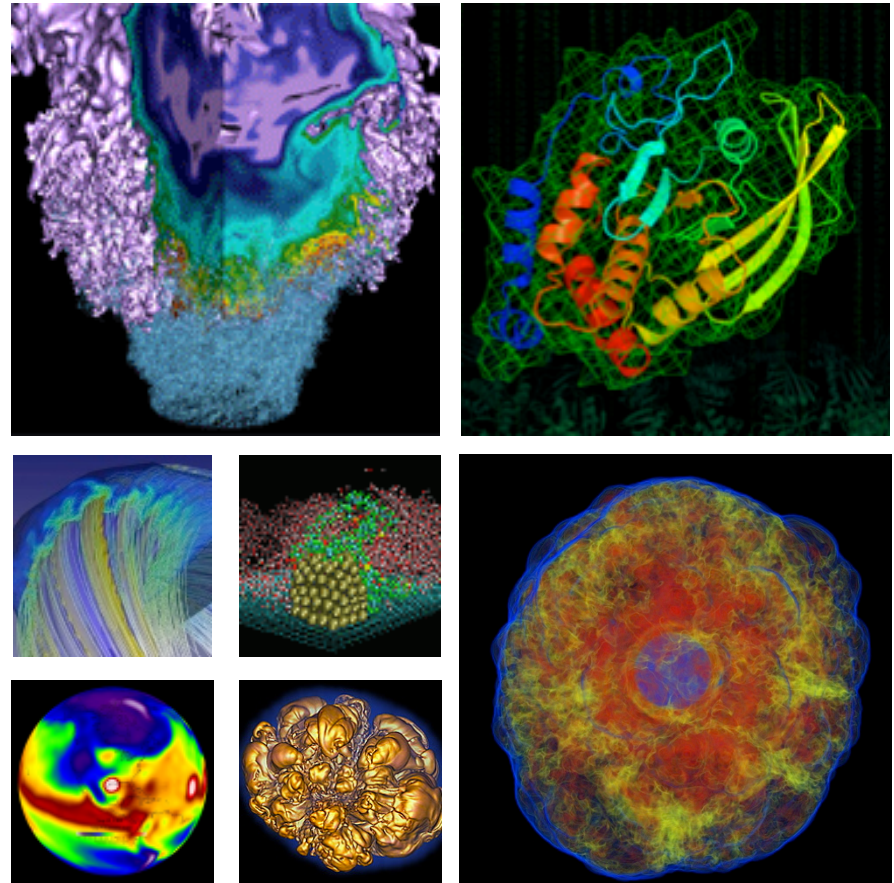


# Introduction to Intel Xeon Phi (“Knights Landing”) on Cori



**Brandon Cook**  
**Brian Friesen**

2017 June 9

# Knights Landing is here!



- **KNL nodes installed in Cori in 2016**
- **“Pre-production” for ~ 1 year**
  - No charge for using KNL nodes
  - Limited access (until now!)
  - Intermittent downtime
  - Frequent software changes
- **KNL nodes on Cori will soon enter production**
  - Charging begins 2017 July 1

# Knights Landing overview



## Knights Landing: Next Intel® Xeon Phi™ Processor

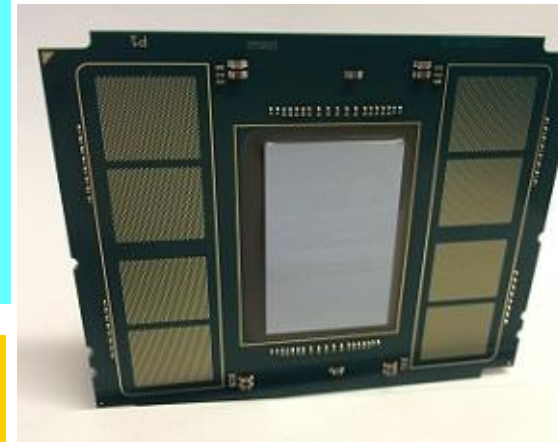
Intel® Many-Core Processor targeted for HPC and Supercomputing

First **self-boot** Intel® Xeon Phi™ processor that is **binary compatible** with main line IA. Boots standard OS.

