Parallel Processing with Pragmas” on page 39

Related References

“smp” on page 232

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

“Built-in Functions Used for Parallel Processing” on page 364

For complete information about the OpenMP Specification, see:

- OpenMP Web site
- OpenMP Specification

#pragma ibm critical



Description

The **critical** pragma identifies a critical section of program code that must only be run by one process at a time.

Syntax

```
#pragma ibm critical [(name)]  
<statement>
```

where *name* can be used to optionally identify the critical region. Identifiers naming a critical region have external linkage.

Notes

The compiler reports an error if you try to branch into or out of a critical section. Some situations that will cause an error are:

- A critical section that contains the **return** statement.
- A critical section that contains **goto**, **continue**, or **break** statements that transfer program flow outside of the critical section.
- A **goto** statement outside a critical section that transfers program flow to a label defined within a critical section.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma ibm independent_calls



Description

The **independent_calls** pragma asserts that specified function calls within the chosen loop have no loop-carried dependencies. This information helps the compiler perform dependency analysis.

Syntax

```
#pragma ibm independent_calls [(identifier [,identifier] ... )]  
<countable for/while/do loop>
```

where *identifier* represents the name of a function.

Notes

identifier cannot be the name of a pointer to a function.

If no function identifiers are specified, the compiler assumes that all functions inside the loop are free of carried dependencies.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma ibm independent_loop

► C

Description

The **independent_loop** pragma asserts that iterations of the chosen loop are independent, and that the loop can be parallelized.

Syntax

```
#pragma ibm independent_loop [if (exp)]  
<countable for/while/do loop>
```

where *exp* represents a scalar expression.

Notes

When the **if** argument is specified, loop iterations are considered independent only as long as *exp* evaluates to TRUE at run-time.

This pragma can be combined with the **schedule** pragma to select a specific parallel process scheduling algorithm. For more information, see the description for the **schedule** pragma.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma ibm schedule” on page 329

#pragma ibm iterations



Description

The **iterations** pragma specifies the approximate number of loop iterations for the chosen loop.

Syntax

```
#pragma ibm iterations (iteration-count)  
<countable for/while/do loop>
```

where *iteration-count* represents a positive integral constant expression.

Notes

The compiler uses the information in the *iteration-count* variable to determine if it is efficient to parallelize the loop.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma ibm parallel_loop

► C

Description

The **parallel_loop** pragma explicitly instructs the compiler to parallelize the chosen loop.

Syntax

```
#pragma ibm parallel_loop [if (exp)] [schedule (sched-type)]  
<countable for/while/do loop>
```

where *exp* represents a scalar expression, and *sched-type* represents any scheduling algorithm as valid for the *schedule* directive.

Notes

When the *if* argument is specified, the loop executes in parallel only if *exp* evaluates to TRUE at run-time. Otherwise the loop executes sequentially. The loop will also run sequentially if it is in a critical section.

This pragma can be applied to a wide variety of C loops, and the compiler will try to determine if a loop is countable or not.

Program sections using the **parallel_loop** pragma must be able to produce a correct result in both sequential and parallel mode. For example, loop iterations must be independent before the loop can be parallelized. Explicit parallel programming techniques involving condition synchronization are not permitted.

The compiler will not automatically detect reductions on loops marked with this pragma. To properly parallelize loops with reductions, use:

- **#pragma omp parallel for** and specify reductions explicitly, or,
- **#pragma ibm independent_loop**, which will let the compiler discover the reductions.

This pragma can be combined with the **schedule** pragma to select a specific parallel process scheduling algorithm. For more information, see the description for the **schedule** pragma.

A warning is generated if this pragma is not followed by a countable loop.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma ibm schedule” on page 329

#pragma ibm permutation



Description

The **permutation** pragma asserts that specified arrays in the chosen loop contain no repeated values.

Syntax

```
#pragma ibm permutation (identifier [,identifier] ... )  
<countable for/while/do loop>
```

where *identifier* represents the name of an array.

Notes

identifier cannot be a function parameter or the name of a pointer.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma ibm schedule



Description

The `schedule` pragma specifies the scheduling algorithms used for parallel processing.

Syntax

```
#pragma ibm schedule (sched-type)  
<countable for/while/do loop>
```

where *sched-type* represents one of the following options:

affinity	Iterations of a loop are initially divided into local partitions of size ceiling (<i>number_of_iterations/number_of_threads</i>). Each local partition is then further subdivided into chunks of size ceiling (<i>number_of_iterations_remaining_in_partition/2</i>). When a thread becomes available, it takes the next chunk from its local partition. If there are no more chunks in the local partition, the thread takes an available chunk from the partition of another thread.
affinity, <i>n</i>	As above, except that each local partition is subdivided into chunks of size <i>n</i> . <i>n</i> must be an integral assignment expression of value 1 or greater.
dynamic	Iterations of a loop are divided into chunks of size 1. Chunks are assigned to threads on a first-come, first-serve basis as threads become available. This continues until all work is completed.
dynamic, <i>n</i>	As above, except that all chunks are set to size <i>n</i> . <i>n</i> must be an integral assignment expression of value 1 or greater.
guided	Chunks are made progressively smaller until a chunk size of one is reached. The first chunk is of size ceiling (<i>number_of_iterations/number_of_threads</i>). Remaining chunks are of size ceiling (<i>number_of_iterations_remaining/number_of_threads</i>). Chunks are assigned to threads on a first-come, first-serve basis as threads become available. This continues until all work is completed.
guided, <i>n</i>	As above, except the minimum chunk size is set to <i>n</i> . <i>n</i> must be an integral assignment expression of value 1 or greater.
runtime	Scheduling policy is determined at run-time.
static	Iterations of a loop are divided into chunks of size ceiling (<i>number_of_iterations/number_of_threads</i>). Each thread is assigned a separate chunk. This scheduling policy is also known as <i>block scheduling</i> .
static, <i>n</i>	Iterations of a loop are divided into chunks of size <i>n</i> . Each chunk is assigned to a thread in <i>round-robin</i> fashion. <i>n</i> must be an integral assignment expression of value 1 or greater.
static,1	This scheduling policy is also known as <i>block cyclic scheduling</i> . Iterations of a loop are divided into chunks of size 1. Each chunk is assigned to a thread in <i>round-robin</i> fashion. This scheduling policy is also known as <i>cyclic scheduling</i> .

Notes

Scheduling algorithms for parallel processing can be specified using any of the methods shown below. If used, methods higher in the list override entries lower in the list.

- pragma statements
- compiler command line options
- run-time command line options
- run-time default options

Scheduling algorithms can also be specified using the **schedule** argument of the **parallel_loop** and **independent_loop** pragma statements. For example, the following sets of statements are equivalent:

```
#pragma ibm parallel_loop
#pragma ibm schedule (sched_type)
<countable for|while|do loop>
and
#pragma ibm parallel_loop (sched_type)
<countable for|while|do loop>
```

If different scheduling types are specified for a given loop, the last one specified is applied.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma ibm independent_loop” on page 325

“#pragma ibm parallel_loop” on page 327

#pragma ibm sequential_loop

► C

Description

The **sequential_loop** pragma explicitly instructs the compiler to execute the chosen loop sequentially.

Syntax

```
#pragma ibm sequential_loop  
<countable for/while/do loop>
```

Notes

This pragma disables automatic parallelization of the chosen loop, and is always respected by the compiler.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp atomic

Description

The **omp atomic** directive identifies a specific memory location that must be updated atomically and not be exposed to multiple, simultaneous writing threads.

Syntax

```
#pragma omp atomic
  <statement_block>
```

where *statement* is an expression statement of scalar type that takes one of the forms that follow:

<i>statement</i>	Conditions
$x \text{ bin_op} = \text{expr}$	where: <i>bin_op</i> is one of: + * - / & ^ << >> <i>expr</i> is an expression of scalar type that does not reference <i>x</i> .
$x++$	
$++x$	
$x--$	
$--x$	

Notes

Load and store operations are atomic only for object *x*. Evaluation of *expr* is not atomic.

All atomic references to a given object in your program must have a compatible type.

Objects that can be updated in parallel and may be subject to race conditions should be protected with the **omp atomic** directive.

Examples

```
extern float x[], *p = x, y;
/* Protect against race conditions among multiple updates. */
#pragma omp atomic
x[index[i]] += y;
/* Protect against races with updates through x. */
#pragma omp atomic
p[i] -= 1.0f;
```

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp parallel

Description

The **omp parallel** directive explicitly instructs the compiler to parallelize the chosen segment of code.

Syntax

```
#pragma omp parallel [clause[[,] clause] ...]  
<statement_block>
```

where *clause* is any of the following:

<i>if</i> (<i>exp</i>)	When the if argument is specified, the program code executes in parallel only if the scalar expression represented by <i>exp</i> evaluates to a non-zero value at run-time. Only one if clause can be specified.
<i>private</i> (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Data variables in <i>list</i> are separated by commas.
<i>firstprivate</i> (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Each new private object is initialized with the value of the original variable as if there was an implied declaration within the statement block. Data variables in <i>list</i> are separated by commas.
<i>num_threads</i> (<i>int_exp</i>)	The value of <i>int_exp</i> is an integer expression that specifies the number of threads to use for the parallel region. If dynamic adjustment of the number of threads is also enabled, then <i>int_exp</i> specifies the maximum number of threads to be used.
<i>shared</i> (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be shared across all threads.
<i>default</i> (<i>shared</i> <i>none</i>)	Defines the default data scope of variables in each thread. Only one default clause can be specified on an omp parallel directive. Specifying default(shared) is equivalent to stating each variable in a shared(list) clause. Specifying default(none) requires that each data variable visible to the parallelized statement block must be explicitly listed in a data scope clause, with the exception of those variables that are: <ul style="list-style-type: none">• const-qualified,• specified in an enclosed data scope attribute clause, or,• used as a loop control variable referenced only by a corresponding omp for or omp parallel for directive.
<i>copyin</i> (<i>list</i>)	For each data variable specified in <i>list</i> , the value of the data variable in the master thread is copied to the thread-private copies at the beginning of the parallel region. Data variables in <i>list</i> are separated by commas. Each data variable specified in the copyin clause must be a threadprivate variable.

reduction (*operator: list*) Performs a reduction on all scalar variables in *list* using the specified *operator*. Reduction variables in *list* are separated by commas.

A private copy of each variable in *list* is created for each thread. At the end of the statement block, the final values of all private copies of the reduction variable are combined in a manner appropriate to the operator, and the result is placed back into the original value of the shared reduction variable.

Variables specified in the **reduction** clause:

- must be of a type appropriate to the operator.
- must be shared in the enclosing context.
- must not be const-qualified.
- must not have pointer type.

Notes

When a parallel region is encountered, a logical team of threads is formed. Each thread in the team executes all statements within a parallel region except for work-sharing constructs. Work within work-sharing constructs is distributed among the threads in a team.

Loop iterations must be independent before the loop can be parallelized. An implied barrier exists at the end of a parallelized statement block.

Nested parallel regions are always serialized.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp for” on page 335

“#pragma omp parallel for” on page 340

“#pragma omp parallel sections” on page 343

#pragma omp for

Description

The **omp for** directive instructs the compiler to distribute loop iterations within the team of threads that encounters this work-sharing construct.

