Experiences benchmarking and optimizing GTC on High Performance Computers

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The Gyrokinetic Toroidal Code

- 3D particle-in-cell code to study microturbulence in magnetically confined fusion plasmas.
- Solves the gyro-averaged Vlasov equation.
- Gyrokinetic Poisson equation solved in real space.
- Low noise δf method.
- Global code (full torus as opposed to only a flux tube).
- Massively parallel: typical runs done on 1024 processors.
- Electrostatic approximation with adiabatic electrons.
- Nonlinear and fully self-consistent.
- Written in Fortran 90/95
- Well optimized for superscalar processors







- Particles sample distribution function.
- The particles interact via a grid, on which the potential is calculated from deposited charges.



The PIC Steps

- "SCATTER", or deposit, charges on the grid (nearest neighbors)
- Solve Poisson equation
- "GATHER" forces on each particle from potential
- Move particles (PUSH)
- Repeat...



Charge Deposition for charged rings: 4-point average method

Point-charge particles replaced by charged rings due to gyro-averaging

Charge Deposition Step (SCATTER operation)



Classic PIC

4-Point Average GK (W.W. Lee)



- Fast particle motion along the magnetic field lines leads to a quasi-2D structure in the electrostatic potential
- Poisson equation needs only to be solved on 2D poloidal plane





GTC mesh and geometry



Field-line following coordinates



Original parallel model in GTC: 1D toroidal domain decomposition

- Uses Message Passing Interface (MPI)
- Each MPI process holds a toroidal section
- Most of the communications due to particles moving in and out of the toroidal domains (10% of particles at each time step)
- Efficient "ring-type" communication when moving particles.
- Scales perfectly but limited to about 64 or 128 domains due to (Landau) damping of shorter wavelength modes.







Y-axis: the number of particles (in milions) which move 1 step in 1 second



- The arrival of the IBM SP Power 3 Seaborg at NERSC opened new possibilities for higher performance.
- First step: port GTC from the T3E to the SP and optimize single processor performance
 - Larger memory allowed us to reuse calculations done in the charge deposition subroutine
- The Symmetric Multi-Processing (SMP) nodes of the IBM SP gave an easy path to higher concurrency for GTC: Shared memory programming
- With 16 processors per node, Mixed-model MPI+OpenMP would allow GTC to run on 1,024 processors instead of only 64



New level of parallelism in GTC: Loop-level





- VERY EASY TO IMPLEMENT...
- Although one has to watch out for potential conflicts between threads (processors) trying to write to the same memory location at the same time
 - Easily solved by using thread-private copies of conflicting arrays
- 85% of the work in GTC reside in 4 loops over the number of particles on each MPI process.
- Adding the other loops pushes the amount of computational work in parallel loops beyond 90%.
- The bigger the loops (problem size), the more efficient is the calculation (we saw 98% on large simulations).



OpenMP example of loop-level parallelism

• Simple but powerful OpenMP directives

```
!$omp parallel do private(psitmp,thetatmp,zetatmp,weight,&
!$omp&rhoi,r,ip,jt, ipjt,wz1,kk,wz0,larmor,rdum,ii,wp1,wp0,&
!$omp& tflr,im,tdum,j00,wt10,wt00,j01,wt11,wt01,ij)
    do m=1,mp
        psitmp=phase(1,m)
        thetatmp=phase(2,m)
        zetatmp=phase(2,m)
        weight=phase(3,m)
        weight=phase(5,m)
        rhoi=phase(6,m)*g_inv
        ...
```

enddo



- With mixed-model a single MPI process is assigned to each SMP node on Seaborg
 - Large amount of memory per MPI process (32 GB/proc!)
 - Had to wait for 64-bit MPI to access it though...
- Allowed size scaling study of turbulent transport in tokamaks, including ITER size:
 - 1 billion particles
 - 125 million grid points
 - 1,024 processors
 - largest GTC run at the time





Interesting benchmark of OpenMP on IBM SP and SGI Origin 2000

- SGI O2k has really only 2 processors that share local memory symmetrically.
- The NUMA architecture performs poorly unless processor placement is used.
- The symmetric memory access for the processors on the IBM SP node is ideally adapted to the mixed-model algorithm.