**Significant improvement in scalar and vector performance**

Integration of **Memory on package**: innovative memory architecture for high bandwidth and high capacity

Integration of **Fabric on package**



### Three products

**KNL Self-Boot**  
(Baseline)

**KNL Self-Boot w/ Fabric**  
(Fabric Integrated)

**KNL Card**  
(PCIe-Card)

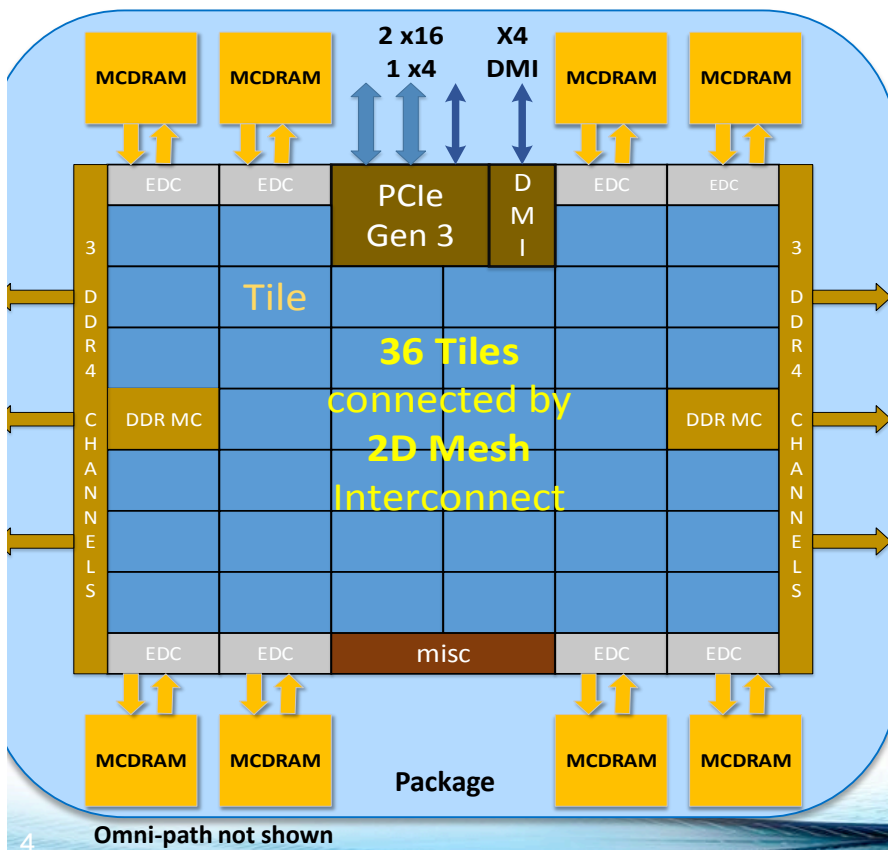
Potential future options subject to change without notice.

All timeframes, features, products and dates are preliminary forecasts and subject to change without further notification.

# Knights Landing overview



## Knights Landing Overview



### TILE

2 VPU	CHA	2 VPU
Core	1MB L2	Core

**Chip: 36 Tiles** interconnected by 2D Mesh

**Tile: 2 Cores + 2 VPU/core + 1 MB L2**

**Memory: MCDRAM:** 16 GB on-package; High BW  
**DDR4:** 6 channels @ 2400 up to 384GB

**IO:** 36 lanes PCIe Gen3. 4 lanes of DMI for chipset

**Node:** 1-Socket only

**Fabric:** Omni-Path on-package (not shown)

**Vector Peak Perf:** 3+TF DP and 6+TF SP Flops

**Scalar Perf:** ~3x over Knights Corner

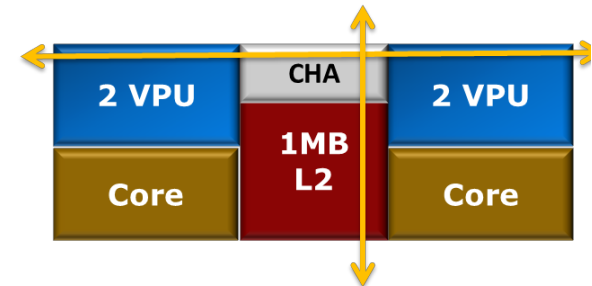
**Streams Triad (GB/s):** MCDRAM : 400+; DDR: 90+

Source Intel: All products, computer systems, dates and figures specified are preliminary based on current expectations, and are subject to change without notice. KNL data are preliminary based on current expectations and are subject to change without notice. <sup>1</sup>Binary Compatible with Intel Xeon processors using Haswell Instruction Set (except TSX). <sup>2</sup>Bandwidth numbers are based on STREAM-like memory access pattern when MCDRAM used as flat memory. Results have been estimated based on internal Intel analysis and are provided for informational purposes only. Any difference in system hardware or software design or usage scenario may affect actual performance.

# Knights Landing overview



**KNL Tile:** 2 Cores, each with 2 VPU  
1M L2 shared between two Cores



**Core:** Changed from Knights Corner (KNC) to KNL. Based on 2-wide OoO Silvermont™ Microarchitecture, but with many changes for HPC.