Syntax

```
#pragma omp for [clause[[,] clause] ...]  
<for_loop>
```

where *clause* is any of the following:

private (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Data variables in <i>list</i> are separated by commas.
firstprivate (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Each new private object is initialized as if there was an implied declaration within the statement block. Data variables in <i>list</i> are separated by commas.
lastprivate (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. The final value of each variable in <i>list</i> , if assigned, will be the value assigned to that variable in the last iteration. Variables not assigned a value will have an indeterminate value. Data variables in <i>list</i> are separated by commas.
reduction (<i>operator:list</i>)	Performs a reduction on all scalar variables in <i>list</i> using the specified <i>operator</i> . Reduction variables in <i>list</i> are separated by commas. A private copy of each variable in <i>list</i> is created for each thread. At the end of the statement block, the final values of all private copies of the reduction variable are combined in a manner appropriate to the operator, and the result is placed back into the original value of the shared reduction variable. Variables specified in the reduction clause: <ul style="list-style-type: none">• must be of a type appropriate to the operator.• must be shared in the enclosing context.• must not be const-qualified.• must not have pointer type.
ordered	Specify this clause if an ordered construct is present within the dynamic extent of the omp for directive.

schedule (<i>type</i>)	<p>Specifies how iterations of the for loop are divided among available threads. Acceptable values for <i>type</i> are:</p> <p>dynamic Iterations of a loop are divided into chunks of size ceiling(<i>number_of_iterations</i>/<i>number_of_threads</i>). Chunks are dynamically assigned to threads on a first-come, first-serve basis as threads become available. This continues until all work is completed.</p> <p>dynamic,<i>n</i> As above, except chunks are set to size <i>n</i>. <i>n</i> must be an integral assignment expression of value 1 or greater.</p> <p>guided Chunks are made progressively smaller until the default minimum chunk size is reached. The first chunk is of size ceiling(<i>number_of_iterations</i>/<i>number_of_threads</i>). Remaining chunks are of size ceiling(<i>number_of_iterations_left</i>/<i>number_of_threads</i>). The minimum chunk size is 1. Chunks are assigned to threads on a first-come, first-serve basis as threads become available. This continues until all work is completed.</p> <p>guided,<i>n</i> As above, except the minimum chunk size is set to <i>n</i>. <i>n</i> must be an integral assignment expression of value 1 or greater.</p> <p>runtime Scheduling policy is determined at run-time. Use the OMP_SCHEDULE environment variable to set the scheduling type and chunk size.</p> <p>static Iterations of a loop are divided into chunks of size ceiling(<i>number_of_iterations</i>/<i>number_of_threads</i>). Each thread is assigned a separate chunk. This scheduling policy is also known as <i>block scheduling</i>.</p> <p>static,<i>n</i> Iterations of a loop are divided into chunks of size <i>n</i>. Each chunk is assigned to a thread in <i>round-robin</i> fashion. <i>n</i> must be an integral assignment expression of value 1 or greater. This scheduling policy is also known as <i>block cyclic scheduling</i>.</p> <p>static,1 Iterations of a loop are divided into chunks of size 1. Each chunk is assigned to a thread in <i>round-robin</i> fashion. This scheduling policy is also known as <i>cyclic scheduling</i>.</p>
nowait	<p>Use this clause to avoid the implied barrier at the end of the for directive. This is useful if you have multiple independent work-sharing sections or iterative loops within a given parallel region. Only one nowait clause can appear on a given for directive.</p>

and where *for_loop* is a **for** loop construct with the following canonical shape:

```
for (init_expr; exit_cond; incr_expr)  
  statement
```

where:

<i>init_expr</i>	takes form:	<i>iv</i> = <i>b</i> <i>integer-type iv</i> = <i>b</i>
<i>exit_cond</i>	takes form:	<i>iv</i> <= <i>ub</i> <i>iv</i> < <i>ub</i> <i>iv</i> >= <i>ub</i> <i>iv</i> > <i>ub</i>
<i>incr_expr</i>	takes form:	++ <i>iv</i> <i>iv</i> ++ -- <i>iv</i> <i>iv</i> -- <i>iv</i> += <i>incr</i> <i>iv</i> -= <i>incr</i> <i>iv</i> = <i>iv</i> + <i>incr</i> <i>iv</i> = <i>incr</i> + <i>iv</i> <i>iv</i> = <i>iv</i> - <i>incr</i>

and where:

<i>iv</i>	Iteration variable. The iteration variable must be a signed integer not modified anywhere within the for loop. It is implicitly made private for the duration of the for operation. If not specified as lastprivate , the iteration variable will have an indeterminate value after the operation completes.
<i>b, ub, incr</i>	Loop invariant signed integer expressions. No synchronization is performed when evaluating these expressions and evaluated side effects may result in indeterminate values.

Notes

Program sections using the **omp for** pragma must be able to produce a correct result regardless of which thread executes a particular iteration. Similarly, program correctness must not rely on using a particular scheduling algorithm.

The **for** loop iteration variable is implicitly made private in scope for the duration of loop execution. This variable must not be modified within the body of the **for** loop. The value of the increment variable is indeterminate unless the variable is specified as having a data scope of **lastprivate**.

An implicit barrier exists at the end of the **for** loop unless the **nowait** clause is specified.

Restrictions are:

- The **for** loop must be a structured block, and must not be terminated by a **break** statement.
- Values of the loop control expressions must be the same for all iterations of the loop.
- An **omp for** directive can accept only one **schedule** clauses.
- The value of *n* (chunk size) must be the same for all threads of a parallel region.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp parallel for” on page 340

#pragma omp ordered

Description

The **omp ordered** directive identifies a structured block of code that must be executed in sequential order.

Syntax

```
#pragma omp ordered
    statement_block
```

Notes

The **omp ordered** directive must be used as follows:

- It must appear within the extent of a **omp for** or **omp parallel for** construct containing an **ordered** clause.
- It applies to the statement block immediately following it. Statements in that block are executed in the same order in which iterations are executed in a sequential loop.
- An iteration of a loop must not execute the same **omp ordered** directive more than once.
- An iteration of a loop must not execute more than one distinct **omp ordered** directive.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp for” on page 335

“#pragma omp parallel for” on page 340

#pragma omp parallel for

Description

The **omp parallel for** directive effectively combines the **omp parallel** and **omp for** directives. This directive lets you define a parallel region containing a single **for** directive in one step.

Syntax

```
#pragma omp parallel for [clause[[,] clause] ...]  
<for_loop>
```

Notes

With the exception of the **nowait** clause, clauses and restrictions described in the **omp parallel** and **omp for** directives also apply to the **omp parallel for** directive.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp for” on page 335

“#pragma omp parallel” on page 333

#pragma omp section, #pragma omp sections

Description

The **omp sections** directive distributes work among threads bound to a defined parallel region.

Syntax

```
#pragma omp sections [clause[ clause] ...]
{
  [#pragma omp section]
  statement-block
  [#pragma omp section]
  statement-block
  .
  .
  .
}
```

where *clause* is any of the following:

private (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Data variables in <i>list</i> are separated by commas.
firstprivate (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Each new private object is initialized as if there was an implied declaration within the statement block. Data variables in <i>list</i> are separated by commas.
lastprivate (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. The final value of each variable in <i>list</i> , if assigned, will be the value assigned to that variable in the last section . Variables not assigned a value will have an indeterminate value. Data variables in <i>list</i> are separated by commas.
num_threads (<i>int_exp</i>)	The value of <i>int_exp</i> is an integer expression that specifies the number of threads to use for the parallel region. If dynamic adjustment of the number of threads is also enabled, then <i>int_exp</i> specifies the maximum number of threads to be used.
reduction (<i>operator</i> : <i>list</i>)	Performs a reduction on all scalar variables in <i>list</i> using the specified <i>operator</i> . Reduction variables in <i>list</i> are separated by commas. A private copy of each variable in <i>list</i> is created for each thread. At the end of the statement block, the final values of all private copies of the reduction variable are combined in a manner appropriate to the operator, and the result is placed back into the original value of the shared reduction variable. Variables specified in the reduction clause: <ul style="list-style-type: none">• must be of a type appropriate to the operator.• must be shared in the enclosing context.• must not be const-qualified.• must not have pointer type.
nowait	Use this clause to avoid the implied barrier at the end of the sections directive. This is useful if you have multiple independent work-sharing sections within a given parallel region. Only one nowait clause can appear on a given sections directive.

Notes

The **omp section** directive is optional for the first program code segment inside the **omp sections** directive. Following segments must be preceded by an **omp section**

directive. All **omp section** directives must appear within the lexical construct of the program source code segment associated with the **omp sections** directive.

When program execution reaches a **omp sections** directive, program segments defined by the following **omp section** directive are distributed for parallel execution among available threads. A barrier is implicitly defined at the end of the larger program region associated with the **omp sections** directive unless the **nowait** clause is specified.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp parallel sections” on page 343

#pragma omp parallel sections

Description

The **omp parallel sections** directive effectively combines the **omp parallel** and **omp sections** directives. This directive lets you define a parallel region containing a single **sections** directive in one step.

Syntax

```
#pragma omp parallel sections [clause[[,] clause] ...]
{
  [#pragma omp section]
  statement-block
  [#pragma omp section]
  statement-block
  .
  .
  ]
}
```

Notes

All clauses and restrictions described in the **omp parallel** and **omp sections** directives apply to the **omp parallel sections** directive.

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp parallel” on page 333

“#pragma omp section, #pragma omp sections” on page 341

#pragma omp single

Description

The **omp single** directive identifies a section of code that must be run by a single available thread.

Syntax

```
#pragma omp single [clause[[, clause] ...]  
    statement_block
```

where *clause* is any of the following:

private (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Data variables in <i>list</i> are separated by commas. A variable in the private clause must not also appear in a copyprivate clause for the same omp single directive.
copyprivate (<i>list</i>)	Broadcasts the values of variables specified in <i>list</i> from one member of the team to other members. This occurs after the execution of the structured block associated with the omp single directive, and before any of the threads leave the barrier at the end of the construct. For all other threads in the team, each variable in the <i>list</i> becomes defined with the value of the corresponding variable in the thread that executed the structured block. Data variables in <i>list</i> are separated by commas. Usage restrictions for this clause are: <ul style="list-style-type: none">• A variable in the copyprivate clause must not also appear in a private or firstprivate clause for the same omp single directive.• If an omp single directive with a copyprivate clause is encountered in the dynamic extent of a parallel region, all variables specified in the copyprivate clause must be private in the enclosing context.• Variables specified in copyprivate clause within dynamic extent of a parallel region must be private in the enclosing context.• A variable that is specified in the copyprivate clause must have an accessible and unambiguous copy assignment operator.• The copyprivate clause must not be used together with the nowait clause.
firstprivate (<i>list</i>)	Declares the scope of the data variables in <i>list</i> to be private to each thread. Each new private object is initialized as if there was an implied declaration within the statement block. Data variables in <i>list</i> are separated by commas. A variable in the firstprivate clause must not also appear in a copyprivate clause for the same omp single directive.
nowait	Use this clause to avoid the implied barrier at the end of the single directive. Only one nowait clause can appear on a given single directive. The nowait clause must not be used together with the copyprivate clause.

Notes

An implied barrier exists at the end of a parallelized statement block unless the **nowait** clause is specified.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp master

Description

The **omp master** directive identifies a section of code that must be run only by the master thread.

Syntax

```
#pragma omp master  
    statement_block
```

Notes

Threads other than the master thread will not execute the statement block associated with this construct.

No implied barrier exists on either entry to or exit from the master section.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp critical

Description

The **omp critical** directive identifies a section of code that must be executed by a single thread at a time.

Syntax

```
#pragma omp critical [(name)]  
    statement_block
```

where *name* can optionally be used to identify the critical region. Identifiers naming a critical region have external linkage and occupy a namespace distinct from that used by ordinary identifiers.