Seaborg allows GTC to routinely run on 1000+ processors



Y-axis: the number of particles (in milions) which move 1 step in 1 second



- Newer, bigger, and faster computers continuously emerge.
- The 2002 record-breaking performance of the Earth Simulator vector computer took everybody by surprise.
- It prompted a renewed interest in vector processing.
- Cray introduced the X1 vector machine soon after.
- I was invited to participate in a study of modern vector architectures compared to current superscalar ones such as the IBM SP.
- The study was lead by Dr Leonid Oliker of the Future Technologies Group at LBL.



- Started on the single-node NEC SX-6 at ARSC
- Porting GTC was very easy although the first tests on a single processor gave a very low performance
- Real work starts: profiling, vectorizing, optimizing, test, and... repeat several times
- Multi-processor optimization done on to the Earth Simulator and CRAY X1





- The charge deposition step (scatter operation) writes to the charge accumulation array in a random fashion (particle positions are random), producing dependencies and memory conflicts whenever 2 or more particles have a common neighboring grid point \rightarrow this prevents vectorization
- In 1D, the charge deposition step with linear interpolation looks like this:

```
do i=1,nparticles
```

```
x = particle_position(i) Indirect addressing!
ix_grid = int(x) Potential Conflicts
dx = x - real(ix_grid)
charge(ix_grid) = charge(ix_grid)+q*(1-dx)
charge(ix_grid+1) = charge(ix_grid+1)+q*dx
end do
```



Avoiding memory dependencies: The work-vector method (Nishiguchi '85)

Example of loop with indirect addressing similar to charge deposition:

```
DO i=1,np
    charge(ix(i))=charge(ix(i)) + q(i)
END DO
```

Fully vectorizable loop using multiple copies (vector length of 256):

```
ALLOCATE(charge_tmp(256,ngrid)) Uses 256*ngrid*sizeof(charge_tmp)
DO i=1,np,256 of extra memory! (can be 1GB)
DO j=1,min(256,np-i+1)
charge_tmp(j,ix(i+j-1))=charge_tmp(j,ix(i+j-1)) + q(i+j-1)
END DO
END DO
END DO
DO i=1,256
DO ig=1,ngrid
charge(ig)=charge(ig) + charge_tmp(i,igrid)
END DO
END DO
END DO
```



- Each Earth Simulator node has 8 vector processors sharing 16 GBytes of memory, allowing us to use GTC's mixed-model MPI+OpenMP.
- However, loop-level work splitting with OpenMP reduces the number of loop operations, which in turn degrades vector efficiency → Lower performance
- Charge deposition loop with OpenMP requires private copies of the grid array for each processor on the node.
- Combined with the 256 copies of the same grid array needed for vectorization, the loop-level OpenMP requires too much memory.



Cache-less memory access issues on the SX-6 and ES

- Better memory access is the secret to higher performance
- True for STORING to memory as well as FETCHING from it!

```
do m=1,mi
    psitmp=zion(1,m)
    thetatmp=zion(2,m)
    zetatmp=zion(3,m)
    rhoi=zion(6,m)*smu_inv
    r=sqrt(2.0*psitmp)
    ip=max(0,min(mpsi,int((r-a0)*delr+0.5)))
    jt=max(0,min(mtheta(ip),int(thetatmp*pi2_inv*delt(ip)+0.5)))
    ipjt=igrid(ip)+jt
    wz1=(zetatmp-zetamin)*delz
    ...
Repeatedly accessing the same
memory bank before the bank busy
time is over from the last access
leads to poor memory performance!
```

Duplicate small arrays like "igrid" and "mtheta": !\$duplicate 37% improvement on chargei, but uses even more memory...



Vector performance of main routines on the Earth Simulator

ORIGINAL CODE BEFORE MODIFICATIONS:

PROG.UNIT	EXCLUSIVE TIME[sec](%)	MFLOPS	V.OP RATIO	AVER. V.LEN	BANK CONF
chargei	282.677(54.4)	62.0	0.65	98.1	0.0000
pushi	125.211(24.1)	320.1	67.51	196.8	4.3336
poisson	57.878(11.1)	418.9	94.26	107.2	0.3158

Note: the 2 tests do not have the same number of time steps so the times are different

CODE AFTER MODIFICATIONS TO CHARGEI, PUSHI, POISSON:

PROG.UNIT	EXCLUSIVE TIME[sec](%)	MFLOPS	V.OP RATIO	AVER. V.LEN	BANK CONF
chargei	89.924(33.3)	1314.3	99.65	248.1	6.5002
pushi	93.877(34.7)	2426.6	99.38	255.9	8.8139
poisson	26.239(9.7)	918.1	99.71	252.7	3.2485