4 thread/core. Deeper OoO. Better RAS. Higher bandwidth. Larger TLBs.

**2 VPU:** 2x AVX512 units. 32SP/16DP per unit. X87, SSE, AVX1, AVX2 and EMU

**L2:** 1MB 16-way. 1 Line Read and ½ Line Write per cycle. Coherent across all Tiles

**CHA:** Caching/Home Agent. Distributed Tag Directory to keep L2s coherent. MESIF protocol. 2D-Mesh connections for Tile

# Knights Landing overview



## Many Trailblazing Improvements in KNL

Improvements	What/Why
Self Boot Processor	No PCIe bottleneck
Binary Compatibility with Xeon	Runs all legacy software. No recompilation.
New Core: Atom™ based	~3x higher ST performance over KNC
Improved Vector density	3+ TFLOPS (DP) peak per chip
New AVX 512 ISA	New 512-bit Vector ISA with Masks
Scatter/Gather Engine	Hardware support for gather and scatter
New memory technology: MCDRAM + DDR	Large High Bandwidth Memory → MCDRAM Huge bulk memory → DDR
New on-die interconnect: Mesh	High BW connection between cores and memory
Integrated Fabric: Omni-Path	Better scalability to large systems. Lower Cost

Software and workloads used in performance tests may have been optimized for performance only on Intel microprocessors. Performance tests, such as SYSmark and MobileMark, are measured using specific computer systems, components, software, operations and functions. Any change to any of those factors may cause the results to vary. You should consult other information and performance tests to assist you in fully evaluating your contemplated purchases, including the performance of that product when combined with other products. For more information go to <http://www.intel.com/performance>. Results have been estimated based on internal Intel analysis and are provided for informational purposes only. Any difference in system hardware or software design or configuration may affect actual performance.

## What is different about Cori?

### Edison (Xeon E7-4680 v2; “Ivy Bridge”):

- 5576 nodes
- 12 physical cores/socket
- 24 logical cores/socket
- 2.4 - 3.2 GHz
- 8 double precision operations per cycle
- 2.5 GB DRAM/core (64 GB/node)
- ~100 GB/s memory bandwidth

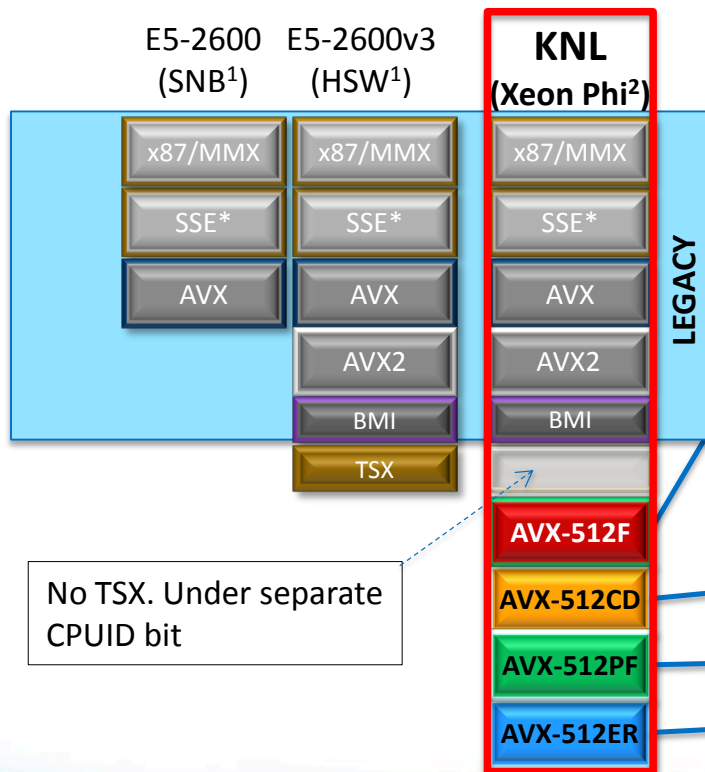
### Cori (Xeon Phi 7250; “Knights Landing”):

- 9688 nodes
- 68 physical cores/socket
- 272 logical cores/socket
- 1.2-1.4 GHz (up to 1.6 GHz in Turbo Mode)
- 32 double precision operations per cycle
- 16 GB of fast memory (MCDRAM)  
96GB of DDR4 DRAM
- ~115 GB/s DDR4 bandwidth
- ~400 GB/s MCDRAM bandwidth