Notes

A thread waits at the start of a critical region identified by a given name until no other thread in the program is executing a critical region with that same name. Critical sections not specifically named by **omp critical** directive invocation are mapped to the same unspecified name.

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp barrier

Description

The **omp barrier** directive identifies a synchronization point at which threads in a parallel region will wait until all other threads in that section reach the same point. Statement execution past the **omp barrier** point then continues in parallel.

Syntax

```
#pragma omp barrier
```

Notes

The **omp barrier** directive must appear within a block or compound statement. For example:

```
if (x!=0) {  
    #pragma omp barrier    /* valid usage    */  
}  
if (x!=0)  
    #pragma omp barrier    /* invalid usage */
```

Related References

“Pragmas to Control Parallel Processing” on page 321

#pragma omp flush

Description

The **omp flush** directive identifies a point at which the compiler ensures that all threads in a parallel region have the same view of specified objects in memory.

Syntax

```
#pragma omp flush [ (list) ]
```

where *list* is a comma-separated list of variables that will be synchronized.

Notes

If *list* includes a pointer, the pointer is flushed, not the object being referred to by the pointer. If *list* is not specified, all shared objects are synchronized except those inaccessible with automatic storage duration.

An implied **flush** directive appears in conjunction with the following directives:

- **omp barrier**
- Entry to and exit from **omp critical**.
- Exit from **omp parallel**.
- Exit from **omp for**.
- Exit from **omp sections**.
- Exit from **omp single**.

The **omp flush** directive must appear within a block or compound statement. For example:

```
if (x!=0) {  
    #pragma omp flush /* valid usage */  
}  
if (x!=0)  
    #pragma omp flush /* invalid usage */
```

Related References

“Pragmas to Control Parallel Processing” on page 321

“#pragma omp barrier” on page 347

“#pragma omp critical” on page 346

“#pragma omp for” on page 335

“#pragma omp parallel” on page 333

“#pragma omp parallel for” on page 340

“#pragma omp parallel sections” on page 343

“#pragma omp section, #pragma omp sections” on page 341

“#pragma omp single” on page 344

#pragma omp threadprivate

Description

The **omp threadprivate** directive makes the named file-scope, namespace-scope, or static block-scope variables private to a thread.

Syntax

```
#pragma omp threadprivate (list)
```

where *list* is a comma-separated list of variables.

Notes

Each copy of an **omp threadprivate** data variable is initialized once prior to first use of that copy. If an object is changed before being used to initialize a **threadprivate** data variable, behavior is unspecified.

A thread must not reference another thread's copy of an **omp threadprivate** data variable. References will always be to the master thread's copy of the data variable when executing serial and master regions of the program.

Use of the **omp threadprivate** directive is governed by the following points:

- An **omp threadprivate** directive must appear at file scope outside of any definition or declaration.
- The **omp threadprivate** directive is applicable to static-block scope variables and may appear in lexical blocks to reference those block-scope variables. The directive must appear in the scope of the variable and not in a nested scope, and must precede all references to variables in its list.
- A data variable must be declared with file scope prior to inclusion in an **omp threadprivate** directive *list*.
- An **omp threadprivate** directive and its *list* must lexically precede any reference to a data variable found in that *list*.
- A data variable specified in an **omp threadprivate** directive in one translation unit must also be specified as such in all other translation units in which it is declared.
- Data variables specified in an **omp threadprivate** *list* must not appear in any clause other than the **copyin**, **copyprivate**, **if**, **num_threads**, and **schedule** clauses.
- The address of a data variable in an **omp threadprivate** *list* is not an address constant.
- A data variable specified in an **omp threadprivate** *list* must not have an incomplete or reference type.

Related References

“Pragmas to Control Parallel Processing” on page 321

Acceptable Compiler Mode and Processor Architecture Combinations

You can use the `-q32`, `-q64`, `-qarch`, and `-qtune` compiler options to optimize the output of the compiler to suit:

- the broadest possible selection of target processors,
- a range of processors within a given processor architecture family,
- a single specific processor.

Generally speaking, the options do the following:

- `-q32` selects 32-bit execution mode.
- `-q64` selects 64-bit execution mode.
- `-qarch` selects the general family processor architecture for which instruction code should be generated. Certain `-qarch` settings produce code that will run *only* on RS/6000 systems that support *all* of the instructions generated by the compiler in response to a chosen `-qarch` setting.
- `-qtune` selects the specific processor for which compiler output is optimized. Some `-qtune` settings can also be specified as `-qarch` options, in which case they do not also need to be specified as a `-qtune` option. The `-qtune` option influences only the performance of the code when running on a particular system but does not determine where the code will run.

There are three main families of RS/6000 machines:

- POWER
- POWER2
- PowerPC

All RS/6000 machines share a common set of instructions, but may also include additional instructions unique to a given processor or processor family.

For example, the POWER2 instruction set is a superset of the POWER instructions set. The PowerPC instruction set includes some instructions not available on POWER systems but does not support all of the POWER instruction set. It also includes a number of POWER2 instructions not available in the POWER instruction set. Also, some features found in the POWER2 instruction set may or may not be implemented on particular PowerPC processors. These optional feature groups include:

- support for the graphics instruction group
- support for the sqrt instruction group
- support for 64-bit mode (`-q64` compiler option)

The table below shows some selected processors, and the various features they may or may not support:

Processor	graphics support	sqrt support	64-bit support	large page support
601	no	no	no	no
603	yes	no	no	no
604	yes	no	no	no
rs64a	no	no	yes	no
rs64b	yes	yes	yes	no
rs64c	yes	yes	yes	no
pwr3	yes	yes	yes	no
pwr4	yes	yes	yes	yes, with AIX v5.1D or later

Processor	graphics support	sqrt support	64-bit support	large page support
pwr5	yes	yes	yes	yes, with AIX v5.1D or later

If you want to generate code that will run across a variety of processors, use the following guidelines to select the appropriate **-qarch** and/or **-qtune** compiler options. Code compiled with:

- **-qarch=com** will run on any system.
- **-qarch=pwr** will run on any POWER or POWER2 machine.
- **-qarch=pwr2** (or **pwr2s**, **pwrx**, **p2sc**) will run only on POWER2 machines.
- **-qarch=pwr4** will run only on POWER4 machines.
- **-qarch=pwr5** will run only on POWER5 machines.
- **-qarch=ppc** will run on any PowerPC system.
- **-q64** will run only on PowerPC machines with 64-bit support
- Other **-qarch** options that refer to specific processors will run on any functionally equivalent PowerPC machine. For example, the table that follows shows that code compiled with **-qarch=pwr3** will also run on a **rs64c** but not on a **rs64a**. Similarly, code compiled with **-qarch=603** will run on a **pwr3** but not on a **rs64a**.

If you want to generate code optimized specifically for a particular processor, acceptable combinations of **-q32**, **-q64**, **-qarch**, and **-qtune** compiler options are shown in the following table.

The default execution mode is 32-bit unless the **OBJECT_MODE** environment variable is set to 64 or **-q64** is specified on the command line.

Related Tasks

- “Specify Compiler Options for Architecture-Specific, 32- or 64-bit Compilation” on page 33
- “Set Environment Variables to Select 64- or 32-bit Modes” on page 24

Related References

- “Compiler Command Line Options” on page 51
- “32, 64” on page 60
- “arch” on page 70
- “tune” on page 253

Acceptable -qarch/-qtune Combinations for 32-bit Execution Mode			
-qarch option	Predefined Macro(s)	-qtune default	Available -qtune setting(s)
com	_ARCH_COM	pwr3	auto pwr pwr2 pwr3 pwr4 pwr5 ppc970 p2sc 601 602 603 604 403 rs64a rs64b rs64c
pwr	_ARCH_COM _ARCH_PWR	pwr3	auto pwr pwr2 pwr2s p2sc 601
	_ARCH_COM _ARCH_PWR _ARCH_PWR2	pwr2	auto pwr2 pwr2s p2sc
	_ARCH_COM _ARCH_PWR _ARCH_PWR2 _ARCH_PWR2S	pwr2s	auto pwr2s
	_ARCH_COM _ARCH_PWR _ARCH_PWR2 _ARCH_P2SC	p2sc	auto p2sc
	_ARCH_COM _ARCH_PPC	pwr3	auto 601 602 603 604 403 rs64a rs64b rs64c pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PWR _ARCH_PPC _ARCH_601	601	auto 601
	_ARCH_COM _ARCH_PPC _ARCH_602	602	auto 602
	_ARCH_COM _ARCH_PPC _ARCH_403	403	auto 403
	_ARCH_COM _ARCH_PPC _ARCH_PPCGR	604	auto 603 604 rs64b rs64c pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPC _ARCH_PPCGR _ARCH_603	603	auto 603
ppcgr	_ARCH_COM _ARCH_PPC _ARCH_PPCGR _ARCH_604	604	auto 604
	_ARCH_COM _ARCH_PPC _ARCH_PPC64	pwr3	auto rs64a rs64b rs64c pwr3 pwr4 pwr5 ppc970
ppc64	_ARCH_COM _ARCH_PPC _ARCH_PPC64	rs64a	auto rs64a
	_ARCH_COM _ARCH_PPC _ARCH_PPC64 _ARCH_RS64A	pwr3	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	pwr3	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	pwr3	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64GRSQ	pwr3	auto pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR _ARCH_PPC64GRSQ _ARCH_PWR3	pwr3	auto pwr3 pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR _ARCH_PPC64GRSQ _ARCH_PWR3 _ARCH_PWR4	pwr4	auto pwr4 pwr5 ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR _ARCH_PPC64GRSQ _ARCH_PWR3 _ARCH_PWR4 _ARCH_PWR5	pwr5	auto pwr5
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64GRSQ _ARCH_PWR3 _ARCH_PWR4 _ARCH_PWR4 _ARCH_PPC970	ppc970	auto ppc970
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64GRSQ _ARCH_RS64B	rs64b	auto rs64b
ppc64gr	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64GRSQ _ARCH_RS64B	rs64c	auto rs64c
	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64GRSQ _ARCH_RS64C		

Acceptable -qarch/-qtune Combinations for 64-bit Execution Mode			
-qarch option	Predefined Macro(s)	-qtune default Available -qtune setting(s)	
com	_ARCH_COM	auto pwr pwr2 pwr3 pwr4 pwr5 ppc970 p2sc 601 602 603 604 403 rs64a rs64b rs64c	
ppc	_ARCH_COM _ARCH_PPC	auto 601 602 603 604 403 rs64a rs64b rs64c pwr3 pwr4 pwr5 ppc970	
	ppc64	auto rs64a rs64b rs64c pwr3 pwr4 pwr5 ppc970	
	rs64a	auto rs64a	
	ppc64gr	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970	
	ppc64grsq	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970
		_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto rs64b rs64c pwr3 pwr4 pwr5 ppc970
		_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto pwr3 pwr4 pwr5 ppc970
		_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto pwr4 pwr5 ppc970
		_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto pwr5
		_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto ppc970
rs64b	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto rs64b	
rs64c	_ARCH_COM _ARCH_PPCGR _ARCH_PPC64 _ARCH_PPC64GR	auto rs64c	

Compiler Messages

This section outlines some of the basic reporting mechanisms the compiler uses to describe compilation errors.

- “Message Severity Levels and Compiler Response”
- “Compiler Return Codes”
- “Compiler Message Format” on page 356

Message Severity Levels and Compiler Response

The following table shows the compiler response associated with each level of message severity.