Total = 1.412 Gflops per proc



- Must deal with multi-streaming on top of vectorization
- Same vectorizations apply.
- Easier to prevent vectorization of small inner loops
- Also needs the work-vector method with the same dimensions of 256 in MSP mode:
 - 4 streams x 64 (vector length)
 - Uses as much extra memory as the Earth Simulator
- Unvectorized and unstreamed loop in "shifti" slows down the calculation to a crawl
 - 54% of the time spent in that routine according to "pat"
 - Was only 11% on the ES



• "Unstreamed" and "unvectorized" loop due to nested if blocks:

```
do m=m0,mi
   zetaright=min(2.0*pi,zion(3,m))-zetamax
   zetaleft=zion(3,m)-zetamin
   if( zetaright*zetaleft > 0 )then
      zetaright=zetaright*0.5*pi_inv
      zetaright=zetaright-real(floor(zetaright))
      msend=msend+1
      kzi(msend)=m
      if( zetaright < 0.5 )then
         msendright(1)=msendright(1)+1
         iright(msendright(1))=m
      else
         msendleft(1)=msendleft(1)+1
         ileft(msendleft(1))=m
      endif
   endif
enddo
```



New loop in shift

```
!dir$ preferstream
    do imm=1,4
!dir$ prefervector
        do m = (imm-1) * mi/4 + 1, imm * mi/4
           zetaright=min(2.0*pi,zion(3,m))-zetamax
           zetaleft=zetamin-zion(3,m)
           alpha=pi2*aint(1.0-pi4 inv*zetaleft)
           beta=pi2*aint(1.0-pi4 inv*zetaright)
           kappa=pi2*aint(1.0+zetaleft*zetaright*pi2sg inv)
           aright=(alpha+zetaleft) - (beta+zetaright) - kappa
           aleft=(alpha+zetaleft) - (beta+zetaright) + kappa
           if( aright > 0.0 )then
              msend r(imm)=msend r(imm)+1
              kzi r(msend r(imm),imm)=m
           endif
           if( aleft < 0.0 )then
              msend l(imm)=msend l(imm)+1
              kzi l(msend l(imm),imm)=m
           endif
        enddo
    enddo
```



- Yes, the overall time spent in shifti went from 54% to only 4%!!
- On the Earth Simulator, the compiler can only deal with a single conditional statement within a loop in order to vectorize that loop. Solution: split the loop in 2 parts.
- Improved performance on the ES but not as dramatic as on the X1.



Memory used by the vectorized version of GTC (per processor)



- For micell=10 memory on the ES is up to 8 times more than one the Power 3!
- It gets better as the number of particles per cell increases



Flops/sec for higher particle resolution





- The algorithm to "push" the particles is very efficient on vector machines.
- This allows us to run high resolution simulations using a large number of particles.
- More particles means
 - more phase space resolution (velocity + configuration space)
 - lower discrete particle noise/fluctuations
 - longer simulations
 - higher efficiency on vector computers
- Without using OpenMP, we are back to only 1D domain decomposition and a max of 64-128 processors
- How to improve concurrency?



New parallel model: Domain decomposition + particle splitting

- 1D Domain decomposition:
 - Several MPI processes can now sit a section of the torus
- Particle splitting method
 - The particles in a toroidal section a: equally divided between several MPI processes
- Particles randomly distributed between processors within a toroidal domain.
- <u>No OpenMP</u>
- Pure MPI version







- New large scale computers such as Blue Gene/L and Cray XT3 allow only message passing for communication between processors.
- The MPI-only version of GTC has been very successful on those platforms.
- It achieved the highest performance of 7.2 TFlops on the Earth Simulator using 4,096 processors.
- Used over 16,000 processors on the Blue Gene/L computer at IBM Watson.
- Largest GTC production simulation recently carried out on 4,800 processors of the Cray XT3 at ORNL using 28 billion particles.



Latest benchmark: weak scaling study with fixed device size





- Benchmarking and optimizing work never ends
- New platforms with more processors and different characteristics are continuously being developed.
- To cope with the changes while wanting to achieve top performance, codes must be flexible and developers must be willing to modify their codes.
- Secret to high performance (in my opinion...)
 - Data access (must feed the processor as fast as possible)
 - Fast data access = good data layout
 - Minimize communications
- Speed is not a substitute to "right answer"...