# Knights Landing overview



## KNL ISA



No TSX. Under separate CPUID bit

### KNL implements all legacy instructions

- Legacy binary runs w/o recompilation
- KNC binary requires recompilation

### KNL introduces AVX-512 Extensions

- 512-bit FP/Integer Vectors
- 32 registers, & 8 mask registers
- Gather/Scatter

**Conflict Detection**: Improves Vectorization

**Prefetch**: Gather and Scatter Prefetch

**Exponential and Reciprocal** Instructions

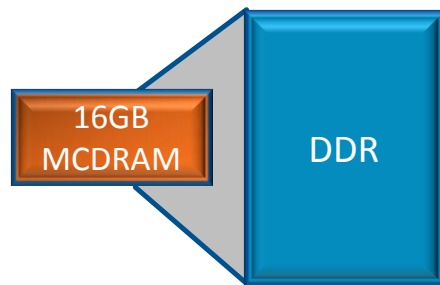
<sup>1</sup> Previous Code name Intel® Xeon® processors  
<sup>2</sup> Xeon Phi = Intel® Xeon Phi™ processor



## Memory Modes

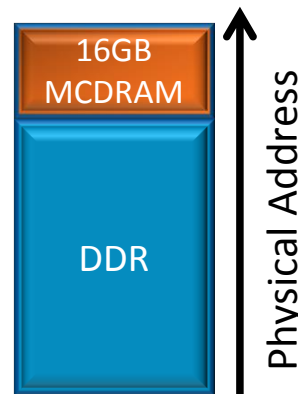
Three Modes. Selected at boot

### Cache Mode



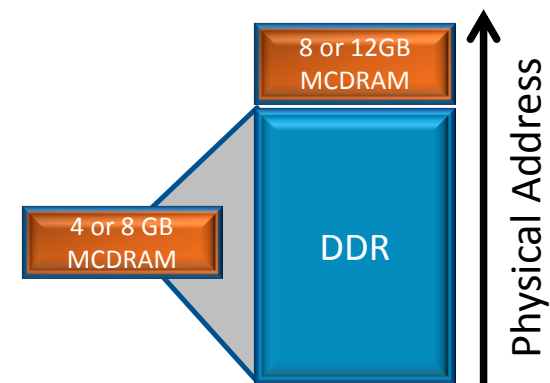
- SW-Transparent, Mem-side cache
- Direct mapped. 64B lines.
- Tags part of line
- Covers whole DDR range

### Flat Mode



- MCDRAM as regular memory
- SW-Managed
- Same address space

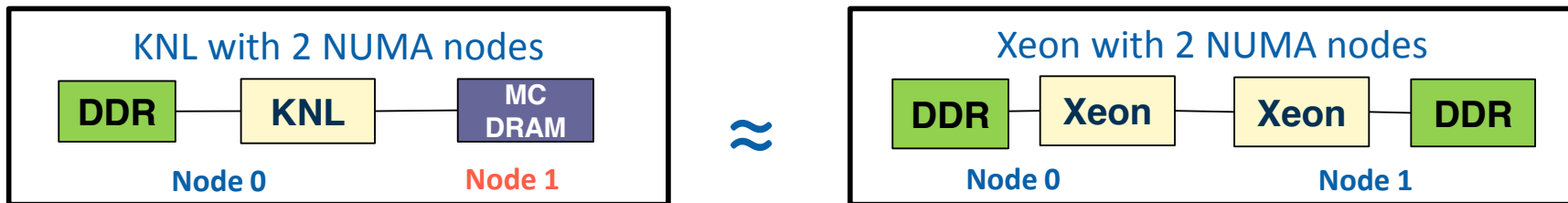
### Hybrid Mode



- Part cache, Part memory
- 25% or 50% cache
- Benefits of both

## Flat MCDRAM: SW Architecture

MCDRAM exposed as a separate NUMA node



Memory allocated in DDR by default → Keeps non-critical data out of MCDRAM.

Apps explicitly allocate critical data in MCDRAM. Using two methods:

- “**Fast Malloc**” functions in High BW library (<https://github.com/memkind>)
  - Built on top to existing *libnuma* API
- “**FASTMEM**” Compiler Annotation for Intel Fortran

**Flat MCDRAM with existing NUMA support in Legacy OS**

## Flat MCDRAM SW Usage: Code Snippets

### C/C++ ([\\*https://github.com/memkind](https://github.com/memkind))

Allocate into DDR

```
float *fv;  
fv = (float *)malloc(sizeof(float)*100);
```



Allocate into MCDRAM

```
float *fv;  
fv = (float *)hbw_malloc(sizeof(float) * 100);
```