Letter	Severity	Compiler Response
I	Informational	Compilation continues. The message reports conditions found during compilation.
W	Warning	Compilation continues. The message reports questionable and possibly unintended conditions. The program will run as written.
E	Error	Compilation continues and object code is generated. Error conditions exist that the compiler can correct, but the program might not run correctly.
S	Severe error	Compilation continues, but object code is not generated. Error conditions exist that the compiler cannot correct.
U	Unrecoverable error	The compiler halts. An internal compiler error has occurred. <ul style="list-style-type: none">• If the message indicates a resource limit (for example, file system full or paging space full), provide additional resources and recompile.• If the message indicates that different compiler options are needed, recompile using them.• Check for and correct any other errors reported prior to the unrecoverable error.• If the unrecoverable error persists, report the message to your IBM service representative.

Related Concepts

“Compiler Message and Listing Information” on page 8

Related References

“Compiler Return Codes”

“Compiler Message Format” on page 356

“halt” on page 132

“maxerr” on page 188

Compiler Return Codes

At the end of compilation, the compiler sets the return code to zero under any of the following conditions:

- No messages are issued.

- The highest severity level of all errors diagnosed is less than the setting of the **-qhalt** compiler option, and the number of errors did not reach the limit set by the **-qmaxerr** compiler option.

Otherwise, the compiler sets the return code to one of the following values:

Return Code	Error Type
1	Any error with a severity level higher than the setting of the halt compiler option has been detected.
40	An option error or an unrecoverable error has been detected.
41	A configuration file error has been detected.
250	An out-of-memory error has been detected. The <code>xlccommand</code> cannot allocate any more memory for its use.
251	A signal-received error has been detected. That is, an unrecoverable error or interrupt signal has occurred.
252	A file-not-found error has been detected.
253	An input/output error has been detected: files cannot be read or written to.
254	A fork error has been detected. A new process cannot be created.
255	An error has been detected while the process was running.

Note: Return codes may also be displayed for run-time errors. For example, a run-time return code of **99** indicates that a static initialization has failed.

Related Concepts

“Compiler Message and Listing Information” on page 8

Related References

“Message Severity Levels and Compiler Response” on page 355

“Compiler Message Format”

“halt” on page 132

“maxerr” on page 188

Compiler Message Format

Diagnostic messages have the following format when the **-qnosrcmsg** option is active (which is the default):

`"file", line line_number.column_number: 15dd-nnn (severity) text.`

where:

file is the name of the C source file with the error.

line_number is the line number of the error.

column_number is the column number of the error.

15 is the compiler product identifier.

<i>dd</i>	is a two-digit code indicating the XL C Enterprise Edition component that issued the message. <i>dd</i> can have the following values:
00	- code generating or optimizing message
01	- compiler services message.
05	- message specific to the C compiler
06	- message specific to the C compiler
47	- message specific to munch utility
86	- message specific to interprocedural analysis (IPA).
<i>nnn</i>	is the message number.
<i>severity</i>	is a letter representing the severity of the error.
<i>text</i>	is a message describing the error.

Diagnostic messages have the following format when the **-qsrcmsg** option is specified:

x - 15dd-nnn(severity) text.

where *x* is a letter referring to a finger in the finger line.

Related Concepts

“Compiler Message and Listing Information” on page 8

Related References

“Message Severity Levels and Compiler Response” on page 355

“Compiler Return Codes” on page 355

“halt” on page 132

“maxerr” on page 188

Parallel Processing Support

This section contains information on environment variables and built-in functions used to control parallel processing. Topics in this section are:

- “IBM SMP Run-time Options for Parallel Processing”
- “OpenMP Run-time Options for Parallel Processing” on page 362
- “Built-in Functions Used for Parallel Processing” on page 364

IBM SMP Run-time Options for Parallel Processing

Run-time options affecting SMP parallel processing can be specified with the XLSMPOPTS environment variable. This environment variable must be set before you run an application, and uses basic syntax of the form:

```
→→ XLSMPOPTS = [ : option_and_args ] →→
```

Parallelization run-time options can also be specified using OMP environment variables. When run-time options specified by OMP- and XLSMPOPTS-specific environment variables conflict, OMP options will prevail.

Note: You must use thread-safe compiler mode invocations when compiling parallelized program code.

SMP run-time option settings for the XLSMPOPTS environment variable are shown below, grouped by category:

Scheduling Algorithm Options

**XLSMPOPTS
Environment Variable
Option**

`schedule=algorithm=[n]`

Description

This option specifies the scheduling algorithm used for loops not explicitly assigned a scheduling algorithm with the **ibm schedule** pragma.

Valid options for *algorithm* are:

- guided
- affinity
- dynamic
- static

If specified, the chunk size *n* must be an integer value of 1 or greater.

The default scheduling algorithm is **static**.

See **#pragma ibm schedule** for a description of these algorithms.

Parallel Environment Options

XLSPMPOPTS Environment Variable Option	Description
<code>parthds=num</code>	<p><i>num</i> represents the number of parallel threads requested, which is usually equivalent to the number of processors available on the system.</p> <p>Some applications cannot use more threads than the maximum number of processors available. Other applications can experience significant performance improvements if they use more threads than there are processors. This option gives you full control over the number of user threads used to run your program.</p> <p>The default value for <i>num</i> is the number of processors available on the system.</p>
<code>usrthds=num</code>	<p><i>num</i> represents the number of user threads expected.</p> <p>This option should be used if the program code explicitly creates threads, in which case <i>num</i> should be set to the number of threads created.</p> <p>The default value for <i>num</i> is 0.</p>
<code>stack=num</code>	<p><i>num</i> specifies the largest amount of space required for a thread's stack.</p> <p>The default value for <i>num</i> is 4194304.</p>

Performance Tuning Options

XLSPMPOPTS Environment Variable Option	Description
<code>spins=num</code>	<p><i>num</i> represents the number of loop spins, or iterations, before a yield occurs.</p> <p>When a thread completes its work, the thread continues executing in a tight loop looking for new work. One complete scan of the work queue is done during each busy-wait state. An extended busy-wait state can make a particular application highly responsive, but can also harm the overall responsiveness of the system unless the thread is given instructions to periodically scan for and yield to requests from other applications.</p> <p>A complete busy-wait state for benchmarking purposes can be forced by setting both <code>spins</code> and <code>yields</code> to 0.</p> <p>The default value for <i>num</i> is 100.</p>
<code>yields=num</code>	<p><i>num</i> represents the number of yields before a sleep occurs.</p> <p>When a thread sleeps, it completely suspends execution until another thread signals that there is work to do. This provides better system utilization, but also adds extra system overhead for the application.</p> <p>The default value for <i>num</i> is 100.</p>

XLSMPOPTS Environment Variable Option	Description
delays= <i>num</i>	<p><i>num</i> represents a period of do-nothing delay time between each scan of the work queue. Each unit of delay is achieved by running a single no-memory-access delay loop.</p> <p>The default value for <i>num</i> is 500.</p>

Dynamic Profiling Options

XLSMPOPTS Environment Variable Option	Description
profilefreq= <i>num</i>	<p><i>num</i> represents the sampling rate at which each loop is revisited to determine appropriateness for parallel processing.</p> <p>The run-time library uses dynamic profiling to dynamically tune the performance of automatically-parallelized loops. Dynamic profiling gathers information about loop running times to determine if the loop should be run sequentially or in parallel the next time through. Threshold running times are set by the parthreshold and seqthreshold dynamic profiling options, described below.</p> <p>If <i>num</i> is 0, all profiling is turned off, and overheads that occur because of profiling will not occur. If <i>num</i> is greater than 0, running time of the loop is monitored once every <i>num</i> times through the loop.</p> <p>The default for <i>num</i> is 16. The maximum sampling rate is 32. Values of <i>num</i> exceeding 32 are changed to 32.</p>
parthreshold= <i>mSec</i>	<p><i>mSec</i> specifies the expected running time in milliseconds below which a loop must be run sequentially. <i>mSec</i> can be specified using decimal places.</p> <p>If parthreshold is set to 0, a parallelized loop will never be serialized by the dynamic profiler.</p> <p>The default value for <i>mSec</i> is 0.2 milliseconds.</p>
seqthreshold= <i>mSec</i>	<p><i>mSec</i> specifies the expected running time in milliseconds beyond which a loop that has been serialized by the dynamic profiler must revert to being run in parallel mode again. <i>mSec</i> can be specified using decimal places.</p> <p>The default value for <i>mSec</i> is 5 milliseconds.</p>

Related Concepts

- “Program Parallelization” on page 11
- “IBM SMP Directives” on page 11
- “OpenMP Directives” on page 12

Related References

- “OpenMP Run-time Options for Parallel Processing” on page 362
- “smp” on page 232
- “Pragmas to Control Parallel Processing” on page 321
- “Built-in Functions Used for Parallel Processing” on page 364

For complete information about the OpenMP Specification, see:
OpenMP Web site at www.openmp.org
OpenMP Specification at www.openmp.org/specs

OpenMP Run-time Options for Parallel Processing

OpenMP run-time options affecting parallel processing are set by specifying OMP environment variables. These environment variables, use syntax of the form:

►►—*env_variable*—=—*option_and_args*—◄◄

If an OMP environment variable is not explicitly set, its default setting is used.

Parallelization run-time options can also be specified by the XLSMPOPTS environment variable. When OMP and XLSMPOPTS run-time options conflict, OMP options will prevail.

Note: You must use thread-safe compiler mode invocations when compiling parallelized program code.

OpenMP run-time options fall into different categories as described below:

Scheduling Algorithm Environment Variable

OMP_SCHEDULE=*algorithm*

This option specifies the scheduling algorithm used for loops not explicitly assigned a scheduling algorithm with the **omp schedule** directive. For example:

```
OMP_SCHEDULE="guided, 4"
```

Valid options for *algorithm* are:

- dynamic[, *n*]
- guided[, *n*]
- runtime
- static[, *n*]

If specifying a chunk size with *n*, the value of *n* must be an integer value of 1 or greater.

The default scheduling algorithm is **static**.

See “Scheduling Algorithm Options” on page 359 for a description of these algorithms.

Parallel Environment Environment Variables

`OMP_NUM_THREADS=num` *num* represents the number of parallel threads requested, which is usually equivalent to the number of processors available on the system.

This number can be overridden during program execution by calling the `omp_set_num_threads()` runtime library function.

Some applications cannot use more threads than the maximum number of processors available. Other applications can experience significant performance improvements if they use more threads than there are processors. This option gives you full control over the number of user threads used to run your program.

The default value for *num* is the number of processors available on the system.

You can override the setting of `OMP_NUM_THREADS` for a given parallel section by using the `num_threads` clause available in several `#pragma omp` directives.

`OMP_NESTED=TRUE|FALSE`

This environment variable enables or disables nested parallelism. The setting of this environment variable can be overridden by calling the `omp_set_nested()` runtime library function.

If nested parallelism is disabled, nested parallel regions are serialized and run in the current thread.

In the current implementation, nested parallel regions are always serialized. As a result, `OMP_SET_NESTED` does not have any effect, and `omp_get_nested()` always returns 0. If `-qsmp=nested_par` option is on (only in non-strict OMP mode), nested parallel regions may employ additional threads as available. However, no new team will be created to run nested parallel regions.

The default value for `OMP_NESTED` is `FALSE`.

Dynamic Profiling Environment Variable

`OMP_DYNAMIC=TRUE|FALSE` This environment variable enables or disables dynamic adjustment of the number of threads available for running parallel regions.

If set to `TRUE`, the number of threads available for executing parallel regions may be adjusted at runtime to make the best use of system resources. See the description for `profilefreq=num` in “Dynamic Profiling Options” on page 361 for more information.

If set to `FALSE`, dynamic adjustment is disabled.

The default setting is `TRUE`.