### Intel Fortran

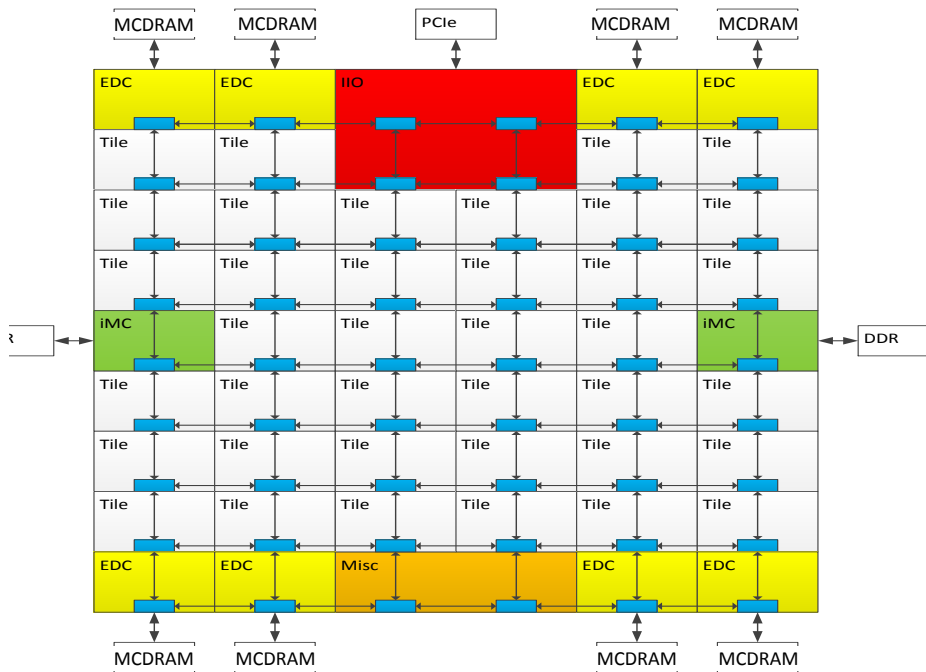
Allocate into MCDRAM

```
c Declare arrays to be dynamic  
REAL, ALLOCATABLE :: A(:)  
  
!DEC$ ATTRIBUTES, FASTMEM :: A  
  
NSIZE=1024  
c allocate array 'A' from MCDRAM  
c  
ALLOCATE (A(1:NSIZE))
```

# Knights Landing overview



## KNL Mesh Interconnect



### Mesh of Rings

- Every row and column is a (half) ring
- YX routing: Go in Y → Turn → Go in X
- Messages arbitrate at injection and on turn

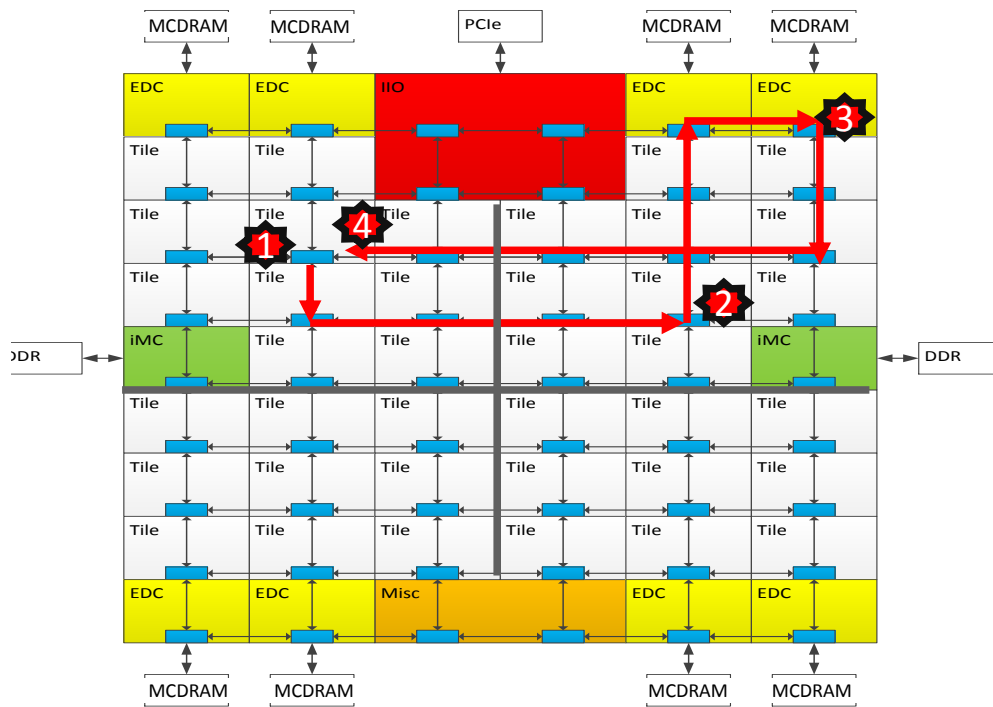
### Cache Coherent Interconnect

- MESIF protocol (F = Forward)
- Distributed directory to filter snoops

### Three Cluster Modes

- (1) All-to-All (2) Quadrant (3) Sub-NUMA Clustering

## Cluster Mode: Quadrant



Chip divided into four virtual Quadrants

Address hashed to a Directory in the same quadrant as the Memory

Affinity between the Directory and Memory

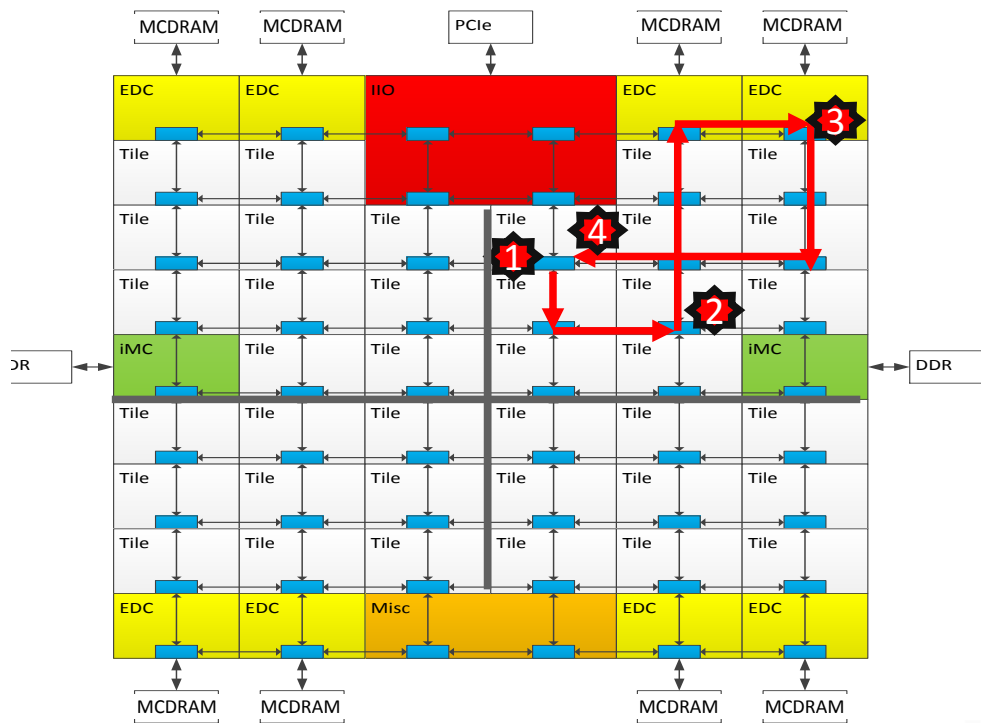
Lower latency and higher BW than all-to-all. SW Transparent.

1) L2 miss, 2) Directory access, 3) Memory access, 4) Data return

# Knights Landing overview



## Cluster Mode: Sub-NUMA Clustering (SNC)



Each Quadrant (Cluster) exposed as a separate NUMA domain to OS.

Looks analogous to 4-Socket Xeon

Affinity between Tile, Directory and Memory

Local communication. Lowest latency of all modes.

SW needs to NUMA optimize to get benefit.

1) L2 miss, 2) Directory access, 3) Memory access, 4) Data return

16

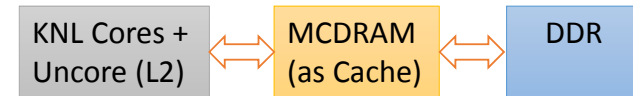
# Multi-channel DRAM overview



## MCDRAM Modes

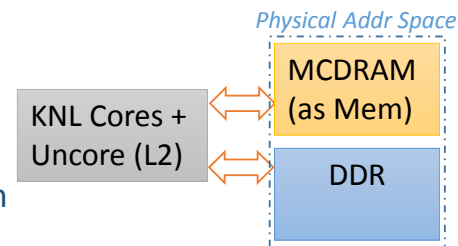
### Cache mode

- No source changes needed to use
- Misses are expensive (higher latency)
  - Needs MCDRAM access + DDR access