Related Concepts

“Program Parallelization” on page 11

“IBM SMP Directives” on page 11

“OpenMP Directives” on page 12

Related References

“IBM SMP Run-time Options for Parallel Processing” on page 359

“smp” on page 232

“Pragmas to Control Parallel Processing” on page 321

“Built-in Functions Used for Parallel Processing”

For complete information about the OpenMP Specification, see:

- OpenMP Web site (www.openmp.org)
- OpenMP Specification (www.openmp.org/specs)

Built-in Functions Used for Parallel Processing

Use these built-in functions to obtain information about the parallel environment. Function definitions for the **omp_** functions can be found in the **omp.h** header file.

Function Prototype	Description
<code>int __parthds(void);</code>	This function returns the value of the parthds run-time option. If the parthds option is not explicitly set by the user, the function returns the default value set by the run-time library. If the -qsmp compiler option was not specified during program compilation, this function returns 1 regardless of run-time options selected.
<code>int __usrthds(void);</code>	This function returns the value of the usrthds run-time option. If the usrthds option is not explicitly set by the user, or the -qsmp compiler option was not specified during program compilation, this function returns 0 regardless of run-time options selected.
<code>int omp_get_num_threads(void);</code>	This function returns the number of threads currently in the team executing the parallel region from which it is called.
<code>void omp_set_num_threads(int num_threads);</code>	This function overrides the setting of the OMP_NUM_THREADS environment variable, and specifies the number of threads to use in parallel regions following this directive. The value <code>num_threads</code> must be a positive integer. If the <code>num_threads</code> clause is present, then for the parallel region it is applied to, it supersedes the number of threads requested by the <code>omp_set_num_threads</code> library function or the OMP_NUM_THREADS environment variable. Subsequent parallel regions are not affected by it.
<code>int omp_get_max_threads(void);</code>	This function returns the maximum value that can be returned by calls to <code>omp_get_num_threads</code> .

Function Prototype	Description
<code>int omp_get_thread_num(void);</code>	This function returns the thread number, within its team, of the thread executing the function. The thread number lies between 0 and <code>omp_get_num_threads()-1</code> , inclusive. The master thread of the team is thread 0.
<code>int omp_get_num_procs(void);</code>	This function returns the maximum number of processors that could be assigned to the program.
<code>int omp_in_parallel(void);</code>	This function returns non-zero if it is called within the dynamic extent of a parallel region executing in parallel; otherwise, it returns 0.
<code>void omp_set_dynamic(int dynamic_threads);</code>	This function enables or disables dynamic adjustment of the number of threads available for execution of parallel regions.
<code>int omp_get_dynamic(void);</code>	This function returns non-zero if dynamic thread adjustment is enabled, and returns 0 otherwise.
<code>void omp_set_nested(int nested);</code>	This function enables or disables nested parallelism.
<code>int omp_get_nested(void);</code>	This function returns non-zero if nested parallelism is enabled and 0 if it is disabled.
<code>void omp_init_lock(omp_lock_t *lock);</code> <code>void omp_init_nest_lock(omp_nest_lock_t *lock);</code>	These functions provide the only means of initializing a lock. Each function initializes the lock associated with the parameter <i>lock</i> for use in subsequent calls.
<code>void omp_destroy_lock(omp_lock_t *lock);</code> <code>void omp_destroy_nest_lock(omp_nest_lock_t *lock);</code>	These functions ensure that the specified lock variable <i>lock</i> is uninitialized.
<code>void omp_set_lock(omp_lock_t *lock);</code> <code>void omp_set_nest_lock(omp_nest_lock_t *lock);</code>	Each of these functions blocks the thread executing the function until the specified lock is available and then sets the lock. A simple lock is available if it is unlocked. A nestable lock is available if it is unlocked or if it is already owned by the thread executing the function.
<code>void omp_unset_lock(omp_lock_t *lock);</code> <code>void omp_unset_nest_lock(omp_nest_lock_t *lock);</code>	These functions provide the means of releasing ownership of a lock.
<code>int omp_test_lock(omp_lock_t *lock);</code> <code>int omp_test_nest_lock(omp_nest_lock_t *lock);</code>	These functions attempt to set a lock but do not block execution of the thread.
<code>double omp_get_wtime(void);</code>	Returns the time elapsed from a fixed starting time. The value of the fixed starting time is determined at the start of the current program, and remains constant throughout program execution.
<code>double omp_get_wtick(void);</code>	Returns the number of seconds between clock ticks.

Note: In the current implementation, nested parallel regions are always serialized. As a result, `omp_set_nested` does not have any effect, and `omp_get_nested` always returns 0.

For complete information about OpenMP run-time library functions, refer to the OpenMP C/C++ Application Program Interface specification.

Related Concepts

“Program Parallelization” on page 11

Related Tasks

“Set Parallel Processing Run-time Options” on page 24

“Control Parallel Processing with Pragmas” on page 39

Related References

“Pragmas to Control Parallel Processing” on page 321

“smp” on page 232

“IBM SMP Run-time Options for Parallel Processing” on page 359

“OpenMP Run-time Options for Parallel Processing” on page 362

Appendix B, “Built-in Functions,” on page 371

Part 4. Appendixes

Appendix A. Predefined Macros

Predefined macros fall into two major categories: those related to the AIX platform and those related to language features. Platform-specific macros are described here.

For information about language-specific macros, see the *Preprocessor Directives* section of the *C/C++ Language Reference*.

Macros related to the platform

The following predefined macros are provided to facilitate porting applications between platforms.

Predefined Macro Name	Description
__BASE_FILE__	Defined to the fully qualified filename of the primary source file.
__BIG_ENDIAN	Defined to 1.
__BIG_ENDIAN__	Defined to 1.
__CALL_SYSV	Defined to 1.
__OPTIMIZE__	Defined to 2 for optimization level -O or -O2 , or to 3 for optimization level -O3 or higher.
__OPTIMIZE_SIZE__	Defined to 1 if the options -qcompact and -O are set. Otherwise it is not defined.
__powerpc	Defined to 1.
__powerpc__	Defined to 1.
__powerpc64__	Defined to 1 when compiling in 64-bit mode. Otherwise it is not defined.
__PPC	Defined to 1.
__PPC__	Defined to 1.
__ppc64	Defined to 1 when compiling in 64-bit mode. Otherwise it is not defined.
__PPC64__	Defined to 1 when compiling in 64-bit mode. Otherwise it is not defined.
__PTRDIFF_TYPE__	Defined to the underlying type of ptrdiff_t on this platform. In 32-bit mode, the value is int . In 64-bit mode, the value is long int .
__SIZE_TYPE__	Defined to the underlying type of size_t on this platform. On this platform, the macro is defined as long unsigned int . In 32-bit mode, the macro is defined as unsigned int . In 64-bit mode, the macro is defined as long int . The compile mode is controlled by the -q32 and -q64 options.
__unix	Defined to 1 on all UNIX-like platforms. Otherwise it is not defined.
__unix__	Defined to 1 on all UNIX-like platforms. Otherwise it is not defined.

Appendix B. Built-in Functions

The compiler provides you with a selection of built-in functions to help you write more efficient programs. This section summarizes the various built-in functions available to you.

You can also find additional built-in functions described at “Built-in Functions Used for Parallel Processing” on page 364.

Name	Prototype	Description
<code>__alignx</code>	<code>void __alignx(int <i>alignment</i>, const void *<i>address</i>);</code>	Informs the compiler that the specified <i>address</i> is aligned at a known compile-time offset. <i>Alignment</i> must be a positive constant integer with a value greater than zero and of a power of two.
<code>__assert1</code>	<code>int __assert1(int, int, int);</code>	Generates TRAP instructions for kernel debugging. See <code>/usr/include/sys/syspest.h</code> .
<code>__assert2</code>	<code>void assert2(int);</code>	Generates TRAP instructions for kernel debugging. See <code>/usr/include/sys/syspest.h</code> .
<code>__bcopy</code>	<code>void __bcopy(char *, char *, int);</code>	Block copy
<code>__bzero</code>	<code>void __bzero(char *, int);</code>	Block zero
<code>__check_lock_mp</code>	<code>unsigned int __check_lock_mp (const int* <i>addr</i>, int <i>old_value</i>, int <i>new_value</i>);</code>	<p>Check Lock on MultiProcessor systems.</p> <p>Conditionally updates a single word variable atomically. <i>addr</i> specifies the address of the single word variable. <i>old_value</i> specifies the old value to be checked against the value of the single word variable. <i>new_value</i> specifies the new value to be conditionally assigned to the single word variable. The word variable must be aligned on a full word boundary.</p> <p>Return values:</p> <ol style="list-style-type: none">1. A return value of false indicates that the single word variable was equal to the old value and has been set to the new value.2. A return value of true indicates that the single word variable was not equal to the old value and has been left unchanged.

Name	Prototype	Description
__check_lockd_mp	unsigned int __check_lockd_mp (const long long int* <i>addr</i> , long long int <i>old_value</i> , long long int <i>new_value</i>);	<p>Check Lock Doubleword on MultiProcessor systems.</p> <p>Conditionally updates a double word variable atomically. <i>addr</i> specifies the address of the double word variable. <i>old_value</i> specifies the old value to be checked against the value of the double word variable. <i>new_value</i> specifies the new value to be conditionally assigned to the double word variable. The double word variable must be aligned on a double word boundary.</p> <p>Return values:</p> <ol style="list-style-type: none"> 1. A return value of false indicates that the double word variable was equal to the old value and has been set to the new value. 2. A return value of true indicates that the double word variable was not equal to the old value and has been left unchanged.
__check_lock_up	unsigned int __check_lock_up (const int* <i>addr</i> , int <i>old_value</i> , int <i>new_value</i>);	<p>Check Lock on UniProcessor systems.</p> <p>Conditionally updates a single word variable atomically. <i>addr</i> specifies the address of the single word variable. <i>old_value</i> specifies the old value to be checked against the value of the single word variable. <i>new_value</i> specifies the new value to be conditionally assigned to the single word variable. The word variable must be aligned on a full word boundary.</p> <p>Return values:</p> <ul style="list-style-type: none"> • A return value of false indicates that the single word variable was equal to the old value, and has been set to the new value. • A return value of true indicates that the single word variable was not equal to the old value and has been left unchanged.