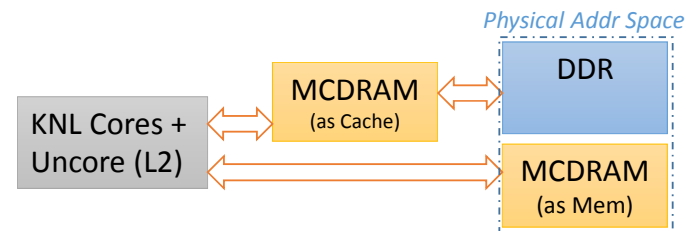
### Flat mode

- MCDRAM mapped to physical address space
- Exposed as a NUMA node
  - Use `numactl --hardware, lscpu` to display configuration
- Accessed through `memkind` library or `numactl`



### Hybrid

- Combination of the above two
  - E.g., 8 GB in cache + 8 GB in Flat Mode

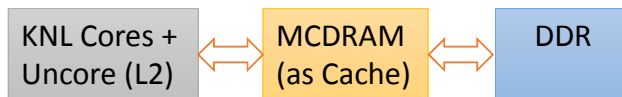


# Multi-channel DRAM overview



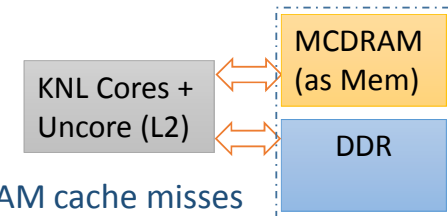
## MCDRAM as Cache

- Upside
  - No software modifications required
  - Bandwidth benefit (over DDR)
- Downside
  - Higher latency for DDR access
    - i.e., for cache misses
  - Misses limited by DDR BW
  - All memory is transferred as:
    - DDR -> MCDRAM -> L2
  - Less addressable memory



## MCDRAM as Flat Mode

- Upside
  - Maximum BW
  - Lower latency
    - i.e., no MCDRAM cache misses
  - Maximum addressable memory
  - Isolation of MCDRAM for high-performance application use only
- Downside
  - Software modifications (or interposer library) required
    - to use DDR and MCDRAM in the same app
  - Which data structures should go where?
  - MCDRAM is a finite resource and tracking it adds complexity





## Accessing MCDRAM in Flat Mode

- Option A: Using numactl
  - Works best if the whole app can fit in MCDRAM
- Option B: Using libraries
  - Memkind Library
    - Using library calls or Compiler Directives (Fortran\*)
    - Needs source modification
  - AutoHBW (interposer library based on memkind)
    - No source modification needed (based on size of allocations)
    - No fine control over *individual* allocations
- Option C: Direct OS system calls
  - mmap(1), mbind(1)
  - Not the preferred method
    - Page-only granularity, OS serialization, no pool management

\*Other names and brands may be claimed as the property of others.

## Option A: Using numactl to Access MCDRAM

- MCDRAM is exposed to OS/software as a NUMA node
- Utility `numactl` is standard utility for NUMA system control
  - See “man numactl”
  - Do “numactl --hardware” to see the NUMA configuration of your system
- If the total memory footprint of your app is smaller than the size of MCDRAM
  - Use numactl to allocate all of its memory from MCDRAM
  - `numactl --membind=mcdram_id <your_command>`
    - Where `mcdram_id` is the ID of MCDRAM “node”
- If the total memory footprint of your app is larger than the size of MCDRAM
  - You can still use numactl to allocate *part* of your app in MCDRAM
    - `numactl --preferred=mcdram_id <your_command>`
      - Allocations that don’t fit into MCDRAM spills over to DDR
    - `numactl --interleave=nodes <your_command>`
      - Allocations are interleaved across all *nodes*

## Option B.1: Using Memkind Library to Access MCDRAM

### Allocate 1000 floats from DDR

```
float   *fv;  
fv = (float *)malloc(sizeof(float) * 1000);
```

### Allocate 1000 floats from MCDRAM

```
#include <hbwmalloc.h>  
float   *fv;  
fv = (float *)hbw_malloc(sizeof(float) * 1000);
```

### Allocate arrays from MCDRAM and DDR in Intel® Fortran Compiler

```
c   Declare arrays to be dynamic  
   REAL, ALLOCATABLE :: A(:), B(:), C(:)  
  
   !DEC$ ATTRIBUTES FASTMEM :: A  
  
   NSIZE=1024  
c  
c   allocate array 'A' from MCDRAM  
c  
   ALLOCATE (A(1:NSIZE))  
c  
c   Allocate arrays that will come from DDR  
c  
   ALLOCATE (B(NSIZE), C(NSIZE))
```