Name	Prototype	Description
__check_lockd_up	unsigned int __check_lockd_up (const long long int* <i>addr</i> , long long int <i>old_value</i> , int long long <i>new_value</i>);	<p>Check Lock Doubleword on UniProcessor systems.</p> <p>Conditionally updates a double word variable atomically. <i>addr</i> specifies the address of the double word variable. <i>old_value</i> specifies the old value to be checked against the value of the double word variable. <i>new_value</i> specifies the new value to be conditionally assigned to the double word variable. The double word variable must be aligned on a double word boundary.</p> <p>Return values:</p> <ul style="list-style-type: none"> • A return value of false indicates that the double word variable was equal to the old value, and has been set to the new value. • A return value of true indicates that the double word variable was not equal to the old value and has been left unchanged.
__cimag (C99)	double __cimag(__complex__ double);	Returns the imaginary part of a complex number.
__cimagf (C99)	float __cimagf(__complex__ float);	Returns the imaginary part of a complex number.
__cimagl (C99)	long double __cimagl(__complex__ long double);	Returns the imaginary part of a complex number.
__clear_lock_mp	void __clear_lock_mp (const int* <i>addr</i> , int <i>value</i>);	<p>Clear Lock on MultiProcessor systems.</p> <p>Atomic store of the <i>value</i> into the single word variable at the address <i>addr</i>. The word variable must be aligned on a full word boundary.</p>
__clear_lockd_mp	void __clear_lockd_mp (const long long int* <i>addr</i> , long long int <i>value</i>);	<p>Clear Lock Doubleword on MultiProcessor systems.</p> <p>Atomic store of the <i>value</i> into the double word variable at the address <i>addr</i>. The double word variable must be aligned on a double word boundary.</p>
__clear_lock_up	void __clear_lock_up (const int* <i>addr</i> , int <i>value</i>);	<p>Clear Lock on UniProcessor systems.</p> <p>Atomic store of the <i>value</i> into the single word variable at the address <i>addr</i>. The word variable must be aligned on a full word boundary.</p>

Name	Prototype	Description
<code>__clear_lockd_up</code>	<code>void __clear_lockd_up (const long long int* <i>addr</i>, long long int <i>value</i>);</code>	Clear Lock Doubleword on UniProcessor systems. Atomic store of the <i>value</i> into the double word variable at the address <i>addr</i> . The double word variable must be aligned on a double word boundary.
<code>__cntlz4</code>	<code>unsigned int __cntlz4(unsigned int);</code>	Count Leading Zeros, 4-Byte Integer
<code>__cntlz8</code>	<code>unsigned int __cntlz8(unsigned long long);</code>	Count Leading Zeros, 8-Byte Integer
<code>__cnttz4</code>	<code>unsigned int __cnttz4(unsigned int);</code>	Count Trailing Zeros, 4-Byte Integer
<code>__cnttz8</code>	<code>unsigned int __cnttz8(unsigned long long);</code>	Count Trailing Zeros, 8-Byte Integer
<code>__conj (C99)</code>	<code>__complex__ double conj(__complex__ double);</code>	Generates the complex conjugate.
<code>__conjf (C99)</code>	<code>__complex__ float conjf(__complex__ float);</code>	Generates the complex conjugate.
<code>__conjl(C99)</code>	<code>__complex__ long double conjl(__complex__ long double);</code>	Generates the complex conjugate.
<code>__cos</code>	<code>double __cos(double);</code>	Returns a cosine value.
<code>__cosf (C99)</code>	<code>float __cosf(float);</code>	Returns a cosine value.
<code>__cosl (C99)</code>	<code>long double __cosl(long double);</code>	Returns a cosine value.
<code>__creal (C99)</code>	<code>double __creal(__complex__ double);</code>	Returns the real part of a complex number.
<code>__crealf (C99)</code>	<code>float __crealf(__complex__ float);</code>	Returns the real part of a complex number.
<code>__creall (C99)</code>	<code>long double __creall(__complex__ long double);</code>	Returns the real part of a complex number.
<code>__dcbt()</code>	<code>void __dcbt (void *);</code>	Data Cache Block Touch. Loads the block of memory containing the specified address into the data cache.
<code>__dcbz()</code>	<code>void __dcbz (void *);</code>	Data Cache Block set to Zero. Sets the specified address in the data cache to zero (0).
<code>__eieio</code>	<code>void __eieio();</code>	Extra name for the existing <code>__iospace_eieio</code> built-in. Compiler will recognize <code>__eieio</code> built-in. Everything except for the name is exactly same as for <code>__iospace_eieio</code> . <code>__eieio</code> is consistent with the corresponding PowerPC instruction name.
<code>__exp</code>	<code>double __exp(double);</code>	Returns the exponential value
<code>__fabs</code>	<code>double __fabs(double);</code>	Returns the absolute value

Name	Prototype	Description
<code>__fabss</code>	<code>float __fabss(float);</code>	Returns the short floating point absolute value
<code>__fcfid</code>	<code>double __fcfid (double);</code>	Floating Convert From Integer Doubleword. The 64-bit signed fixedpoint operand is converted to a double-precision floating-point.
<code>__fctid</code>	<code>double __fctid (double);</code>	Floating Convert to Integer Doubleword. The floating-point operand is converted into 64-bit signed fixed-point integer, using the rounding mode specified by $FPSCR_{RN}$ (Floating-Point Rounding Control field in the Floating-Point Status and Control Register).
<code>__fctidz</code>	<code>double __fctidz (double);</code>	Floating Convert to Integer Doubleword with Rounding towards Zero. The floating-point operand is converted into 64-bit signed fixed-point integer, using the rounding mode Round toward Zero.
<code>__fctiw</code>	<code>double __fctiw (double);</code>	Floating Convert To Integer Word. The floating-point operand is converted to a 32-bit signed fixed-point integer, using the rounding mode specified by $FPSCR_{RN}$ (Floating-Point Rounding Control field in the Floating-Point Status and Control Register).
<code>__fctiwz</code>	<code>double __fctiwz (double);</code>	Floating Convert To Integer Word with Rounding towards Zero. The floating-point operand is converted to a 32-bit signed fixed-point integer, using the rounding mode Round toward Zero.
<code>__fmadd</code>	<code>double __fmadd(double, double, double);</code>	Floating point multiply-add
<code>__fmadds</code>	<code>float __fmadds(float, float, float);</code>	Floating point multiply-add short
<code>__fmsub</code>	<code>double __fmsub(double, double, double);</code>	Floating point multiply-subtract
<code>__fmsubs</code>	<code>float __fmsubs(float, float, float);</code>	Floating point multiply-subtract
<code>__fnabs</code>	<code>double __fnabs(double);</code>	Floating point negative absolute
<code>__fnabss</code>	<code>float __fnabss(float);</code>	Floating point negative absolute

Name	Prototype	Description
<code>__fnmadd</code>	<code>double __fnmadd(double, double, double);</code>	Floating point negative multiply-add
<code>__fnmadds</code>	<code>float __fnmadds (float, float, float);</code>	Floating point negative multiply-add
<code>__fnmsub</code>	<code>double __fnmsub(double, double, double);</code>	Floating point negative multiply-subtract
<code>__fnmsubs</code>	<code>float __fnmsubs (float, float, float);</code>	<code>__fnmsubs (a, x, y) = [- (a * x - y)]</code>
<code>__fre</code>	<code>double __fre (double);</code>	Floating point reciprocal <code>__fre (x) = [(estimate of) 1.0/x]</code>
<code>__fres</code>	<code>float __fres (float);</code>	Floating point reciprocal <code>__fres (x) = [(estimate of) 1.0/x]</code>
<code>__frsqrte</code>	<code>double __frsqrte (double);</code>	Floating point reciprocal square root <code>__frsqrte (x) = [(estimate of) 1.0/sqrt(x)]</code>
<code>__frsqrtes</code>	<code>float __frsqrtes (float);</code>	Floating point reciprocal square root <code>__frsqrtes (x) = [(estimate of) 1.0/sqrt(x)]</code>
<code>__fsel</code>	<code>double __fsel (double, double, double);</code>	Floating point select if ($a \geq 0.0$) then <code>__fsel (a, x, y) = x</code> ; else <code>__fsel (a, x, y) = y</code>
<code>__fsels</code>	<code>float __fsels (float, float, float);</code>	Floating point select if ($a \geq 0.0$) then <code>__fsels (a, x, y) = x</code> ; else <code>__fsels (a, x, y) = y</code>
<code>__fsqrt</code>	<code>double __fsqrt (double);</code>	Floating point square root <code>__fsqrt (x) = square root of x</code>
<code>__fsqrts</code>	<code>float __fsqrts (float);</code>	Floating point square root <code>__fsqrts (x) = square root of x</code>
<code>__iospace_eieio</code>	<code>void __iospace_eieio(void);</code>	Generates an EIEIO instruction
<code>__iospace_lwsync</code>	(equivalent to: <code>void __iospace_lwsync(void);</code>)	Generates a lwsync instruction
<code>__iospace_sync</code>	(equivalent to: <code>void __iospace_sync(void);</code>)	Generates a sync instruction
<code>__isync</code>	<code>void __isync(void);</code>	Waits for all previous instructions to complete and then discards any prefetched instructions, causing subsequent instructions to be fetched (or refetched) and executed in the context established by previous instructions.
<code>__load2r</code>	<code>unsigned short __load2r(unsigned short*);</code>	Load halfword byte reversed
<code>__load4r</code>	<code>unsigned int __load4r(unsigned int*);</code>	Load word byte reversed

Name	Prototype	Description
<code>__lwsync</code>	<code>void __lwsync();</code>	Extra name for the existing <code>__iospace_lwsync</code> built-in. Compiler will recognize <code>__lwsync</code> built-in. Everything except for the name is exactly same as for <code>__iospace_lwsync</code> . <code>__lwsync</code> is consistent with the corresponding PowerPC instruction name. This function is supported only by the PowerPC 970 processor. Value must be known at compile time.
<code>__mfocr</code>	<code>unsigned __mfocr(const int);</code>	Returns value of device control register. This is a privileged instruction valid only for PowerPC 440.
<code>__mtocr</code>	<code>void __mtocr(const int, unsigned long);</code>	Set value of device control register with unsigned long value. Value must be known at compile time. This is a privileged instruction valid only for PowerPC 440.
<code>__mfspr</code>	<code>unsigned __mfspr(const int);</code>	Returns value of special purpose register. Value of "const int" must be known at compile time.
<code>__mtspr</code>	<code>void __mtspr(const int, unsigned long);</code>	Set value of special purpose register with unsigned long value. Values must be known at compile time.
<code>__mfmsr</code>	<code>unsigned __mfmsr();</code>	Returns value of machine state register. This is a privileged instruction on PowerPC systems.
<code>__mtmsr</code>	<code>void (unsigned long);</code>	Set value of machine state register with unsigned long value. This is a privileged instruction.
<code>__mtfsb0</code>	<code>void __mtfsb0(unsigned int <i>bt</i>);</code>	Move to FPSCR Bit 0. Bit <i>bt</i> of the FPSCR is set to 0. <i>bt</i> must be a constant and $0 \leq bt \leq 31$.
<code>__mtfsb1</code>	<code>void __mtfsb1(unsigned int <i>bt</i>);</code>	Move to FPSCR Bit 1. Bit <i>bt</i> of the FPSCR is set to 1. <i>bt</i> must be a constant and $0 \leq bt \leq 31$.
<code>__mtfsf</code>	<code>void __mtfsf(unsigned int <i>flm</i>, unsigned int <i>frb</i>);</code>	Move to FPSCR Fields. The contents of <i>frb</i> are placed into the FPSCR under control of the field mask specified by <i>flm</i> . The field mask <i>flm</i> identifies the 4-bit fields of the FPSCR affected. <i>flm</i> must be a constant 8-bit mask.

Name	Prototype	Description
<code>__mtfsfi</code>	<code>void __mtfsfi(unsigned int <i>bf</i>, unsigned int <i>u</i>);</code>	Move to FPSCR Field Immediate. The value of the <i>u</i> is placed into FPSCR field specified by <i>bf</i> . <i>bf</i> and <i>u</i> must be constants, with $0 \leq bf \leq 7$ and $0 \leq u \leq 15$.
<code>__mulhd</code>	<code>long long int __mulhd(long long int <i>ra</i>, long long int <i>rb</i>);</code>	Multiply High Doubleword Signed. Returns the highorder 64 bits of the 128-bit product of the operands <i>ra</i> and <i>rb</i> .
<code>__mulhdu</code>	<code>unsigned long long int __mulhdu(unsigned long long int <i>ra</i>, unsigned long long int <i>rb</i>);</code>	Multiply High Doubleword Unsigned. Returns the highorder 64 bits of the 128bit product of the operands <i>ra</i> and <i>rb</i> .
<code>__mulhw</code>	<code>int __mulhw(int <i>ra</i>, int <i>rb</i>);</code>	Multiply High Word Signed. Returns the highorder 32 bits of the 64bit product of the operands <i>ra</i> and <i>rb</i> .
<code>__mulhwu</code>	<code>unsigned int __mulhwu(unsigned int <i>ra</i>, unsigned int <i>rb</i>);</code>	Multiply High Word Unsigned. Returns the highorder 32 bits of the 64bit product of the operands <i>ra</i> and <i>rb</i> .
<code>__parthds</code>	<code>int __parthds(void);</code>	Returns the value of the <code>parthds</code> run-time option. If the <code>parthds</code> option is not explicitly set by the user, the function returns the default value set by the run-time library. If the <code>-qsmp</code> compiler option was not specified during program compilation, this function returns 1 regardless of run-time options selected.
<code>__popcnt4</code>	<code>int __popcnt4(unsigned int);</code>	Returns the number of bits set for a 32-bit integer
<code>__popcnt8</code>	<code>int __popcnt8(unsigned long long);</code>	Returns the number of bits set for a 64-bit integer
<code>__popcntb</code>	<code>unsigned long __popcntb(unsigned long);</code>	Returns the number of 1-bits in each byte of the source operand, and places that count into the corresponding byte of the result. Valid only for POWER5 architectures.
<code>__poppar4</code>	<code>int __poppar4(unsigned int);</code>	Returns 1 if there is an odd number of bits set in a 32-bit integer. Returns 0 otherwise.

Name	Prototype	Description
__poppar8	int __poppar8(unsigned long long);	Returns 1 if there is an odd number of bits set in a 64-bit integer. Returns 0 otherwise.
__pow	double __pow(double, double);	Exponentiation built-in.
__prefetch_by_load	void __prefetch_by_load(const void*);	Touch a memory location via explicit load
__prefetch_by_stream	void __prefetch_by_stream(const int, const void*);	Touch a memory location via explicit stream
__protected_stream_count	void __protected_stream_count(unsigned int <i>unit_cnt</i> , unsigned int <i>ID</i>);	Sets <i>unit_cnt</i> number of cache lines for the limited-length protected stream with identifier <i>ID</i> . <i>Unit_cnt</i> must be an integer with value of 0 to 1023. Stream <i>ID</i> must have integer value 0 to 15. (POWER5 processors only)
__protected_stream_go	void __protected_stream_go();	Starts prefetching all limited-length protected streams. (POWER5 processors only)
__protected_stream_set	void __protected_stream_set(unsigned int <i>direction</i> , const void* <i>addr</i> , unsigned int <i>ID</i>);	Establishes a limited-length protected stream using identifier <i>ID</i> , which begins with the cache line at <i>addr</i> and then depending on the value of <i>direction</i> , fetches from either incremental (forward) or decremental (backward) memory addresses. The stream is protected from being replaced by any hardware-detected streams. <i>Direction</i> must have value of 1 (forward) or 3 (backward). Stream <i>ID</i> must have integer value 0 to 15. (POWER5 processors only)
__protected_unlimited_stream_set_go	void __protected_unlimited_stream_set_go(unsigned int <i>direction</i> , const void* <i>addr</i> , unsigned int <i>ID</i>);	Establishes an unlimited-length protected stream using identifier <i>ID</i> , which begins with the cache line at <i>addr</i> and then depending on the value of <i>direction</i> , fetches from either incremental (forward) or decremental (backward) memory addresses. The stream is protected from being replaced by any hardware-detected streams. <i>Direction</i> must have value of 1 (forward) or 3 (backward). Stream <i>ID</i> must have integer value 0 to 15. (PowerPC 970 and POWER5 processors only)

Name	Prototype	Description
__protected_stream_stop	void __protected_stream_stop(unsigned int <i>ID</i>);	Stops prefetching the protected steam with identifier <i>ID</i> . (POWER5 processors only)
__protected_stream_stop_all	void __protected_stream_stop_all();	Stops prefetching all protected steams. (POWER5 processors only)
__rdlam	unsigned long long __rdlam(unsigned long long <i>rs</i> , unsigned int <i>shift</i> , unsigned long long <i>mask</i>);	Rotate Double Left and AND with Mask. The contents of <i>rs</i> are rotated left <i>shift</i> bits. The rotated data is ANDed with the mask and returned as a result. <i>mask</i> must be a constant and represent a contiguous bit field.
__readflm	double __readflm();	Read floating point status/control register
__rldimi	unsigned long long __rldimi(unsigned long long <i>rs</i> , unsigned long long <i>is</i> , unsigned int <i>shift</i> , unsigned long long <i>mask</i>);	Rotate Left Doubleword Immediate then Mask Insert. Rotates <i>rs</i> left <i>shift</i> bits then inserts <i>rs</i> into <i>is</i> under bit mask <i>mask</i> . Shift must be a constant and $0 \leq \text{shift} \leq 63$. <i>mask</i> must be a constant and represent a contiguous bit field.
__rlwimi	unsigned int __rlwimi(unsigned int <i>rs</i> , unsigned int <i>is</i> , unsigned int <i>shift</i> , unsigned int <i>mask</i>);	Rotate Left Word Immidiate then Mask Insert. Rotates <i>rs</i> left <i>shift</i> bits then inserts <i>rs</i> into <i>is</i> under bit mask <i>mask</i> . Shift must be a constant and $0 \leq \text{shift} \leq 31$. <i>mask</i> must be a constant and represent a contiguous bit field.
__rlwnm	unsigned int __rlwnm(unsigned int <i>rs</i> , unsigned int <i>shift</i> , unsigned int <i>mask</i>);	Rotate Left Word then AND with Mask. Rotates <i>rs</i> left <i>shift</i> bits, then ANDs <i>rs</i> with bit mask <i>mask</i> . <i>mask</i> must be a constant and represent a contiguous bit field.
__rotatel4	unsigned int __rotatel4(unsigned int <i>rs</i> , unsigned int <i>shift</i>);	Rotate Left Word. Rotates <i>rs</i> left <i>shift</i> bits.
__setflm	double __setflm(double);	Set Floating Point Status/Control Register
__setrnd	double __setrnd(int);	Set Rounding Mode

Name	Prototype	Description
__stfiw	void __stfiw(const int* <i>addr</i> , double <i>value</i>);	Store Floating-Point as Integer Word. The contents of the loworder 32 bits of <i>value</i> are stored, without conversion, into the word in storage addressed by <i>addr</i> .
__store2r	void __store2r(unsigned short, unsigned short *);	Store 2-byte reversed (opposite endian mode)
__store4r	void __store4r(unsigned int, unsigned int *);	Store 4-byte reversed (opposite endian mode)
__swdiv	double __swdiv(double <i>numerator</i> , double <i>denominator</i>);	Floating-point division of double types. No argument restrictions. With -qstrict in effect, the result is bitwise identical to IEEE division. With -qnostrict in effect, the result may differ slightly from the IEEE result. This function may provide increased performance over the normal divide operator in situations where division is performed repeatedly within a loop.
__swdiv_nochk	double __swdiv_nochk(double <i>numerator</i> , double <i>denominator</i>);	Floating-point division of double types; no range-checking. The following argument types are not permitted: numerators equal to infinity; and denominators equal to infinity, zero, or denormalized. With -qstrict in effect, the result is bitwise identical to IEEE division, except when the numerator is smaller than around 2^{-1000} in magnitude, in which case the result may differ slightly from the IEEE result. With -qnostrict in effect, the result may differ slightly from the IEEE result. This function may provide increased performance over the normal divide operator or <code>__swdiv</code> built-in function in situations where division is performed repeatedly within a loop, and can be used where arguments are known to be within the permitted ranges.

Name	Prototype	Description
__swdivs	float __swdivs(float <i>numerator</i> , float <i>denominator</i>);	<p>Floating-point division of float types. No argument restrictions. The result is bitwise identical to IEEE division.</p> <p>This function may provide increased performance over the normal divide operator in situations where division is performed repeatedly within a loop.</p>
__swdivs_nochk	float __swdivs_nochk(float <i>numerator</i> , float <i>denominator</i>);	<p>Floating-point division of float types; no range-checking. The following argument types are not permitted: numerators equal to infinity; and denominators equal to infinity, zero, or denormalized. The result is bitwise identical to IEEE division.</p> <p>This function may provide increased performance over the normal divide operator or __swdivs built-in function in situations where division is performed repeatedly within a loop, and can be used where arguments are known to be within the permitted ranges.</p>
__sync	void __sync();	<p>Extra name for the existing __iospace_sync built-in.</p> <p>Compiler will recognize __sync built-in. Everything except for the name is exactly same as for __iospace_sync. __sync is consistent with the corresponding PowerPC instruction name.</p>

Name	Prototype	Description
__tdw	void __tdw(long long <i>a</i> , long long <i>b</i> , unsigned int <i>TO</i>);	<p>Trap Doubleword.</p> <p>Operand <i>a</i> is compared with operand <i>b</i>. This comparison results in five conditions which are ANDed with a 5-bit constant <i>TO</i> containing a value of 0 to 31 inclusive.</p> <p>If the result is not 0 the system trap handler is invoked. Each bit position, if set, indicates one or more of the following possible conditions:</p> <p>0 (high-order bit) <i>a</i> Less Than <i>b</i>, using signed comparison.</p> <p>1 <i>a</i> Greater Than <i>b</i>, using signed comparison.</p> <p>2 <i>a</i> Equal <i>b</i></p> <p>3 <i>a</i> Less Than <i>b</i>, using unsigned comparison.</p> <p>4 (low order bit) <i>a</i> Greater Than <i>b</i>, using unsigned comparison.</p>
__trap	void __trap(int);	Trap
__trapd	void __trapd (longlong);	Trap if the parameter is not zero.
__tw	void __tw(int <i>a</i> , int <i>b</i> , unsigned int <i>TO</i>);	<p>Trap Word.</p> <p>Operand <i>a</i> is compared with operand <i>b</i>. This comparison results in five conditions which are ANDed with a 5-bit constant <i>TO</i> containing a value of 0 to 31 inclusive.</p> <p>If the result is not 0 the system trap handler is invoked. Each bit position, if set, indicates one or more of the following possible conditions:</p> <p>0 (high-order bit) <i>a</i> Less Than <i>b</i>, using signed comparison.</p> <p>1 <i>a</i> Greater Than <i>b</i>, using signed comparison.</p> <p>2 <i>a</i> Equal <i>b</i></p> <p>3 <i>a</i> Less Than <i>b</i>, using unsigned comparison.</p> <p>4 (low order bit) <i>a</i> Greater Than <i>b</i>, using unsigned comparison.</p>

Name	Prototype	Description
__usrthds	int __usrthds(void);	<p>Returns the value of the usrthds run-time option.</p> <p>If the usrthds option is not explicitly set by the user, or the -qsmp compiler option was not specified during program compilation, this function returns 0 regardless of run-time options selected.</p>

Appendix C. National Languages Support in XL C Enterprise Edition

This and related pages summarize the national language support (NLS) specific to IBM XL C Enterprise Edition.

For more information, see the following topics in this section:

- “Converting Files Containing Multibyte Data to New Code Pages”
- “Multibyte Character Support”

See also National Language Support in General Programming Concepts: Writing and Debugging Programs.

Converting Files Containing Multibyte Data to New Code Pages

If you have installed new code pages on your system, you can use the AIX `iconv` migration utility to convert files containing multibyte data to use new code pages. This command converts files containing multibyte data from the **IBM-932** code set to the **IBM-euc** code set.

The `iconv` command is described in the *AIX Commands Reference*. Using the NLS code set converters with the `iconv` command is described in “Converters Overview for Programming” in the *AIX General Programming Concepts*.