## Option B.2: AutoHBW

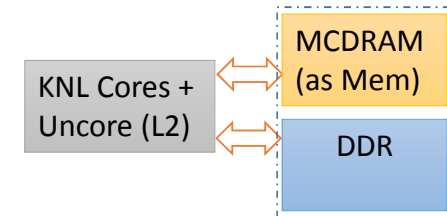
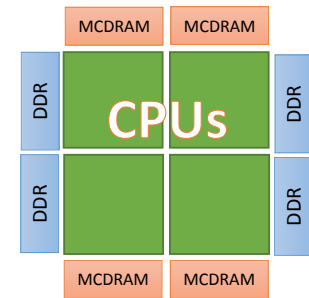
- AutoHBW: Interposer Library that comes with memkind
  - Automatically allocates memory from MCDRAM
    - If a heap allocation (e.g., malloc/calloc) is larger than a given threshold
- Demo
  - see `/examples/autohbw_test.sh`
  - Run `gpp` with AutoHBW
- Environment variables (see [autohbw\\_README](#))
  - `AUTO_HBW_SIZE=x[:y]`
  - `AUTO_HBW_LOG=level`
  - `AUTO_HBW_MEM_TYPE=memory_type` *# useful for interleaving*
- Finding source locations of arrays
  - `export AUTO_HBW_LOG=2`
  - `./app_name > log.txt`
  - `autohbw_get_src_lines.pl log.txt app_name`

# Multi-channel DRAM overview



## Advanced Topic: MCDRAM in SNC4 Mode

- SNC4: Sub-NUMA Clustering
  - KNL die is divided into 4 clusters (similar to a 4-Socket Haswell)
    - SNC4 configured at boot time
    - Use `numactl --hardware` to find out nodes and distances
      - There are **4 DDR (+CPU) nodes + 4 MCDRAM (no CPU) nodes**, in flat mode
- Running 4-MPI ranks is the easiest way to utilize SNC4
  - Each rank allocates from closest DDR node
    - If a rank allocates MCDRAM, it goes to closest MCDRAM node
- If you run only 1 MPI rank and use `numactl` to allocate on MCDRAM
  - Specify all MCDRAM nodes
    - E.g., `numactl -m 4,5,6,7`

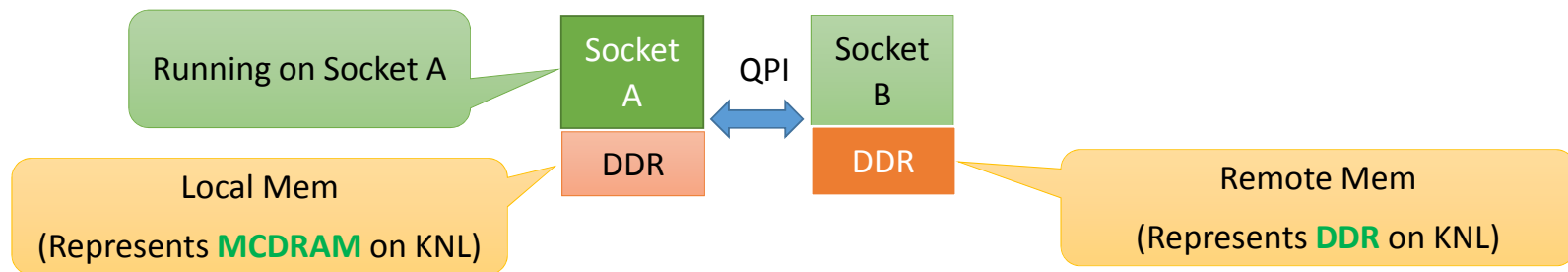


*Compare with 2 NUMA nodes*

## Observing MCDRAM Memory Allocations

- Where is MCDRAM usage printed?
  - `numastat -m`
    - Printed for each NUMA node
    - Includes Huge Pages info
  - `numastat -p <pid>` OR `numastat -p exec_name`
    - Info about process `<pid>`
    - E.g., `watch -n 1 -p numastat exec_name`
  - `cat /sys/devices/system/node/node*/meminfo`
    - Info about each NUMA node
  - `cat /proc/meminfo`
    - Aggregate info for system
- Utilities that provide MCDRAM node info
  - `<memkind_install_dir>/bin/memkind-hbw-nodes`
  - `numactl --hardware`
  - `lscpu`

## MCDRAM Emulation Demo



- Use cores on socket A and DDR on socket B (remote memory)
  - `export MEMKIND_HBW_NODES=0 # HBW allocation go to DDR on Socket A`
  - `numactl --membind=1 --cpunodebind=0 <your_command>`
- Any HBW allocations will allocate DDR on socket A
  - Accesses to DDR on socket A (local memory) has higher BW
    - Also has lower latency (which is an inaccuracy)



**National Energy Research Scientific Computing Center**