Related References

Appendix C, “National Languages Support in XL C Enterprise Edition”
“Multibyte Character Support”

See also:

`iconv` command in Commands Reference, Volume 3: i through m:
Converters Overview for Programming section in AIX 5L Version 5.1 General Programming Concepts: Writing and Debugging Programs

Multibyte Character Support

Support for multibyte characters includes support for wide characters. Generally, wide characters are permitted anywhere multibyte characters are, but they are incompatible with multibyte characters in the same string because their bit patterns differ. Wherever permitted, you can mix single-byte and multibyte characters in the same string.

Note: You must specify the `-qmbcs` option to use multibyte characters anywhere in your program.

In the examples that follow, *multibyte_char* represents any string of one or more multibyte characters.

String Literals and Character Constants

Multibyte characters are supported in string literals and character constants. Strings containing multibyte characters are treated essentially the same way as strings without multibyte characters. Multibyte characters can appear in several contexts:

- Preprocessor directives
- Macro definitions
- The # and ## operators
- The definition of the macro name in the **-D** compiler option

Wide-character strings can be manipulated the same way as single-byte character strings. The system provides equivalent wide-character and single-byte string functions.

For all compiler invocations except **cc** and its derivatives, the default storage type for string literals is read-only. The **-qro** option sets the storage type of string literals to read-only, and the **-qnor** option makes string literals writable.

Note: Because a character constant can store only 1 byte, avoid assigning multibyte characters to character constants. Only the last byte of a multibyte character constant is stored. Use a wide-character representation instead. Wide-character string literals and constants must be prefixed by **L**. For example:

```
wchar_t *a = L"wide_char_string";
wchar_t b = L'c';
```

Preprocessor Directives

The following preprocessor directives permit multibyte-character constants and string literals:

- **#define**
- **#pragma comment**
- **#include**

Macro Definitions

Because string literals and character constants can be part of **#define** statements, multibyte characters are also permitted in both object-like and function-like macro definitions.

Compiler Options

Multibyte characters can appear in the compiler suboptions that take file names as arguments:

- **-B** *prefix*
- **-F** *config_file:stanza*
- **-I** *directory*
- **-L** *directory*
- **-l** *key*
- **-o** *file_name*
- **-qexpfile=***file_name*
- **-qipa=***file_name*
- **-qtempinc=***directory*
- **-qtemplateregistry=***registry_file_name*
- **-qpath** *path_name*

The **-Dname=definition** option permits multibyte characters in the definition of the macro name. In the following example, the first definition is a string literal, and the second is a character constant:

```
-DMYMACRO="kpsmultibyte_chardcs"  
-DMYMACRO='multibyte_char'
```

The **-qmbcs** compiler option permits both double-byte and multibyte characters. In other respects, it is equivalent to the **-qdbcs** option, but it should be used when multibyte characters appear in the program.

The listings produced by the **-qlist** and **-qsource** options display the date and time for the appropriate international language. Multibyte characters in the file name of the C source file also appear in the name of the corresponding list file. For example, a C source file called:

```
multibyte_char.c
```

gives a list file called

```
multibyte_char.lst
```

File Names and Comments

Any file name can contain multibyte characters. The file name can be a relative or absolute path name. For example:

```
#include<multibyte_char/mydir/mysource/multibyte_char.h>  
#include "multibyte_char.h"
```

```
x1C /u/myhome/c_programs/kanji_files/multibyte_char.c -omultibyte_char
```

Multibyte characters are also permitted in comments, if you specify the **-qmbcs** compiler option.

Restrictions

- Multibyte characters are not permitted in identifiers.
- Hexadecimal values for multibyte characters must be in the range of the code page being used.
- You cannot mix wide characters and multibyte characters in macro definitions. For example, a macro expansion that concatenates a wide string and a multibyte string is not permitted.
- Assignment between wide characters and multibyte characters is not permitted.
- Concatenating wide character strings and multibyte character strings is not permitted.

Related References

Appendix C, "National Languages Support in XL C Enterprise Edition," on page 385

"Converting Files Containing Multibyte Data to New Code Pages" on page 385

"mbs, ducs" on page 191

Appendix D. Problem Solving

Topics in this section are:

- “Message Catalog Errors”
- “Correcting Paging Space Errors During Compilation”

Message Catalog Errors

Before the compiler can compile your program, the message catalogs must be installed and the environment variables **LANG** and **NLSPATH** must be set to a language for which the message catalog has been installed.

If you see the following message during compilation, the appropriate message catalog cannot be opened:

```
Error occurred while initializing the message system in
file: message_file
```

where *message_file* is the name of the message catalog that the compiler cannot open. This message is issued in English only.

You should then verify that the message catalogs and the environment variables are in place and correct. If the message catalog or environment variables are not correct, compilation can continue, but diagnostic messages are suppressed and the following message is issued instead:

```
No message text for message_number.
```

where *message_number* is the IBM XL C Enterprise Edition internal message number. This message is issued in English only.

To determine what message catalogs are installed on your system, and assuming that you have installed the compiler to the default installation location, you can list all of the file names for the catalogs by using the following command:

```
ls /usr/lib/nls/msg/%L/*.cat
```

where %L is the current primary language environment (locale) setting. The default locale is **C**. The locale for United States English is **en_US**.

The default message catalogs in **/usr/vacpp/exe/default_msg** are called when the locale has never been changed from the default, **C**.

For more information about the **NLSPATH** and **LANG** environment variables, see your operating system documentation.

Related Tasks

“Set Environment Variables” on page 23

“Set Other Environment Variables” on page 24

Correcting Paging Space Errors During Compilation

If the operating system runs low on paging space during a compilation, the compiler issues one of the following messages:

```
1501-229 Compilation ended due to lack of space.  
1501-224 fatal error in ../exe/xlCcode: signal 9 received.
```

If lack of paging space causes other compiler programs to fail, the following message is displayed:

```
Killed.
```

To minimize paging-space problems, do any of the following and recompile your program:

- Reduce the size of your program by splitting it into two or more source files
- Compile your program without optimization.
- Reduce the number of processes competing for system paging space.
- Increase the system paging space.

To check the current paging-space settings enter the command: **lsps -a** or use the AIX System Management Interface Tool (SMIT) command **smit pgsp**.

See your operating system documentation for more information about paging space and how to allocate it.

Appendix E. ASCII Character Set

XL C Enterprise Edition uses the American National Standard Code for Information Interchange (ASCII) character set.

The following table lists the standard ASCII characters in ascending numerical order, with their corresponding decimal, octal, and hexadecimal values. It also shows the control characters with **Ctrl-** notation. For example, the carriage return (ASCII symbol **CR**) appears as **Ctrl-M**, which you enter by simultaneously pressing the **Ctrl** key and the **M** key.

Decimal Value	Octal Value	Hex Value	Control Character	ASCII Symbol	Meaning
0	0	00	Ctrl-@	NUL	null
1	1	01	Ctrl-A	SOH	start of heading
2	2	02	Ctrl-B	STX	start of text
3	3	03	Ctrl-C	ETX	end of text
4	4	04	Ctrl-D	EOT	end of transmission
5	5	05	Ctrl-E	ENQ	enquiry
6	6	06	Ctrl-F	ACK	acknowledge
7	7	07	Ctrl-G	BEL	bell
8	10	08	Ctrl-H	BS	backspace
9	11	09	Ctrl-I	HT	horizontal tab
10	12	0A	Ctrl-J	LF	new line
11	13	0B	Ctrl-K	VT	vertical tab
12	14	0C	Ctrl-L	FF	form feed
13	15	0D	Ctrl-M	CR	carriage return
14	16	0E	Ctrl-N	SO	shift out
15	17	0F	Ctrl-O	SI	shift in
16	20	10	Ctrl-P	DLE	data link escape
17	21	11	Ctrl-Q	DC1	device control 1
18	22	12	Ctrl-R	DC2	device control 2
19	23	13	Ctrl-S	DC3	device control 3
20	24	14	Ctrl-T	DC4	device control 4
21	25	15	Ctrl-U	NAK	negative acknowledge
22	26	16	Ctrl-V	SYN	synchronous idle
23	27	17	Ctrl-W	ETB	end of transmission block
24	30	18	Ctrl-X	CAN	cancel
25	31	19	Ctrl-Y	EM	end of medium
26	32	1A	Ctrl-Z	SUB	substitute
27	33	1B	Ctrl-[ESC	escape
28	34	1C	Ctrl-\	FS	file separator

Decimal Value	Octal Value	Hex Value	Control Character	ASCII Symbol	Meaning
29	35	1D	Ctrl-]	GS	group separator
30	36	1E	Ctrl-^	RS	record separator
31	37	1F	Ctrl- <u></u>	US	unit separator
32	40	20		SP	digit select
33	41	21		!	exclamation point
34	42	22		"	double quotation mark
35	43	23		#	pound sign, number sign
36	44	24		\$	dollar sign
37	45	25		%	percent sign
38	46	26		&	ampersand
39	47	27		'	apostrophe
40	50	28		(left parenthesis
41	51	29)	right parenthesis
42	52	2A		*	asterisk
43	53	2B		+	addition sign
44	54	2C		,	comma
45	55	2D		-	subtraction sign
46	56	2E		.	period
47	57	2F		/	right slash
48	60	30		0	
49	61	31		1	
50	62	32		2	
51	63	33		3	
52	64	34		4	
53	65	35		5	
54	66	36		6	
55	67	37		7	
56	70	38		8	
57	71	39		9	
58	72	3A		:	colon
59	73	3B		;	semicolon
60	74	3C		<	less than
61	75	3D		=	equal
62	76	3E		>	greater than
63	77	3F		?	question mark
64	100	40		@	at sign
65	101	41		A	
66	102	42		B	
67	103	43		C	

Decimal Value	Octal Value	Hex Value	Control Character	ASCII Symbol	Meaning
68	104	44		D	
69	105	45		E	
70	106	46		F	
71	107	47		G	
72	110	48		H	
73	111	49		I	
74	112	4A		J	
75	113	4B		K	
76	114	4C		L	
77	115	4D		M	
78	116	4E		N	
79	117	4F		O	
80	120	50		P	
81	121	51		Q	
82	122	52		R	
83	123	53		S	
84	124	54		T	
85	125	55		U	
86	126	56		V	
87	127	57		W	
88	130	58		X	
89	131	59		Y	
90	132	5A		Z	
91	133	5B		[left bracket
92	134	5C		\	left slash, backslash
93	135	5D]	right bracket
94	136	5E		^	hat, circumflex, caret
95	137	5F		_	underscore
96	140	60		`	grave accent
97	141	61		a	
98	142	62		b	
99	143	63		c	
100	144	64		d	
101	145	65		e	
102	146	66		f	
103	147	67		g	
104	150	68		h	
105	151	69		i	
106	152	6A		j	
107	153	6B		k	

Decimal Value	Octal Value	Hex Value	Control Character	ASCII Symbol	Meaning
108	154	6C		l	
109	155	6D		m	
110	156	6E		n	
111	157	6F		o	
112	160	70		p	
113	161	71		q	
114	162	72		r	
115	163	73		s	
116	164	74		t	
117	165	75		u	
118	166	76		v	
119	167	77		w	
120	170	78		x	
121	171	79		y	
122	172	7A		z	
123	173	7B		{	left brace
124	174	7C			logical or, vertical bar
125	175	7D		}	right brace
126	176	7E		~	similar, tilde
127	177	7F		DEL	delete

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The following standards are supported:

- The C language is consistent with the International Standard for Information Systems-Programming Language C (ISO/IEC 9899-1999 (E)).
- The C language is consistent with the OpenMP C and C++ Application Programming Interface, Version 2.0.



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